Lecture 6: Assembly Programs

- Today's topics:

  - Procedures
  - Examples

Procedures

- Local variables, AR, $fp, $sp
- Scratchpad and saves/restores
- Arguments and returns
- jal and $ra

2

Procedures

- Each procedure (function, subroutine) maintains a scratchpad of register values – when another procedure is called (the callee), the new procedure takes over the scratchpad – values may have to be saved so we can safely return to the caller

  - parameters (arguments) are placed where the callee can see them
  - control is transferred to the callee
  - acquire storage resources for callee
  - execute the procedure
  - place result value where caller can access it
  - return control to caller

3

Jump-and-Link

- A special register (storage not part of the register file) maintains the address of the instruction currently being executed – this is the *program counter* (PC)

- The procedure call is executed by invoking the jump-and-link (jal) instruction – the current PC (actually, PC+4) is saved in the register $ra and we jump to the procedure's address (the PC is accordingly set to this address)
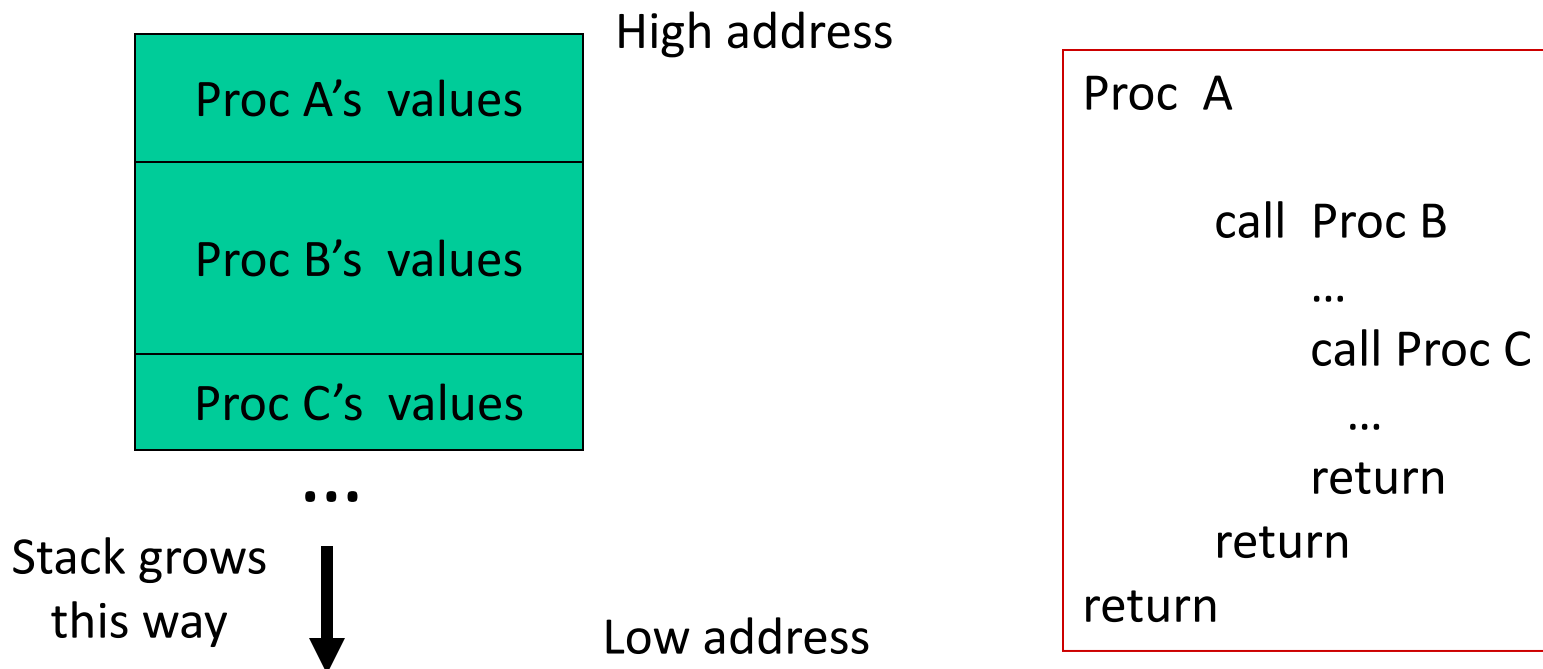
    jal   NewProcedureAddress

- Since jal may over-write a relevant value in $ra, it must be saved somewhere (in memory?) before invoking the jal instruction

- How do we return control back to the caller after completing the callee procedure?

The Stack

The register scratchpad for a procedure seems volatile –
it seems to disappear every time we switch procedures –
a procedure's values are therefore backed up in memory
on a stack

High address

| Proc A's values |
| Proc B's values |
| Proc C's values |

…

Stack grows
this way

Low address

Proc A

      call Proc B
        …
        call Proc C
         …
        return
      return
return

5

Saves and Restores

Storage Management on a Call/Return

- A new procedure must create space for all its variables on the stack

- Before/after executing the jal, the caller/callee must save relevant values in $s0-$s7, $a0-$a3, $ra, $fp, temps into the stack space

- Arguments are copied into $a0-$a3; the jal is executed

- After the callee creates stack space, it updates the value of $sp

- Once the callee finishes, it copies the return value into $v0, frees up stack space, and $sp is incremented

- On return, the caller/callee brings in stack values, ra, temps into registers

- The responsibility for copies between stack and registers may fall upon either the caller or the callee

Example 1 (pg. 98)

int leaf_example (int g, int h, int i, int j)
{
    int f ;
    f = (g + h) − (i + j);
    return f;
}

Notes:
In this example, the callee took care of saving the registers it needs.

The caller took care of saving its $ra and $a0-$a3.

Could have avoided using the stack altogether.

```
leaf_example:
  addi    $sp, $sp, -12
  sw      $t1, 8($sp)
  sw      $t0, 4($sp)
  sw      $s0, 0($sp)
  add     $t0, $a0, $a1
  add     $t1, $a2, $a3
  sub     $s0, $t0, $t1
  add     $v0, $s0, $zero
  lw      $s0, 0($sp)
  lw      $t0, 4($sp)
  lw      $t1, 8($sp)
  addi    $sp, $sp, 12
  jr      $ra
```

Saving Conventions

- Caller saved: Temp registers $t0-$t9 (the callee won't bother saving these, so save them if you care), $ra (it's about to get over-written), $a0-$a3 (so you can put in new arguments), $fp (if being used by the caller)

- Callee saved: $s0-$s7 (these typically contain "valuable" data)

- Read the Notes on the class webpage on this topic

Example 2 (pg. 101)

```
int   fact  (int n)
{
    if (n < 1)  return (1);
        else return (n * fact(n-1));
}
```

Notes:
The caller saves $a0 and $ra
in its stack space.
Temp register $t0 is never saved.

```
fact:
   slti      $t0, $a0, 1
   beq      $t0, $zero, L1
   addi   $v0, $zero, 1
   jr       $ra
L1:
   addi    $sp, $sp, -8
   sw       $ra, 4($sp)
   sw       $a0, 0($sp)
   addi    $a0, $a0, -1
   jal      fact
   lw       $a0, 0($sp)
   lw       $ra, 4($sp)
   addi    $sp, $sp, 8
   mul     $v0, $a0, $v0
   jr       $ra
```

Dealing with Characters

- Instructions are also provided to deal with byte-sized and half-word quantities: lb (load-byte), sb, lh, sh

- These data types are most useful when dealing with characters, pixel values, etc.

- C employs ASCII formats to represent characters – each character is represented with 8 bits and a string ends in the null character (corresponding to the 8-bit number 0); A is 65, a is 97

Example 3 (pg. 108)

Convert to assembly:
void strcpy (char x[], char y[])
{
   int i;
   i=0;
   while ((x[i] = y[i]) != `\0')
     i += 1;
}

Notes:
Temp registers not saved.

```
strcpy:
  addi    $sp, $sp, -4
  sw      $s0, 0($sp)
  add     $s0, $zero, $zero
L1: add  $t1, $s0, $a1
  lb      $t2, 0($t1)
  add     $t3, $s0, $a0
  sb      $t2, 0($t3)
  beq     $t2, $zero, L2
  addi    $s0, $s0, 1
  j       L1
L2: lw    $s0, 0($sp)
  addi    $sp, $sp, 4
  jr      $ra
```