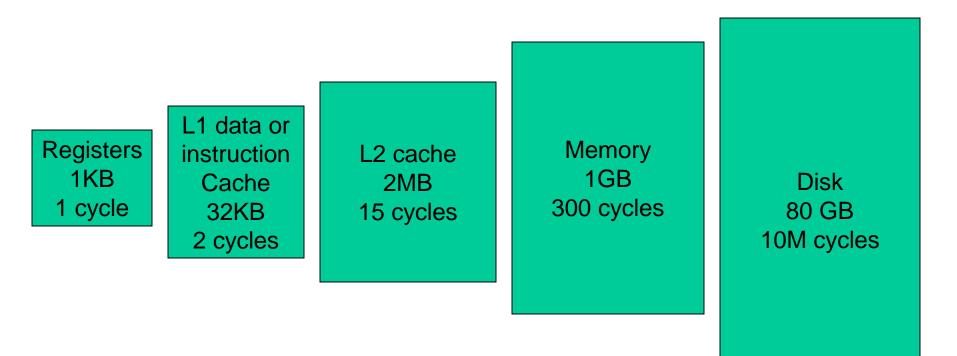
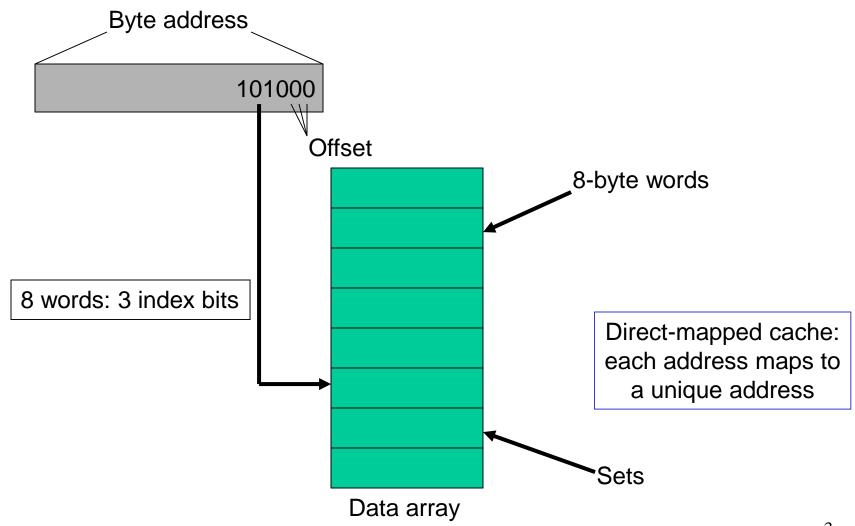
Lecture 21: Memory Hierarchy

- Today's topics:
 - Cache organization
 - Cache hits/misses

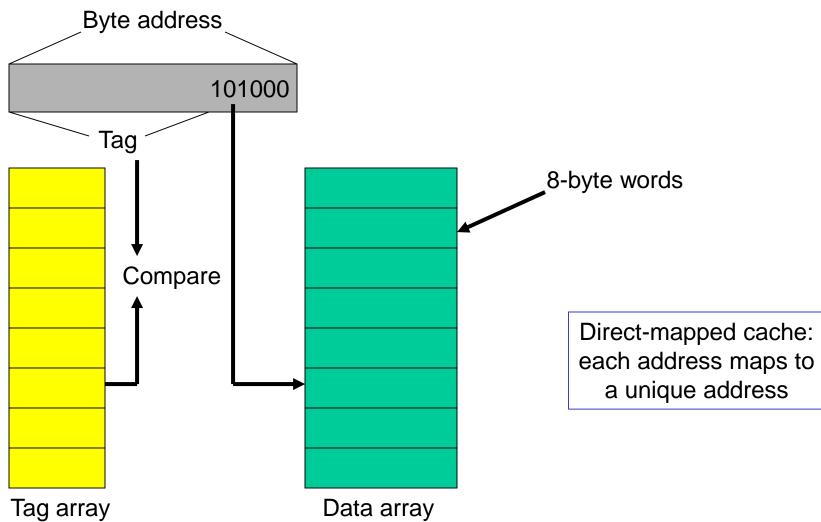
As you go further, capacity and latency increase



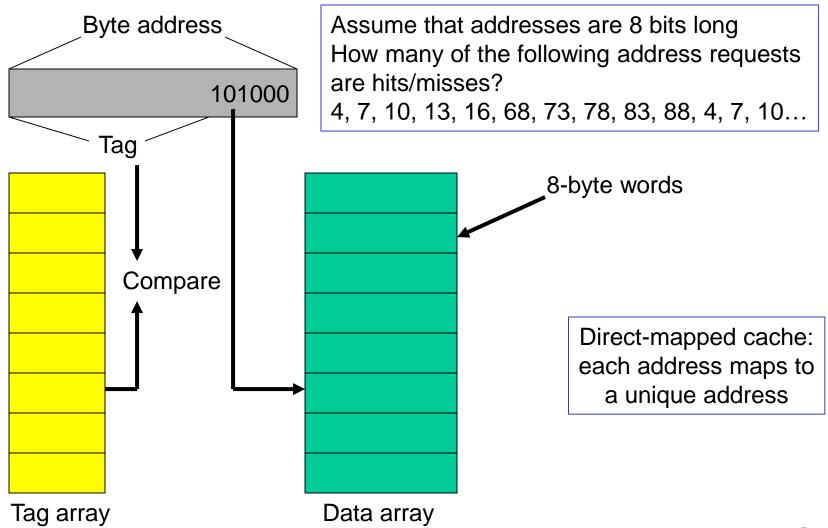
Accessing the Cache



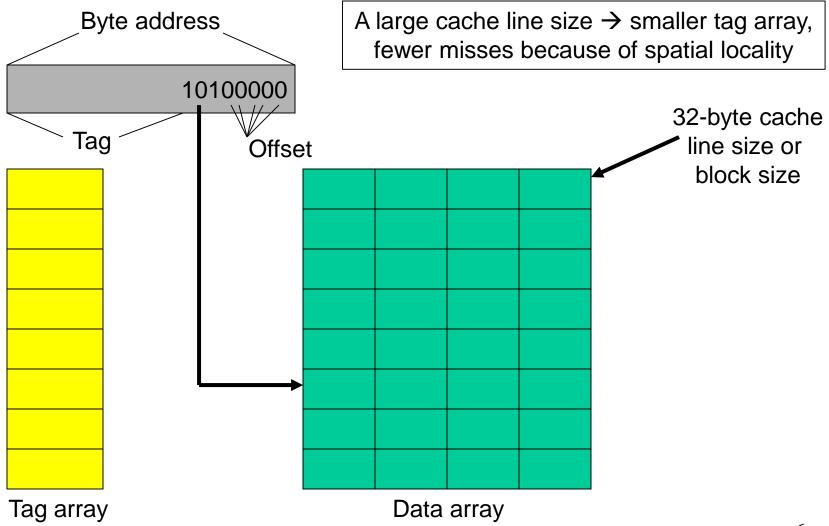
The Tag Array



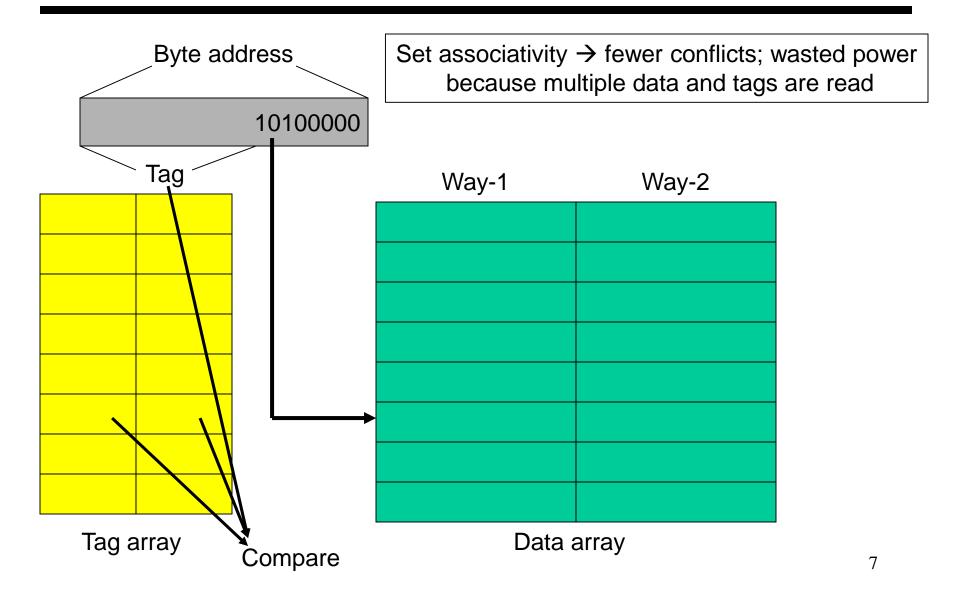
Example Access Pattern



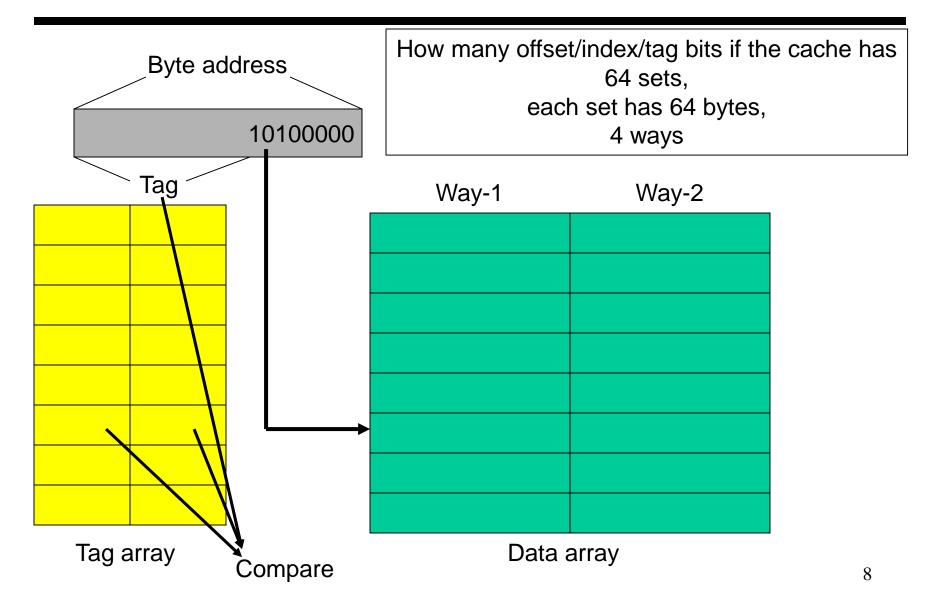
Increasing Line Size



Associativity



Associativity





- 32 KB 4-way set-associative data cache array with 32 byte line sizes
- How many sets?
- How many index bits, offset bits, tag bits?
- How large is the tag array?

- On a write miss, you may either choose to bring the block into the cache (write-allocate) or not (write-no-allocate)
- On a read miss, you always bring the block in (spatial and temporal locality) – but which block do you replace?
 - > no choice for a direct-mapped cache
 - > randomly pick one of the ways to replace
 - replace the way that was least-recently used (LRU)
 - FIFO replacement (round-robin)

Writes

- When you write into a block, do you also update the copy in L2?
 - > write-through: every write to $L1 \rightarrow$ write to L2
 - write-back: mark the block as dirty, when the block gets replaced from L1, write it to L2
- Writeback coalesces multiple writes to an L1 block into one L2 write
- Writethrough simplifies coherency protocols in a multiprocessor system as the L2 always has a current copy of data

- Compulsory misses: happens the first time a memory word is accessed – the misses for an infinite cache
- Capacity misses: happens because the program touched many other words before re-touching the same word – the misses for a fully-associative cache
- Conflict misses: happens because two words map to the same location in the cache – the misses generated while moving from a fully-associative to a direct-mapped cache



Bullet