Week 10: Lecture B Network Denial of Service

Thursday, October 31, 2024



Announcements

Project 3: WebSec released

Deadline: Thursday, November 7th by 11:59PM (next week)

Project 3: Web Security

Deadline: Thursday, November 7 by 11:59PM.

Before you start, review the course syllabus for the Lateness, Collaboration, and Ethical Use policies.

You may optionally work alone, or in teams of **at most two** and submit **one project per team**. If you have difficulties forming a team, post on **Piazza's Search for Teammates** forum. Note that the final exam will cover project material, so you and your partner should collaborate on each part.

The code and other answers your group submits must be entirely your own work, and you are bound by the University's Student Code. You may consult with other students about the conceptualization of the project and the meaning of the questions, but you may not look at any part of someone else's solution or collaborate with anyone outside your group. You may consult published references, provided that you appropriately cite them (e.g., in your code comments). **Don't risk your grade and degree by cheating!**

Complete your work in the **CS 4440 VM**—we will use this same environment for grading. You may not use any **external dependencies**. Use only default Python 3 libraries and/or modules we provide you.

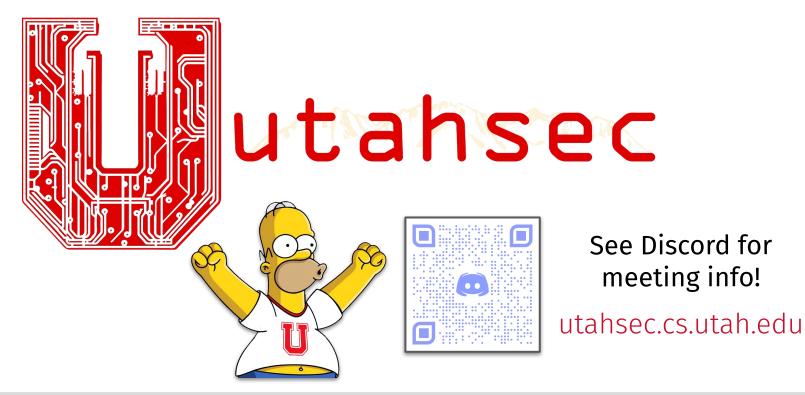
Project 3 progress

Working on Part 1	0%
Finished Part 1, working on Part 2	004
Finished Part 2, working on Part 3	0%
	0%
Finished with everything!	
	0%
Haven't started yet :(0%



Start the presentation to see live content. For screen share software, share the entire screen. Get help at **pollev.com/app**

Announcements



Announcements





Last time on CS 4440...

Application Layer Attacks HTTP Content Injection SMTP Header Spoofing DNS Hijacking



Application Layer Attacks

- Application Layer:
 - ???



Application Layer Attacks

- Application Layer: where network-facing apps send/receive message
 - Application-specific protocols (message semantics, structure, processing rules, etc.)
- Attacking the application layer:
 - ???







Application Layer Attacks

- Application Layer: where network-facing apps send/receive message
 - Application-specific protocols (message semantics, structure, processing rules, etc.)

Attacking the application layer:

- Command Injection
 - SQL injection, CSRF, XSS
- Denial of Service
 - Crash a remote application
 - Prevent others from using it
- Message Tampering / Sniffing
 - Injecting data into messages
 - Capturing unencrypted data
- Other protocol-specific attacks



The HTTP Protocol

What is HTTP?

Helio. Accounts Transfer & Pay	You have the dreams and the funds. You existing home exists for e of roots can fund home improvement projects, large purchases and more.	SCHEDULE AN APPOINTMENT	1
Send Money with Zelle*	DEPOSITS & INVESTMENTS		
	5/3 Essential Checking	Available	
5/3	5/3 Preferred Checking	Available	
	5/3 Essential Checking	hvallable	
(B)	Maxsaver	Available	
vord ©	Roth Ira	ivaliable	
) Remember Me Forgot Login	Ira	Available	
Log In	CREDIT CARDS & LOANS		
	Equity Line	Principal Balance	
NUP 5/3 Find Open Support Find Open			



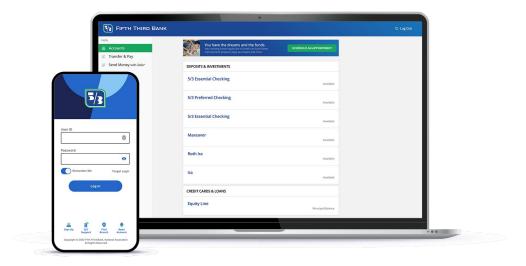
The HTTP Protocol

What is HTTP?

Protocol for transmitting **hypermedia documents** (e.g., web pages)

HTTP's Characteristics:

- Widely used
- Simple
- Unencrypted





HTTP Tampering

Attack: ???



Alter



HTTP Tampering

- Attack: exploit HTTP's insecurity
 - Nothing is encrypted!
- Attacker intercepts requested webpage and modifies it
 - User receives modified webpage
- Attacker capabilities?
 - Inject malicious content
 - Inject malicious code





Thwarting HTTP Injection

Defenses: ???



Thwarting HTTP Injection

• **Defenses:** encrypt everything all the time!

Answer: *completely* ditch **HTTP**!

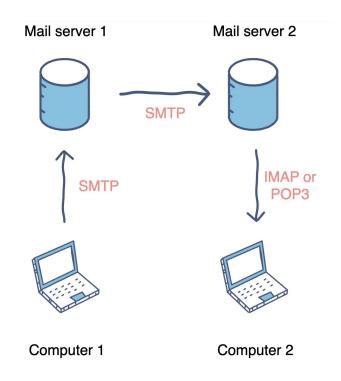
- As web and app developers, enforce strict HTTPS compliance
 - Necessary to prevent HTTPS → HTTP downgrade attacks



- SMTP: Simple Mail Transfer Protocol
 - Implemented in the application layer

Characteristics:

- Text-based
- Connection-oriented
- Uses TCP ports 25/587
- Security guarantees:
 ???





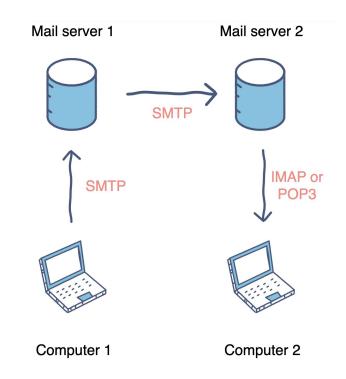
- SMTP: Simple Mail Transfer Protocol
 - Implemented in the application layer

Characteristics:

- Text-based
- Connection-oriented
- Uses TCP ports 25/587

Security guarantees:

- Message integrity—no!
- Confidentiality—no!
- Authentication—no!





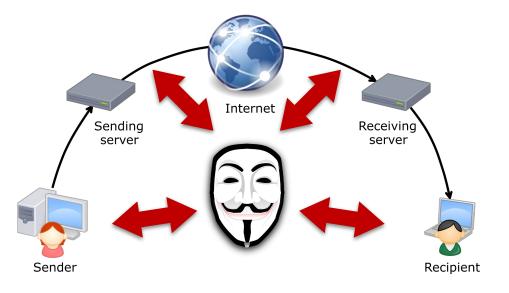
No message integrity
 ???

No confidentiality

· ???

No authentication

???





No message integrity

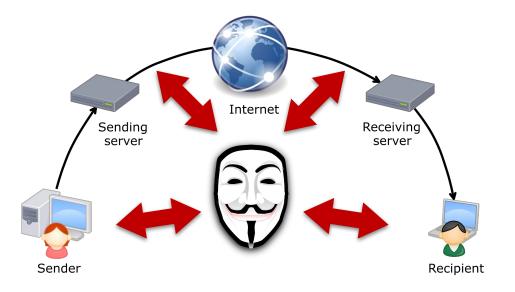
- Tamper with messages
- Block messages

No confidentiality

- Find sender/recipient
- Read message contents

No authentication

Spoof sender identity



SMTP Header Spoofing

Attack: ???

S: 220 attacker.com SMTP Exim C: HELO attacker.com S: 250 Hello attacker.com C: MAIL FROM: <ceo@company.com> S: 250 Ok C: RCPT TO: <bob@company.com> S: 250 Accepted C: DATA S: 354 Enter a message, ending with "." on a line by itself C: Subject: Download this urgently C: From: ceo@company.com C: To: <u>bob@company.com</u> **C:** C: Hi Bob, C: Please download this urgently: https://some-malicious-link.com C: Regards C: . S: 250 OK C: QUIT S: 221 attacker.com closing connection

To: robertbateman@email.com Subject: Hi There From: "Mickey Mouse" <m.mouse@disney.com> X-Priority: 3 (Normal) Importance: Normal Errors-To: m.mouse@disney.com Reply-To: m.mouse@disney.com Content-Type: text/plain



SMTP Header Spoofing

• Attack: spoof SMTP header to **mislead recipient** about sender of the email



```
To: robertbateman@email.com
Subject: Hi There
From: "Mickey Mouse" <m.mouse@disney.com>
X-Priority: 3 (Normal)
Importance: Normal
Errors-To: m.mouse@disney.com
Reply-To: m.mouse@disney.com
Content-Type: text/plain
```



Thwarting Email Spoofing

Defenses: ???



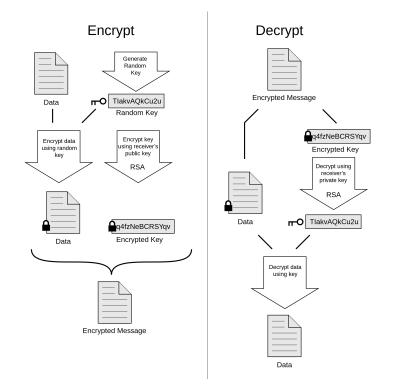
Thwarting Email Spoofing

Checking email bodies

- Included links
- Attached files
- Text analysis (e.g., known spam campaigns)

Checking email headers

- Egress server domain registration
 - Check that sender is who it says it is
- Pretty Good Privacy (PGP)
 - Sender and Receiver authentication
 - Confidentiality
 - Integrity



Identification on the Web

• How do we identify **people**?

- Social security numbers
- Passports, drivers licenses
- Their unique fingerprints
- How can we identify internet hosts?
 - ???

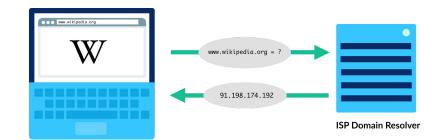
Identification on the Web

- How do we identify people?
 - Social security numbers
 - Passports, drivers licenses
 - Their unique fingerprints
- How can we identify internet hosts?
 - Network layer: location via IP addresses
 - A logical addressing system
 - 32-bit (IPV4) addressing datagrams
 - What you care about: ???



Identification on the Web

- How do we identify **people**?
 - Social security numbers
 - Passports, drivers licenses
 - Their unique fingerprints
- How can we identify internet hosts?
 - Network layer: location via IP addresses
 - A logical addressing system
 - 32-bit (IPV4) addressing datagrams
 - What you care about: the domain name
 - E.g., www.wikipedia.org





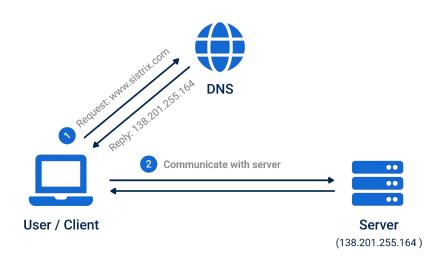


 Distributed database implemented in hierarchy of many name servers

Application-layer protocol:

- Hosts and domain name servers communicate to resolve domain names
 - Address-name translation
- Result: user requests ???
 - But their host really gets ???

- **Distributed database** implemented in hierarchy of many name servers
- Application-layer protocol:
 - Hosts and domain name servers communicate to resolve domain names
 - Address-name translation
- Result: user requests domain name
 - But their host really gets its IP address
 - Convenient!



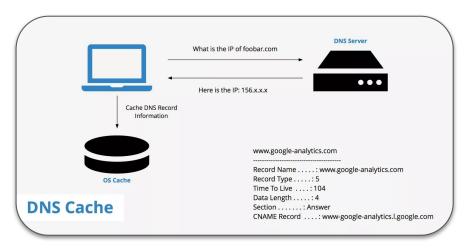
- How can we optimize DNS resolution?
 - ???



How can we optimize DNS resolution?

- Cache look-ups to amortize initial look-up, reduce system load
- Optimization 1: ???

• Optimization 2: ???





How can we optimize DNS resolution?

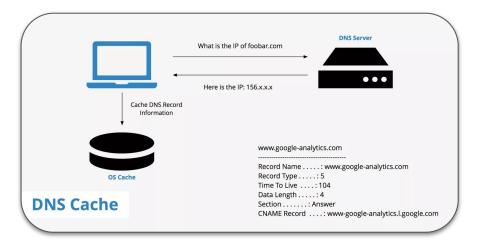
• Cache look-ups to amortize initial look-up, reduce system load

Optimization 1: temporal locality

- www.espn.com/page1
- www.espn.com/page2

Optimization 2: popular domains

- google.com
- Facebook.com



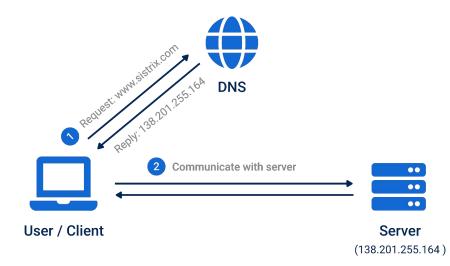


<pre>stefan@cs4440:~\$ time nslookup facebook.com</pre>	<pre>stefan@cs4440:~\$ time nslookup facebook.com</pre>			
Server: 127.0.0.53	Server: 127.0.0.53			
Address: 127.0.0.53#53	Address: 127.0.0.53#53			
Non-authoritative answer:	Non-authoritative answer:			
Name: facebook.com	Name: facebook.com			
Address: 31.13.70.36	Address: 31.13.70.36			
Name: facebook.com	Name: facebook.com			
Address: 2a03:2880:f10d:83:face:b00c:0:25de	Address: 2a03:2880:f10d:83:face:b00c:0:25de			
real 0m0.474s	real 0m0.023s			
user 0m0.000s	user 0m0.000s			
sys 0m0.015s	sys 0m0.011s			

First Lookup (**non-cached**)

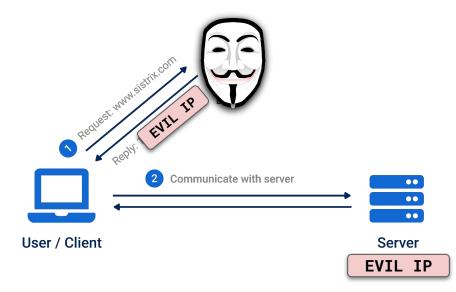
Second Lookup (cached)

What can an attacker do if they control a DNS server?





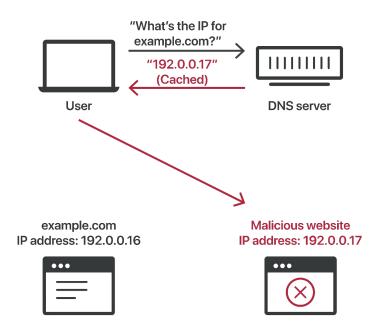
- What can an attacker do if they **control a DNS server**?
 - Control how users of that DNS server view the internet!
 - Assuming they use domain names





DNS Cache Poisoning

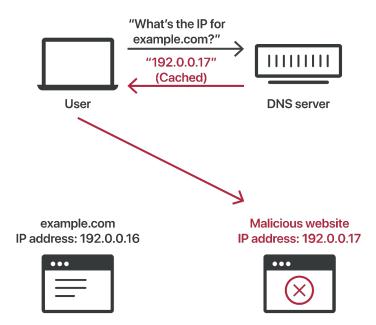
Attack: ???





DNS Cache Poisoning

- Attack: pre-empt DNS lookup by injecting malicious cache contents
 - Exploits DNS lookup optimization!
- Victim performs cache lookup, instead gets malicious domain IP
 - Attacker can redirect the victim's browser to the malicious website



Thwarting DNS Hijacking

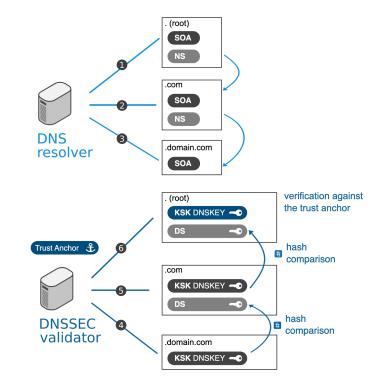
Defenses: ???



Thwarting DNS Hijacking

DNS-level authentication

- DNSSec
- Public-key crypto to "sign" DNS records
- Endpoint authentication
 - Certify that what I am seeing really is bank.com
 - Transport Layer Security (TLS)





Questions?





This time on CS 4440...

Network Availability Denial of Service (DoS) Attacks Transport, Link, Network, and Physical DoS



- Confidentiality: ???
- Authenticity: ???
- Integrity: ???
- Access Control: ???



- Confidentiality: Concealment of information or resources
 Attacks: ???
- Authenticity: Identification and assurance of info origin
 - Attacks: ???
- Integrity: Preventing improper and unauthorized changes
 - Attacks: ???
- Access Control: Enforce who is allowed access to what
 - Attacks: ???



- **Confidentiality:** Concealment of information or resources
 - Attacks: intercept credentials, info
- Authenticity: Identification and assurance of info origin
 - Attacks: SMTP header spoofing
- Integrity: Preventing improper and unauthorized changes
 - Attacks: tampering HTML over HTTP
- Access Control: Enforce who is allowed access to what
 - Attacks: web app code injection

- **Confidentiality:** Concealment of information or resources
 - Attacks: intercept credentials, info
- Authenticity: Identification and assurance of info origin
 - Attacks: SMTP header spoofing
- Integrity: Preventing improper and unauthorized changes
 - Attacks: tampering HTML over HTTP
- Access Control: Enforce who is allowed access to what
 - Attacks: web app code injection
- **Availability:** Ability to use desired information or resource
 - Attacks: ???

- **Confidentiality:** Concealment of information or resources
 - Attacks: intercept credentials, info
- Authenticity: Identification and assurance of info origin
 - Attacks: SMTP header spoofing
- Integrity: Preventing improper and unauthorized changes
 - Attacks: tampering HTML over HTTP
- Access Control: Enforce who is allowed access to what
 - Attacks: web app code injection
- **Availability:** Ability to use desired information or resource
 - Attacks: denial of service

Denial of Service Attacks





DoS: Denial of Service

• Goal: ???



DoS: Denial of Service

- **Goal:** make a service unusable, usually by overloading the server or network
- How?



DoS: Denial of Service

- **Goal:** make a service unusable, usually by overloading the server or network
- How?
 - Trigger the host to **crash**
 - Application-based DoS
 - Memory corruption
 - Consume host's resources
 - TCP SYN floods
 - ICMP ECHO (ping) floods
 - Consume host's **bandwidth**
 - UDP floods
 - ICMP floods

Amazon loses **\$66,000 per minute** of downtime

Higher security makes DoS attacks **more likely**

Common DoS Attacks

- Locally-induced crash
 - Exploit host's OS or server software bug
- Local resource consumption
 - fork() bomb, fill disk, deep directory nesting
- Deny service to individual hosts
 - Force crash or outage of critical services
- Remotely-induced crash
 - "Magic" packets—ping of death, teardrop
- Remote resource consumption
 - Syslog, SYN, fragment flood, UDP storm

1	
This webpage is not available	e
Reload	
	_



Common DoS Attacks (cont.)

- Deny service to an entire network
 - Target vulnerable links, critical network infrastructure
- Remotely-induced network outage
 - Attacks against routers, DNS servers
 - Redirected routes—forged routing information
- Remote network congestion
 - Remote control of compromised hosts ("zombies")
 - Allows for coordinated flooding
 - Distributed Denial of Service (DDoS)







Simple DoS Attacks

- Attacker spoofs their source address to hide origin
 - Defenses: • ???

 Attacker

 Victim
 Victim
 Victim

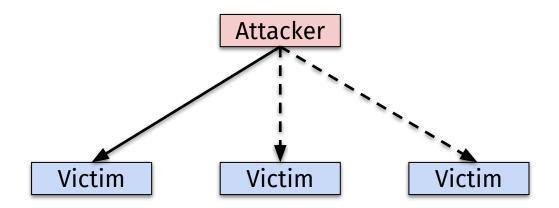


Simple DoS Attacks

Attacker spoofs their source address to hide origin

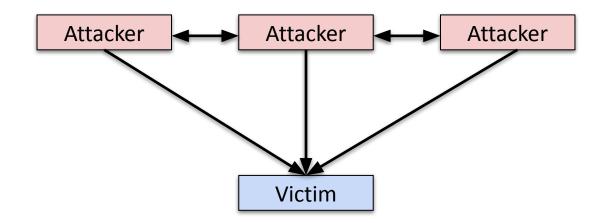
Defenses:

- Block source IP address
- Firewall
- ISP-level blocking
- Ignore requests

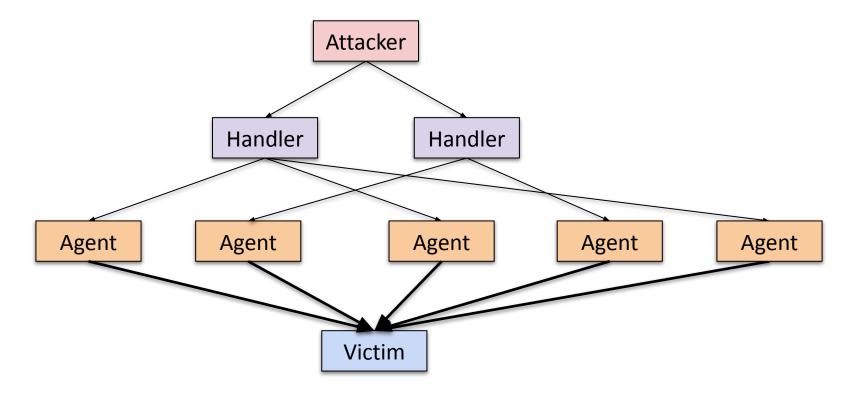


Coordinated DoS Attacks

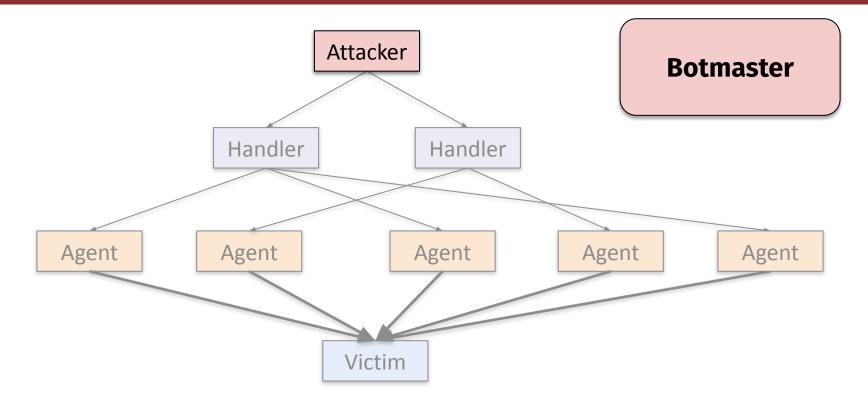
- Multiple willing attackers coordinate an attack on victim(s)
 - Same source-spoofing techniques as before
 - Harder to deal with





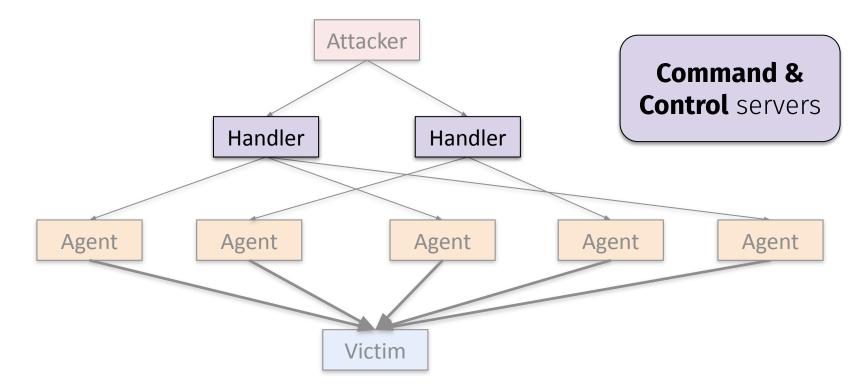


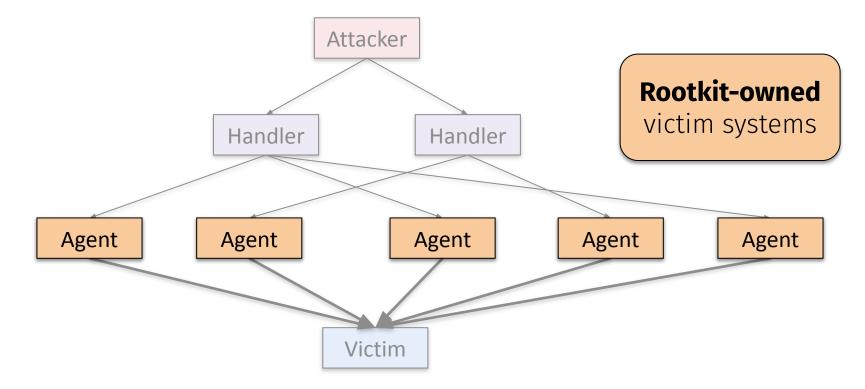






Stefan Nagy





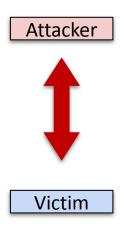
Timeline of a DDoS Attack

- **Goal:** compromise a large number of machines to form a botnet
 - 1. Attacker identifies **exploitable hosts** with scanners or other techniques
 - 2. Attacker gains control over systems via exploits, password cracking, etc.
 - 3. Attacker installs rootkit
 - 4. Attacker **remotely instructs** compromised machines to **attack the target**



Timeline of a DDoS Attack

- **Goal:** compromise a large number of machines to form a botnet
 - 1. Attacker identifies **exploitable hosts** with scanners or other techniques
 - 2. Attacker gains control over systems via exploits, password cracking, etc.
 - 3. Attacker installs rootkit
 - 4. Attacker remotely instructs compromised machines to attack the target



Timeline of a DDoS Attack

- **Goal:** compromise a large number of machines to form a botnet
 - 1. Attacker identifies **exploitable hosts** with scanners or other techniques
 - 2. Attacker gains control over systems via exploits, password cracking, etc.
 - 3. Attacker installs rootkit
 - 4. Attacker remotely instructs compromised machines to attack the target



Real DDoS Botnets

The Mirai botnet explained: How teen scammers and CCTV cameras almost brought down the internet

Mirai took advantage of insecure IoT devices in a simple but clever way. It scanned big blocks of the internet for open Telnet ports, then attempted to log in default passwords. In this way, it was able to amass a botnet army.

Storm: the largest botnet in the world?

Timely spam blasts help spread highly aggressive malware

A tiny botnet launched the largest DDoS attack on record

A small but powerful army of just 5,000 devices generated a record-breaking web attack.









https://www.joeskc.com/

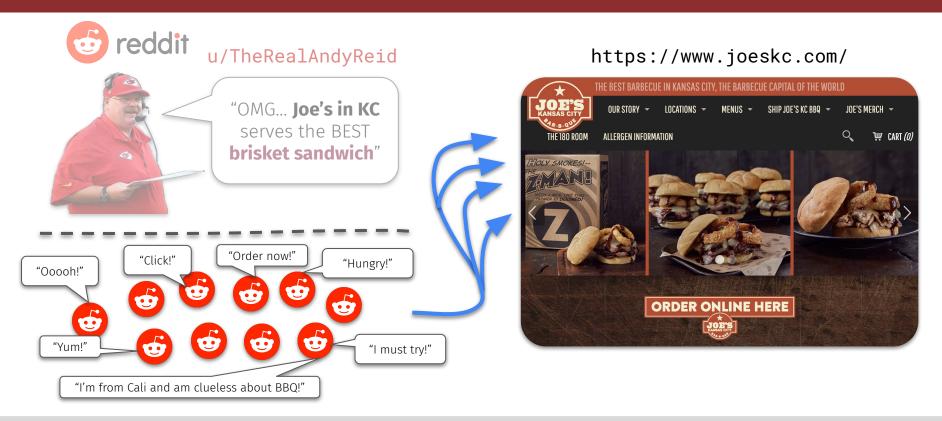


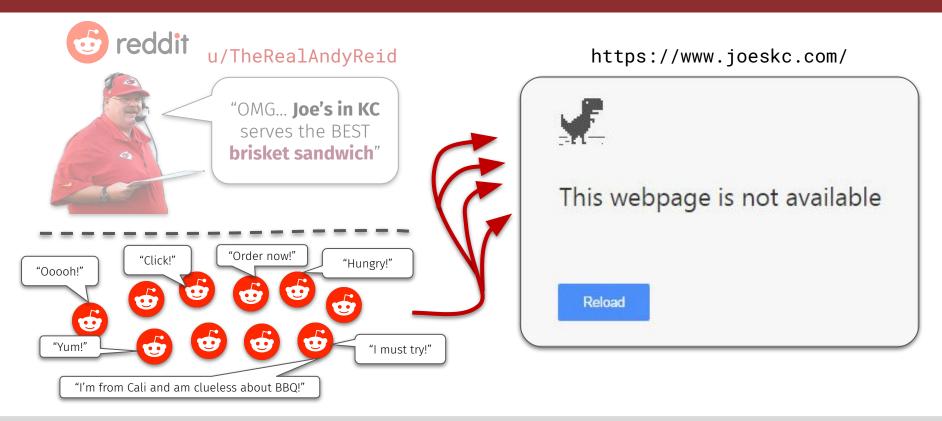




https://www.joeskc.com/







How can we differentiate between Flash Mob traffic and DDoS traffic?



How can we differentiate between Flash Mob traffic and DDoS traffic?

Flash Mob traffic

- Many clients using service legitimately
- "Slashdot Effect", "Reddit Hug of Death"
 - Traffic dies down when the network is flooded
- Sources in flash crowd are clustered
 - Usually by location (e.g., USA)

What Does Slashdot Effect Mean?

The slashdot effect refers to a temporary surge in traffic to a website, which can occur when a high-traffic website posts a link to smaller site or blog, thus directing an unprecedented surge in traffic. If the traffic increase is very large, it slow the site down or make it unreachable. The site is then considered to have been "slashdotted."

It's when someone posts a link to a website saying "Everyone, look at this website!" and everyone *does*. This puts so much traffic on the site in question's servers that they get overloaded and crash, causing the site to be inaccessible until the amount of traffic dies down a bit.



How can we differentiate between Flash Mob traffic and DDoS traffic?

Flash Mob traffic

- Many clients using service legitimately
- "Slashdot Effect", "Reddit Hug of Death"
 - Traffic dies down when the network is flooded
- Sources in flash crowd are clustered
 - Usually by location (e.g., USA)

DDoS traffic

- Attack does not end when host crashes
- Scattered locations (e.g., entire world)

What Does Slashdot Effect Mean?

The slashdot effect refers to a temporary surge in traffic to a website, which can occur when a high-traffic website posts a link to smaller site or blog, thus directing an unprecedented surge in traffic. If the traffic increase is very large, it slow the site down or make it unreachable. The site is then considered to have been "slashdotted."

It's when someone posts a link to a website saying "Everyone, look at this website!" and everyone *does*. This puts so much traffic on the site in question's servers that they get overloaded and crash, causing the site to be inaccessible until the amount of traffic dies down a bit.

Questions?





Transport Layer DoS

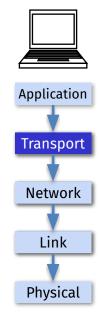


Recap: The Transport Layer

What does it facilitate?

???

- Key protocols?
 - Protocol 1: ???
 - Characteristics: ???
 - Protocol 2: ????
 - Characteristics: ???





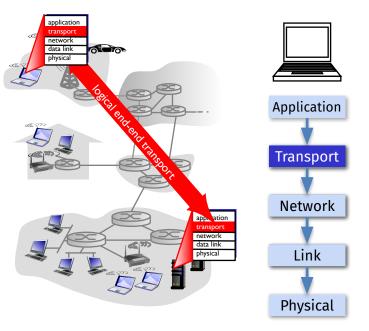
Recap: The Transport Layer

What does it facilitate?

Communication between apps on different hosts

Key protocols?

- Protocol 1: TCP (Transmission Control Protocol)
 - Characteristics: slow/complex but reliable
- Protocol 2: UDP (User Datagram Protocol)
 - Characteristics: fast/simple but unreliable





Recall: TCP is a connection-oriented protocol

Initiate with three-way "handshake": SYN, SYN-ACK, ACK



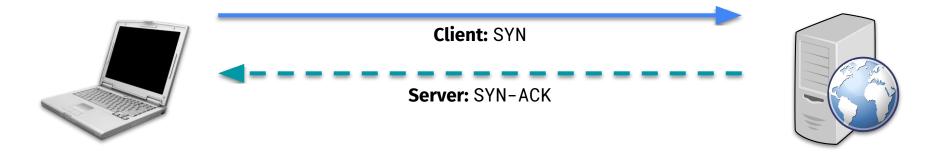
Client: SYN





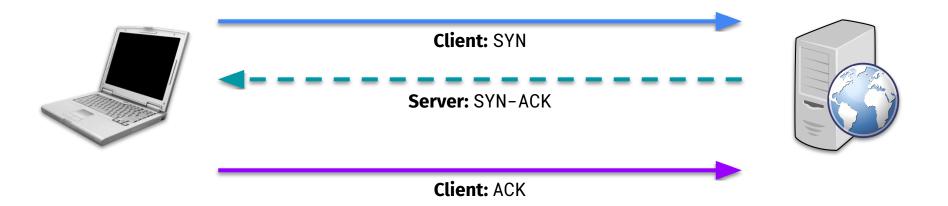
Recall: TCP is a connection-oriented protocol

Initiate with three-way "handshake": SYN, SYN-ACK, ACK



Recall: TCP is a connection-oriented protocol

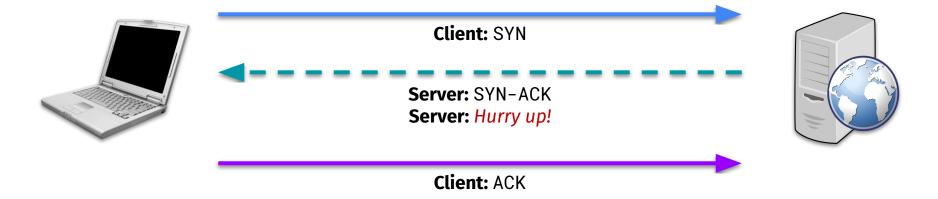
Initiate with three-way "handshake": SYN, SYN-ACK, ACK





Recall: TCP is a connection-oriented protocol

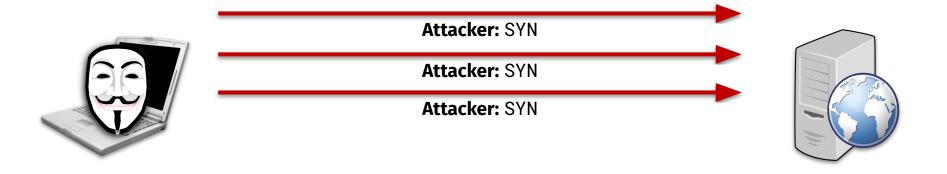
- Initiate with three-way "handshake": SYN, SYN-ACK, ACK
- Server waits until client responds with ACK





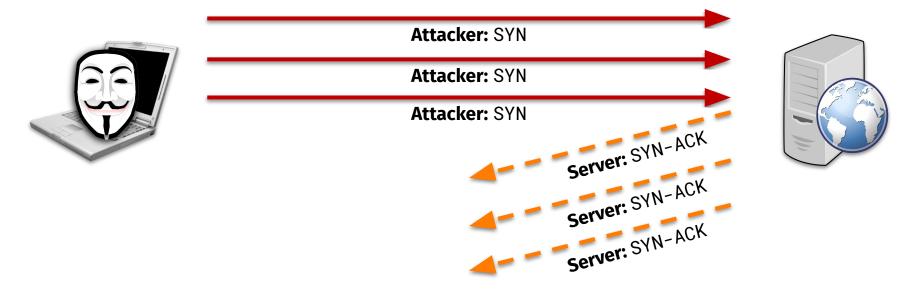
SYN Flooding Attack

Attack: spam SYN packets to server, with spoofed origin address



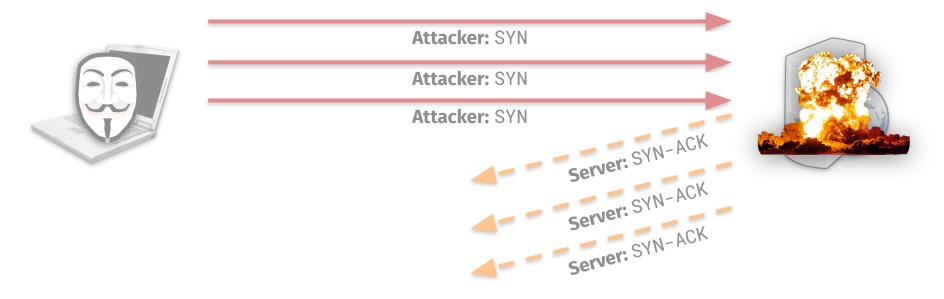
SYN Flooding Attack

Attack: spam SYN packets to server, with spoofed origin address



SYN Flooding Attack

- Attack: spam SYN packets to server, with spoofed origin address
 - Server's resources completely reserved—now can't serve legitimate clients





Thwarting SYN Flooding

Attack: spam SYN packets to server, with spoofed origin address

Server's resources completely reserved—now can't serve legitimate clients





Thwarting SYN Flooding

- Attack: spam SYN packets to server, with spoofed origin address
 - Server's resources completely reserved—now can't serve legitimate clients



How can we **prevent** SYN flooding?

Incorporate **state**—use **SYN cookies!**



Questions?





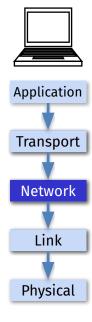
Network Layer DoS



Recap: The Network Layer

- What does it facilitate?
 - ???
- Key functions?
 - Function1: ???
 - Function1: ???
- Addressing?

???





Recap: The Network Layer

What does it facilitate?

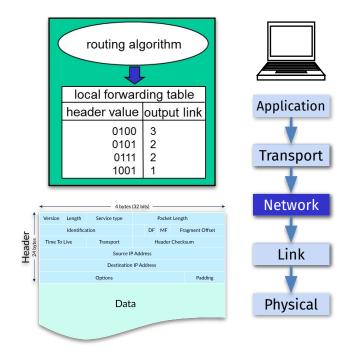
Sending of data from host on one network to another

Key functions?

- **Function1:** Routing: (find the shortest path for a packet)
- Function1: Forwarding (send packet on to the next hop)

Addressing?

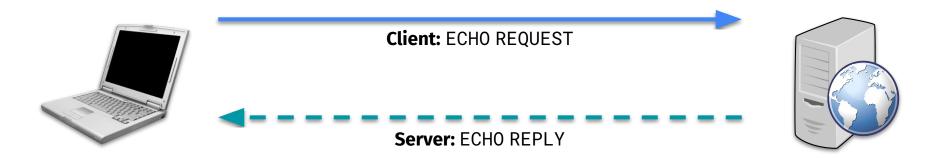
IP addressing (logical addressing)





ICMP: Internet Control Message Protocol

- ICMP: pings to determine whether a system is connected to the Internet
 - Analogous to "Hello, are you still there?"



ICMP Smurf Attacks

• Attack: takes advantage of broadcast-enabled hosts to amplify attack



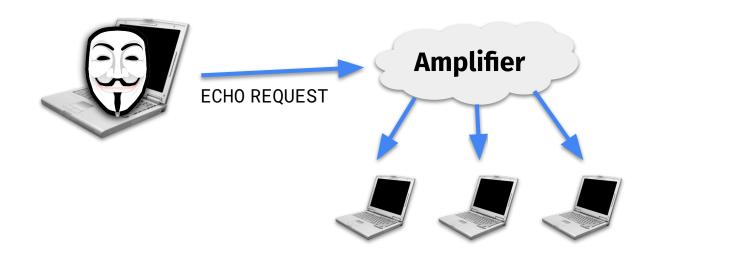






ICMP Smurf Attacks

• Attack: takes advantage of broadcast-enabled hosts to amplify attack

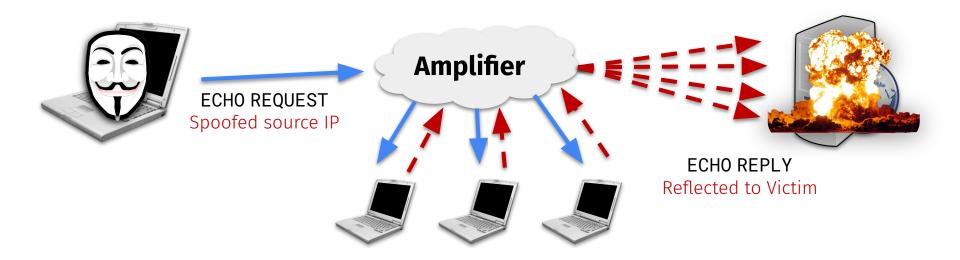






ICMP Smurf Attacks

- Attack: takes advantage of broadcast-enabled hosts to amplify attack
- Attacker spams spoofed-source ICMP requests, reflected to victim's IP





Advanced DoS Strategies

Reflection:

IP spoofing to redirect response to a victim

Amplification:

 Technique that increases the amount of traffic or packet size that the victim sees versus what the attacker originally sent

Common in real-world DDoS attacks

- Harder to detect (source obfuscation)
- Harder to thwart (changing sources)





ICMP Ping of Death Attack

Internet Protocol: IPV4 packets should be less than 65,536 bytes

Packets can be sent in **fragments** and **reassembled** by receiver

IP	ICMP	ICMP
Header	Header	Data
20 bytes	8 bytes	65,508 bytes



ICMP Ping of Death Attack

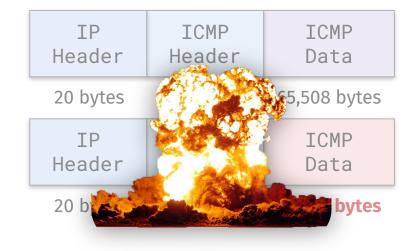
- Internet Protocol: IPV4 packets should be less than 65,536 bytes
 - Packets can be sent in **fragments** and **reassembled** by receiver
- Attack: send packet in fragments that reassemble to 64K+ bytes
 - Many historical computer systems could not handle larger packets

20 bytes	8 bytes	65,510 bytes
IP	ICMP	ICMP
Header	Header	Data
20 bytes	8 bytes	65,508 bytes
IP	ICMP	ICMP
Header	Header	Data



ICMP Ping of Death Attack

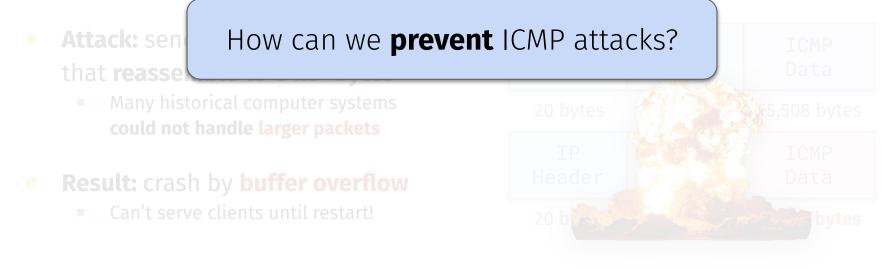
- Internet Protocol: IPV4 packets should be less than 65,536 bytes
 - Packets can be sent in **fragments** and **reassembled** by receiver
- Attack: send packet in fragments that reassemble to 64K+ bytes
 - Many historical computer systems could not handle larger packets
- Result: crash by buffer overflow
 - Can't serve clients until restart!





Thwarting ICMP-based DoS

- Internet Protocol: IPV4 packets should be less than 65,536 bytes
 - Packets can be sent in **fragments** and **reassembled** by receiver





Thwarting ICMP-based DoS

- Internet Protocol: IPV4 packets should be less than 65,536 bytes
 - Packets can be sent in fragments and reassembled by receiver





Questions?





Link Layer DoS





Recap: The Data Link Layer

What does it facilitate?

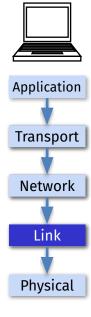
???

Addressing?

???

• Authenticity?

???





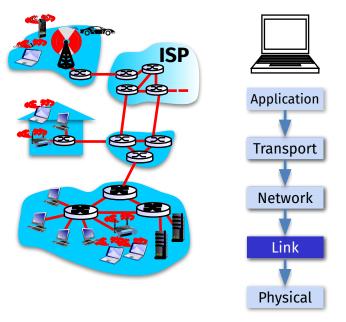
Recap: The Data Link Layer

What does it facilitate?

Responsible for the node-to-node delivery of data

Addressing?

- MAC addresses
 - Physical identifier for hardware
- Authenticity?
 - No—MAC addresses can be changed!

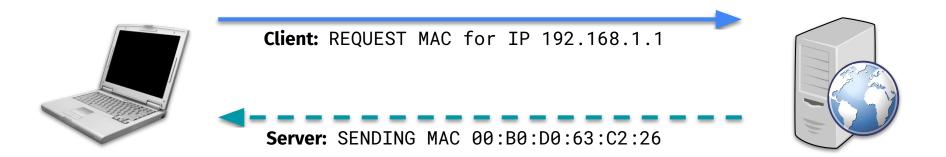




ARP: Address Resolution Protocol

• **ARP:** query to **resolve the MAC address** given a desired host IP

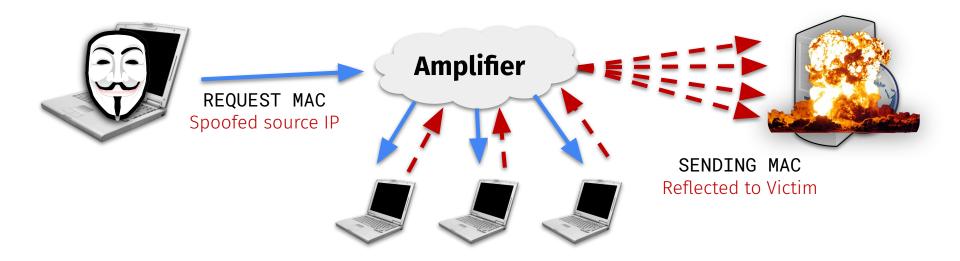
• How we know which **physical** address to transmit data to from its logical address



ARP Flooding Attack

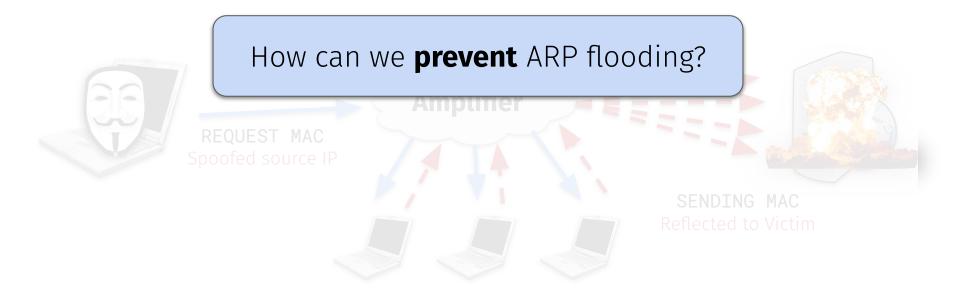
• Attack: same idea as ICMP Smurfing; spoof source to victim and spam away!

Victim gets overwhelmed by ARP replies and bandwith crashes



Thwarting ARP Flooding

- Attack: same idea as ICMP Smurfing; spoof source to victim and spam away!
 - Victim gets overwhelmed by ARP replies and bandwith crashes



Thwarting ARP Flooding

- Attack: same idea as SYN flood; spoof source to victim and spam away!
 - Victim gets overwhelmed by ARP replies and bandwith crashes



Physical Layer DoS

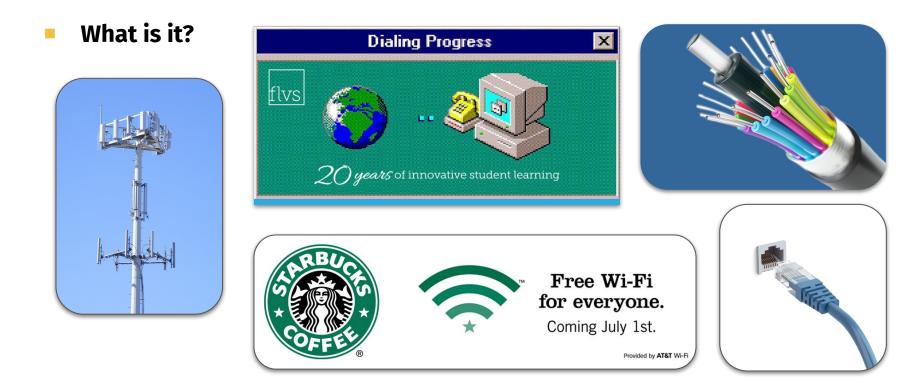


Recap: Physical Layer

What is it?



Recap: Physical Layer



Physical Layer DoS

Russian Spy Submarines Are Tampering with Undersea Cables That Make the Internet Work. Should We Be Worried?

A massive cable attack is probably an over-hyped scenario, at least for a country with as many redundant cables as the United States pitted against a limited number of Russian special-operations submarines.



CNN Exclusive: FBI investigation determined Chinese-made Huawei equipment could disrupt US nuclear arsenal communications

Physical Layer DoS

Russian Spy Submarines Are Tampering with Undersea Cable<u>s That Make the Internet</u>

Work A massive cab A many reduit A mini protests grow

Social media platforms have also been cut off in areas of Tehran and Kurdistan as videos of dissent go viral

Chinese-made Huawei equipment could disrupt US nuclear arsenal communications

Thwarting Physical Layer Attacks





Questions?





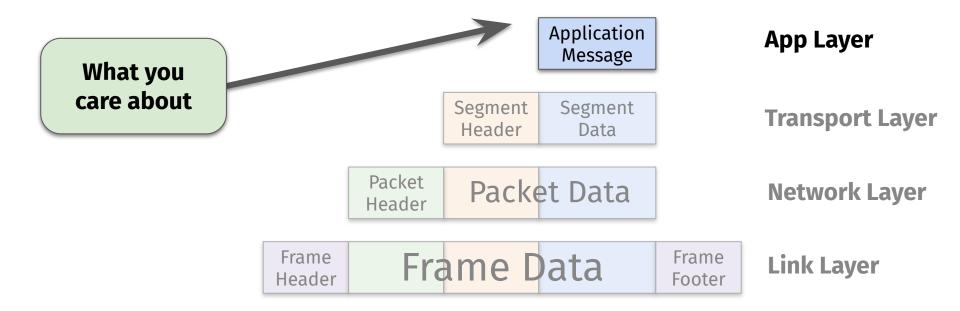
Analyzing Network Packets





Recap: Internet Packet Encapsulation

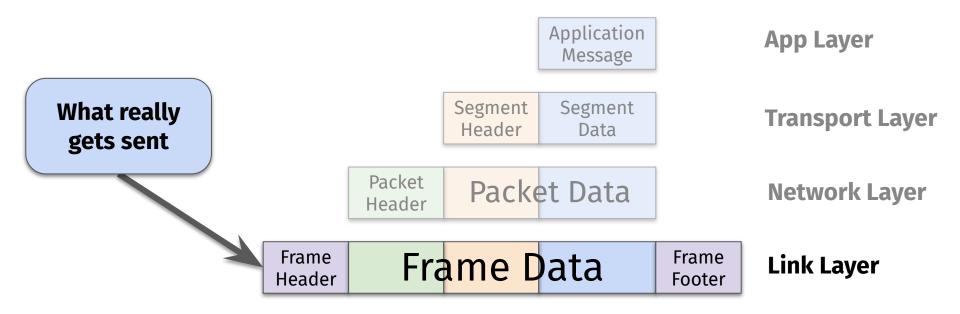
How packets are generated and sent





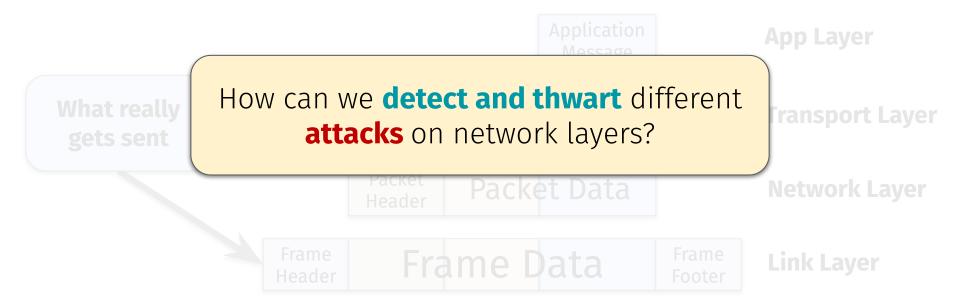
Recap: Internet Packet Encapsulation

How packets are generated and sent



Recap: Internet Packet Encapsulation

How packets are generated and sent



Tools of the Trade

Packet Analyzers:

- Tools for dissecting network packets
- Packet Analyzers allow you to:
 - Identify unusual packets
 - Characterize network activity
 - Pinpoint malicious traffic
- The basis of modern-day network security (e.g., firewalls, antivirus)



Familiarity with packet analysis tools?

I eat NetSec CTF challenges like a kid eats candy on Halloween. 🧙

Some (e.g., Wireshark, DPKT, Scapy, or something else)

None (but that's totally okay!)





0%

0%

0%



Tools of the Trade: Wireshark

- A "graphical interface" for manual packet analysis
 - Completely **open-source and free**
- General workflow:
 - Load up a PCAP (packet capture)
 - Wireshark will display each packet
 - Inspect particular fields of interest

File	Edit	View		- (Canata		Anala		testistis	. Т.I	a a la a u	14	lizalas	s Too		elp		
-ile	Eait	-	-	-												eip		
	E	۲		010			(@		2 0	<u> </u>		6		Q.				
syn	nergy.	packet	type	== 1	DKUP											\times		- +
о.		Time			So	urce				Destir	nation			Pr	otocol	Leng	th Ir	
	428	3.994	768		19	2.16	8.1.1	.06		192.	168.1	1.213		S	nergy	,	68	
	442	4.40	2758		19	2.16	8.1.1	.06		192.	168.1	L.213		5)	nergy		68	
	453	5.04	2758		19	2.16	8.1.1	.06		192.	168.1	L.213		5)	nergy		68	
	464	5.29	740		19	2.16	8.1.1	.06		192.	168.1	L.213		5)	nergy	r (68	
	486	5.820	5760		19	2.16	8.1.1	.06		192.	168.1	L.213		sy	nergy	r I	68	
	494	5.978	3736		19	2.16	8.1.1	.06		192.	168.1	1.213		5)	nergy	r 1	68	
	512	6.18	5737		19	2.16	8.1.1	.06		192.	168.1	L.213		5)	nergy	r (68	
	E10	6.314	737		19	2.16	8.1.1	.06		192.	168.1	1.213		5)	nergy	0	68	
	213																	
Fr	532 ame hern	et II	68 b , Sr	c: /	s on ASUS	wir TekC	_14:f	4 bi 6:e8	(38:	68 by d5:47	:14:f	aptu 6:e8	red (), Ds	544 b t: Ap	its) ple_9 213	on in		ace
> Et > In > Tr	532 ame hern tern ansm	428: et II et Pr issic	68 b , Sr otoc n Co	c: / ol \ ntro	s on ASUS Vers	wir TekC	e (54 _14:f 4, Sr	4 bit 6:e8 c: 19	(38: 92.16	68 by d5:47 8.1.1	tes c :14:f 06, D	aptu 6:e8 0st:	red (), Ds 192.1	544 b t: Ap 68.1.	its) ple_9	on in d:dc	> nterf :83 (ace a0:7
Fr Et In Tr	ame hern tern ansm	428: et II et Pr	68 b , Sr otoc n Co toco	c: / ol \ ntro 1	s on ASUS Vers pl P	wir TekC	e (54 _14:f 4, Sr	4 bit 6:e8 c: 19	(38: 92.16	68 by d5:47 8.1.1	tes c :14:f 06, D	aptu 6:e8 0st:	red (), Ds 192.1	544 b t: Ap 68.1.	its) ple_9 213	on in d:dc	> nterf :83 (ace a0:7
Fr Et In Tr	532 ame hern tern ansm nerg Pacl	428: et II et Pr issic y Pro ket L	68 b , Sr otoc n Co toco engt	c: / ol \ ntro l h: 1	s on ASUS Vers pl P	wir TekC ion	e (54 _14:f 4, Sr	4 bit 6:e8 c: 19 Src P	(38: 92.16	68 by d5:47 8.1.1	tes c :14:f 06, D	aptu 6:e8 0st:	red (), Ds 192.1	544 b t: Ap 68.1.	its) ple_9 213	on in d:dc	> nterf :83 (ace a0:7
Fr Et In Tr Sy	ame hern tern ansm nerg Pacl	428: et II et Pr issic y Pro ket L	68 b , Sr otoc n Co toco engt ype:	c: / ol \ ntro l h: 1 Key	s on ASUS Vers pl P	wir TekC ion	≥ (54 _14:f 4, Sr col,	4 bit 6:e8 c: 19 Src P	(38: 92.16	68 by d5:47 8.1.1	tes c :14:f 06, D	aptu 6:e8 0st:	red (), Ds 192.1	544 b t: Ap 68.1.	its) ple_9 213	on in d:dc	> nterf :83 (ace a0:7
> Fr > Et > In > Tr > Sy	ame hern tern ansm nerg Pacl Rey	428: et II issic y Pro ket L ket T	68 b , Sr otoc n Co toco engt ype: ased	c: / ol \ ntro l h: 1 Key	s on ASUS Vers pl P	wir TekC ion	≥ (54 _14:f 4, Sr col,	4 bit 6:e8 c: 19 Src P	(38: 92.16	68 by d5:47 8.1.1	tes c :14:f 06, D	aptu 6:e8 0st:	red (), Ds 192.1	544 b t: Ap 68.1.	its) ple_9 213	on in d:dc	> nterf :83 (ace a0:7
Fr Et In Tr Sy	532 ame hern tern ansm nerg Pacl Pacl Key	428: et II et Pr issic y Pro ket L ket T Rele	68 b , Sr otoc n Co toco engt ype: ased d: 1	c: / ol N ntro l h: 1 Key 16	s on ASUS Vers ol P LØ / Rei	wir TekC ion roto	≥ (54 _14:f 4, Sr col, ≥d (D	4 bit 6:e8 c: 19 Src P	(38: 92.16	68 by d5:47 8.1.1	tes c :14:f 06, D	aptu 6:e8 0st:	red (), Ds 192.1	544 b t: Ap 68.1.	its) ple_9 213	on in d:dc	> nterf :83 (ace a0:7
Fr Et In Tr Sy	532 ame hern tern ansm nerg Pacl Rey H	428: et II et Pr issic y Pro ket L ket T Rele Key I	68 b , Sr otoc n Co toco engt ype: ased d: 1 odif	c: / ol / ntro l h: 1 Key 16 ier	s on ASUS Vers ol P L0 / Rei	wir TekC ion roto	≥ (54 _14:f 4, Sr col, ≥d (D	4 bit 6:e8 c: 19 Src P	(38: 92.16	68 by d5:47 8.1.1	tes c :14:f 06, D	aptu 6:e8 0st:	red (), Ds 192.1	544 b t: Ap 68.1.	its) ple_9 213	on in d:dc	> nterf :83 (ace a0:7
Fr Et In Tr Sy	532 ame hern tern ansm nerg Pacl Rey H	428: et II et Pr issic y Pro ket L ket T Rele Key I Key M	68 b , Sr otoc n Co toco engt ype: ased d: 1 odif	c: / ol / ntro l h: 1 Key 16 ier	s on ASUS Vers ol P L0 / Rei	wir TekC ion roto	≥ (54 _14:f 4, Sr col, ≥d (D	4 bit 6:e8 c: 19 Src P	(38: 92.16	68 by d5:47 8.1.1	tes c :14:f 06, D	aptu 6:e8 0st:	red (), Ds 192.1	544 b t: Ap 68.1.	its) ple_9 213	on in d:dc	> nterf :83 (ace a0:7
> Fr > Et > In > Tr > Sy	532 ame hern tern ansm Pacl Pacl Key H	428: et II et Pr issic y Prc ket L ket T Rele Key I Key B	68 b , Sr otoc n Co toco engt ype: ased d: 1: odif	c: / ol \ ntro l h: 1 Key 16 ier n: 2	s on ASUS Vers ol P 10 / Re: 20	wire TekC ion a roto lease	≥ (54 _14:f 4, Sr col, ≥d (D	4 bit 6:e8 c: 19 Src P	(38: 92.16 Port:	68 by d5:47 8.1.1 2480	tes c :14:f 06, D 0, Ds	aptu 6:e8 0st: t Po	red (), Ds 192.1 rt: 4	544 b t: Ap 68.1. 9727,	its) ple_9 213 Seq:	on in d:dc	> nterf :83 (ace (a0:7)
<pre>> Fr > Et > In > Tr > Sy </pre>	ame hern tern nerg Pacl Key H H H	428: et II et Pr issic y Pro ket L ket T Rele Key I Key M	68 b , Sr otocc engt ype: ased d: 1 odif utto	c: / ol \ ntro l h: 1 Key 16 ier n: 2 dc	s on ASUS Vers ol P 10 7 Re: Mask 20 83	wirr TekC roto lease c: 81	e (54 _14:f 4, Sr col, ed (D) .92	4 bit 6:e8 c: 19 Src I KUP)	(38: 92.16 Port:	68 by d5:47 8.1.1 2480	tes c :14:f 06, D 0, Ds	aptu 6:e8 9st: t Po 00	red (), Ds 192.1 rt: 4	544 b t: Ap 68.1. 9727,	its) ple_9 213	on in d:dc 109:	> nterf :83 (ace (a0:7)
Fr Et In Tr Sy	ame hern tern nerg Pacl Key H H H H H O 00	428: et II et Pr issic y Prc ket L ket T Rele Key I Key B Key B	68 b , Sr otoco engt ype: ased d: 11 odif utto 7 9d 0 da	c: / ol \ ntro l h: 1 Key 16 ier n: 2 dc 40	s on ASUS Vers ol P 10 7 Re: 20 Mask 20 83 83 83	wirr TekC_ ion - roto lease c: 81	e (54 _14:f 4, Sr col, ed (D .92	4 bit 6:e8 c: 19 Src F KUP)	(38: 92.16 Port:	08 00 01 68	tes c :14:f 06, D 0, Ds 0, Ds 0 45 a c0	aptur 6:e8 0st: t Por 00 a8	red (), Ds 192.1 rt: 4	544 b t: Ap 68.1. 9727,	its) ple_9 213 Seq:	on in d:dc 109: 	> nterf :83 (a0:7

Tools of the Trade: Wireshark

	Time	Source	Destination	Protocol	Length Info
	1 0.000000	10.0.0.2	10.128.0.2	TCP	54 3341 → 80 [SYN] Seq=0 Win=512 Len=0
	2 0.003987	10.128.0.2	10.0.0.2	TCP	58 80 → 3222 [SYN, ACK] Seq=0 Ack=1 Win=29200 Len=0 MSS=1460
	3 0.005514	10.128.0.2	10.0.0.2	TCP	58 80 → 3341 [SYN, ACK] Seq=0 Ack=1 Win=29200 Len=0 MSS=1460
	4 0.008429	10.0.0.2	10.128.0.2	TCP	54 3342 → 80 [SYN] Seq=0 Win=512 Len=0
	5 0.010233	10.128.0.2	10.0.0.2	TCP	58 80 → 3220 [SYN, ACK] Seq=0 Ack=1 Win=29200 Len=0 MSS=1460
	6 0.014072	10.128.0.2	10.0.0.2	TCP	58 80 → 3342 [SYN, ACK] Seq=0 Ack=1 Win=29200 Len=0 MSS=1460
	7 0.016830	10.0.0.2	10.128.0.2	TCP	54 3343 → 80 [SYN] Seq=0 Win=512 Len=0
	8 0.022220	10.128.0.2	10.0.0.2	TCP	58 80 → 3343 [SYN, ACK] Seq=0 Ack=1 Win=29200 Len=0 MSS=1460
	9 0.023496	10.128.0.2	10.0.0.2	TCP	58 80 → 3219 [SYN, ACK] Seq=0 Ack=1 Win=29200 Len=0 MSS=1460
	10 0.025243	10.0.0.2	10.128.0.2	TCP	54 3344 → 80 [SYN] Seq=0 Win=512 Len=0
	11 0.026672	10.128.0.2	10.0.0.2	TCP	58 80 → 3218 [SYN, ACK] Seq=0 Ack=1 Win=29200 Len=0 MSS=1460
	12 0.028038	10.128.0.2	10.0.0.2	TCP	58 80 → 3221 [SYN, ACK] Seq=0 Ack=1 Win=29200 Len=0 MSS=1460
	13 0.030523	10.128.0.2	10.0.0.2	TCP	58 80 → 3344 [SYN, ACK] Seq=0 Ack=1 Win=29200 Len=0 MSS=1460
[estination Por Stream index: : TCP Segment Le	1]			
	Stream index: : TCP Segment Le equence number Next sequence I	1] n: 0] : 0 (relative se number: 0 (relat	equence number) tive sequence number)] tive ack number)	(
	Stream index: : TCP Segment Le equence number Next sequence cknowledgment	1] n: 0] : 0 (relative se number: 0 (relat number: 1 (relat	ive sequence number)] ive ack number)	I	
	Stream index: : TCP Segment Le equence number Next sequence cknowledgment	1] n: 0] : 0 (relative se number: 0 (relat number: 1 (relat <u>der Length</u> : 24 byte	ive sequence number)] ive ack number)		
D([[S(A(0))	Stream index: : TCP Segment Le equence number Next sequence n cknowledgment n 110 = Hea	1] n: 0] : 0 (relative se number: 0 (relat number: 1 (relat der Length: 24 byte YN, ACK)	ive sequence number)] ive ack number)		
D([: S(0: W: W:	Stream index: : TCP Segment Le equence number Next sequence i cknowledgment i 110 = Hea lags: 0x012 (S indow size val Calculated win	1] n: 0] : 0 (relative se number: 0 (relat number: 1 (relat der Length: 24 byte YN, ACK) ue: 29200 dow size: 29200]	ive sequence number)] ive ack number)		
D([[A(0) W: (C	Stream index: : TCP Segment Let equence number Next sequence cknowledgment 110 = Heat lags: 0x012 (S' indow size val Calculated win hecksum: 0x4266	1] n: 0] rowspace of the set of	ive sequence number)] ive ack number)		
D([' S(S) (1) A(0) (1) (1) (1) (1) (1) (1) (1) (1) (1) (1	Stream index: : TCP Segment Le equence number Next sequence : cknowledgment i 110 = Hea lags: 0x012 (S indow size val Calculated win hecksum: 0x426 Checksum Statu	1] n: 0] : 0 (relative se number: 0 (relat number: 1 (relat der Length: 24 byte YN, ACK) ue: 29200 dow size: 29200] 8 [unverified] s: Unverified]	ive sequence number)] ive ack number)		
D([[S(0) 0) W: V: CI CI CI	Stream index: : TCP Segment Le equence number Next sequence cknowledgment 110 = Heaa lags: 0x012 (S indow size val Calculated win hecksum: 0x426 Checksum Statu: rgent pointer:	1] n: 0] : 0 (relative se number: 0 (relat number: 1 (relat der Length: 24 byte YN, ACK) ue: 29200 dow size: 29200] 8 [unverified] s: Unverified] 0	ive sequence number)] ive ack number) es (6)		
D([: [] A(0) () W: () CI [] UI () UI ()	Stream index: TCP Segment Lei equence number Next sequence i cknowledgment i 110 = Hear lags: 0x012 (S indow size val Calculated win hecksum: 0x426 Checksum Statu rgent pointer: ptions: (4 byt	1] n: 0] : 0 (relative se number: 0 (relat number: 1 (relat der Length: 24 byte YN, ACK) ue: 29200 dow size: 29200] 8 [unverified] s: Unverified]	ive sequence number)] ive ack number) es (6)		
D([[S(0) (0) (V) (Cl (Cl (U) (V) (V) (V) (V) (V) (V) (V)	Stream index: : TCP Segment Le equence number Next sequence cknowledgment 110 = Heaa lags: 0x012 (S indow size val Calculated win hecksum: 0x426 Checksum Statu: rgent pointer:	1] n: 0] : 0 (relative se number: 0 (relat number: 1 (relat der Length: 24 byte YN, ACK) ue: 29200 dow size: 29200] 8 [unverified] s: Unverified] 0	ive sequence number)] ive ack number) es (6)		
D([: S(0) () () () () () () () () () () () () ()	Stream index: TCP Segment Lei equence number Next sequence i cknowledgment i 110 = Hear lags: 0x012 (S indow size val Calculated win hecksum: 0x426 Checksum Statu rgent pointer: ptions: (4 byt	1] n: 0] : 0 (relative se number: 0 (relat number: 1 (relat der Length: 24 byte YN, ACK) ue: 29200 dow size: 29200] 8 [unverified] s: Unverified] 0	ive sequence number)] ive ack number) es (6)		
D [: S [] A 0: W: C C C C U U V V V	Stream index: TCP Segment Lei equence number Next sequence i cknowledgment i 110 = Hear lags: 0x012 (S indow size val Calculated win hecksum: 0x426 Checksum Statu rgent pointer: ptions: (4 byt	1] n: 0] : 0 (relative se number: 0 (relat number: 1 (relat der Length: 24 byte YN, ACK) ue: 29200 dow size: 29200] 8 [unverified] s: Unverified] 0	ive sequence number)] ive ack number) es (6)		

Tools of the Trade: Scapy

- Python API for programmatic packet capture and analysis
 - Think of it as "Wireshark in API form"
 - Project 4: you will use Scapy to write your own packet analysis scripts

Scapy API reference		
scapy.asn1 package	希 / Scapy API reference	O Edit on GitHut
scapy.contrib package		
scapy.layers package	Scony ADI reference	
scapy.ansmachine	Scapy API reference	
scapy.as_resolvers	Scapy: create, send, sniff, dissect and manipulate network	packets.
scapy.asn1fields		
scapy.asn1packet	Usable either from an interactive console or as a Python li	ibrary. https://scapy.net
scapy.automaton		
scapy.autorun	Subpackages	
scapy.base_classes	 scapy.asn1 package 	
scapy.config	scapy.contrib package	
scapy.consts	scapy.layers package	
scapy.dadict		
scapy.data	Submodules	
scapy.error		
scapy.fields	 scapy.ansmachine scapy.as resolvers 	
scapy.interfaces	 scapy.as_resolvers scapy.asn1fields 	
scapy.main	scapy.asn1packet	
scapy.packet	scapy.automaton	
scapy.pipetool	• scapy.autorun	
	 scapy.base_classes 	

Tools of the Trade: Scapy

- Python API for programmatic packet capture and analysis
 - Think of it as "Wireshark in API form"
 - Project 4: you will use Scapy to write your own packet analysis scripts
- We'll provide the PCAP traces...
 - You'll write code to analyze them!
 - Examples:
 - Detecting attacks on a network
 - Finding user credentials
 - Sniffing a user's browsing history

Scapy API reference		
scapy.asn1 package	I Scapy API reference	O Edit on GitHub
scapy.contrib package		
scapy.layers package	Scapy API reference	
scapy.ansmachine	Scapy API reference	
scapy.as_resolvers	Scapy: create, send, sniff, dissect and manipulate network packets.	
scapy.asn1fields		
scapy.asn1packet	Usable either from an interactive console or as a Python library. https://scap	y.net
scapy.automaton		
scapy.autorun	Subpackages	
scapy.base_classes	 scapy.asn1 package 	
scapy.config	scapy.asi1 package scapy.contrib package	
scapy.consts	scapy.layers package	
scapy.dadict		
scapy.data	Submodules	
scapy.error		
scapy.fields	 scapy.ansmachine scapy.as resolvers 	
scapy.interfaces	 scapy.as_resolvers scapy.asn1fields 	
scapy.main	• scapy.asn1packet	
scapy.packet	scapy.automaton	
scapy.pipetool	scapy.autorun	
scapy.pipetooi	 scapy.base_classes 	

Questions?





Next time on CS 4440...

Passwords and Secure Authentication

