

# Optimizing Stencil Computations

March 18, 2013



## Administrative

Midterm coming

April 3?

In class March 25, can bring one page of notes

Review notes, readings and review lecture

Prior exams are posted

Design Review

Intermediate assessment of progress on project, oral and short

In class on April 1

Final projects

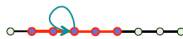
Poster session, April 24 (dry run April 22)

Final report, May 1



## Stencil Computations

A stencil defines the value of a grid point in a d-dimensional spatial grid at time  $t$  as a function of neighboring grid points at recent times before  $t$ .



(a) 1D 5-Point Stencil



(b) 2D 5-Point Stencil



## Stencil Computations, Performance Issues

- Bytes per flop ratio is  $O(1)$
- Most machines cannot supply data at this rate, leading to memory bound computation
- Some reuse, but difficult to exploit fully, and interacts with parallelization

How to maximize performance:

- Avoid extraneous memory traffic, such as cache capacity/ conflict misses
- Bandwidth optimizations to maximize utility of memory transfers
- Maximize in-core performance

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## Learn More: StencilProbe

- See <http://people.csail.mit.edu/skamil/projects/stencilprobe/>
- Several variations of Heat Equation, to be discussed
- Can instantiate to measure performance impact

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## Example: Heat Equation

```
for (t=0; t<timesteps; t++) { // time step loop
  for (k=1; k<nz-1; k++) {
    for (j=1; j<ny-1; j++) {
      for (i=1; i<nx-1; i++) {
        // 3-d 7-point stencil
        B[i][j][k] = A[i][j][k+1] + A[i][j][k-1] +
          A[i][j+1][k] + A[i][j-1][k] + A[i+1][j][k] +
          A[i-1][j][k] - 6.0 * A[i][j][k] / (fac*fac);
      }
    }
  }
  temp_ptr = A;
  A = B;
  B = temp_ptr;
}
```

What if nx,  
ny, nz large?

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## Heat Equation, Add Tiling

```
for (t=0; t<timesteps; t++) { // time step loop
  for (jj = 1; jj < ny-1; jj+=TJ) {
    for (ii = 1; ii < nx - 1; ii+=TI) {
      for (k=1; k<nz-1; k++) {
        for (j = jj; j < MIN(jj+TJ,ny - 1); j++) {
          for (i = ii; i < MIN(ii+TI,nx - 1); i++) {
            // 3-d 7-point stencil
            B[i][j][k] = A[i][j][k+1] + A[i][j][k-1] +
              A[i][j+1][k] + A[i][j-1][k] + A[i+1][j][k] +
              A[i-1][j][k] - 6.0 * A[i][j][k] / (fac*fac);
          }
        }
      }
    }
  }
  temp_ptr = A;
  A = B;
  B = temp_ptr;
}
```

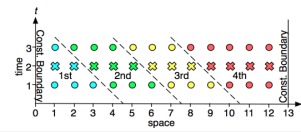
Note the reuse  
across time steps!

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## Heat Equation, Time Skewing

```
for (kk=1; kk < nz-1; kk+=tz) {
  for (jj = 1; jj < ny-1; jj+=ty) {
    for (ii = 1; ii < nx - 1; ii+=tx) {
      for (t=0; t<timesteps; t++) { // time step loop
        ... calculate bounds from t and slope ...
        for (k=blockMin_z; k < blockMax_z; k++) {
          for (j=blockMin_y; j < blockMax_y; j++) {
            for (i=blockMin_x; i < blockMax_x; i++) {
              // 3-d 7-point stencil
              B[i][j][k] = A[i][j][k+1] + A[i][j][k-1] +
                A[i][j+1][k] + A[i][j-1][k] + A[i+1][j][k] +
                A[i-1][j][k] - 6.0 * A[i][j][k] / (fac*fac);
            }
          }
        }
      }
    }
  }
  temp_ptr = A;
  A = B;
  B = temp_ptr;
}
```

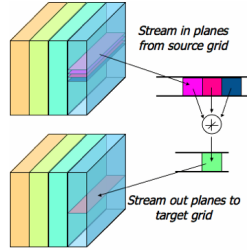


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## Heat Equation, Circular Queue

- See probe\_heat\_circqueue.c

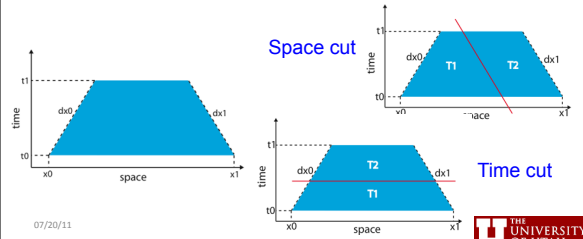


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## Heat Equation, Cache Oblivious

- See probe\_heat\_oblivious.c
- Idea: Recursive decomposition to cutoff point
- Implicit tiling: in both space and time
- Simpler code than complex tiling, but introduces overhead
- Encapsulated in Pochoir DSL (next slide)



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## Example Pochoir Stencil Compiler Specification

```

1 #define mod(r,m) ((r)%m + ((r)<0)? (m):0)
2 Pochoir_Boundary_2D(heat_bv, a, t, x, y)
3 return a.get(1, mod(x,a.size(1)), mod(y,a.size(0)));
4 Pochoir_Boundary_End
5 int main(void) {
6 Pochoir_Shape_2D_2D_five_pt[] = {{1,0,0}, {0,1,0},
7 {0,-1,0}, {0,-1,-1}, {0,0,-1}, {0,0,1}};
8 Pochoir_Array_2D(heat_fn, t, x, y);
9 u.Register_Boundary(heat_bv);
10 heat.Register_Array(a);
11 Pochoir_Kernel_2D(heat_fn, t, x, y)
12 u(t+1, x, y) = CK * (u(t, x+1, y) - 2 * u(t, x,
13 y) + u(t, x-1, y)) + CK * (u(t, x, y+1) - 2
14 * u(t, x, y) + u(t, x, y-1)) + u(t, x, y);
15 Pochoir_Kernel_End
16 for (int x = 0; x < X; ++x)
17 for (int y = 0; y < Y; ++y)
18 u(0, x, y) = rand();
19 heat.Run(T, heat_fn);
20 for (int x = 0; x < X; ++x)
21 for (int y = 0; y < Y; ++y)
22 cout << u(T, x, y);
23 }

```



## Parallel Stencils in Pochoir

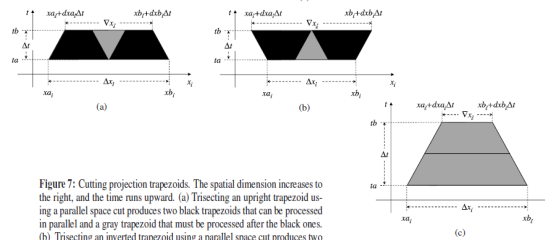


Figure 7: Cutting projection trapezoids. The spatial dimension increases to the right, and the time runs upward. (a) Trisecting an upright trapezoid using a parallel space cut produces two black trapezoids that can be processed in parallel and a gray trapezoid that must be processed after the black ones. (b) Trisecting an inverted trapezoid using a parallel space cut produces two black trapezoids that can be processed in parallel and a gray trapezoid that must be processed before the black ones. (c) A time cut produces a lower and an upper trapezoid where the lower trapezoid must be processed before the upper.

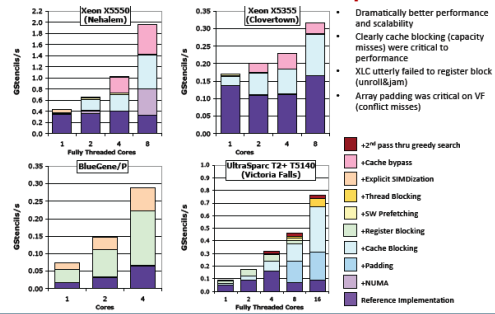


## General Approach to Parallel Stencils

- Always safe to parallelize within a time step
- Circular queue and time skewing encapsulate “tiles” that are independent



## Results for Heat Equation



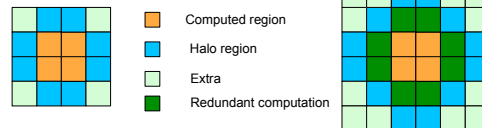
## What about GPUs?

- Two recent papers:
  - “Auto-Generation and Auto-Tuning of 3D Stencil Codes on GPU Clusters,” Zhang and Mueller, CGO 2012.
  - “High-Performance Code Generation for Stencil Computations on GPU Architectures,” Holewinski et al., ICS 2012.
- Key issues:
  - Exploit reuse in shared memory.
  - Avoid fetching from global memory.
  - Thread decomposition to support global memory coalescing.



## Overlapped Tiling for Halo Regions (or Ghost Zones)

- Input data exceeds output result (as in Sobel)
- Halo region or ghost zone extends the per-thread data decomposition to encompass additional input
- An  $(n+2) \times (n+2)$  halo region is needed to compute an  $n \times n$  block if subscript expressions are of the form  $\pm 1$ , for example
- By expanding the halo region, we can trade off redundant computation for reduced access to global memory when parallelizing across time steps.

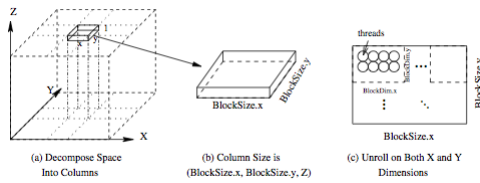


2-d 5-point stencil example



## 2.5D Decomposition

- Partition such that each thread block sweeps over the z-axis and processes one plane at a time.



## Resulting Code

```
#define sizeX (BLOCK_X*halo_x+2)
#define sizeY (BLOCK_Y*halo_y+2)
template <class T>
__kernel__ stencil_iteration(...)
{
    // Initialization Instructions
    g_tx = ...; g_ty = ...;
    ...
    __shared__ T shArr[3][sizeX][sizeY];
    first = 0; second = 1; third = 2;
    // Load first 2 planes
    shArr[0][1][1] = ...;
    shArr[1][1][1] = ...;
    for (k=halo_k; k<sizeY-k+1) {
        // Load third plane to __shared__
        shArr[2][1][1] = ...;
        // Shift registers
        step = middle;
        middle = below;
        // Load third plane to registers
        T below = ...;
        // stencil calculation
        ...
        __syncthreads();
        // Shift planes
        first = (first+1)%3;
        second = (second+1)%3;
        third = (third+1)%3;
    }
}
```

(a) With Corner Accesses (b) Without Corner Accesses

Figure 5: Stencil Kernel Templates

## Other Optimizations

- X dimension delivers coalesced global memory accesses
  - Pad to multiples of 32 stencil elements
- Halo regions are aligned to 128-bit boundaries
- Input (parameter) arrays are padded to match halo region, to share indexing.
- BlockSize.x is maximized to avoid non-coalesced accesses to halo region
- Blocks are square to reduce area of redundancy.
- Use of shared memory for input.
- Use of texture fetch for input.

## Performance Across Stencils and Architectures

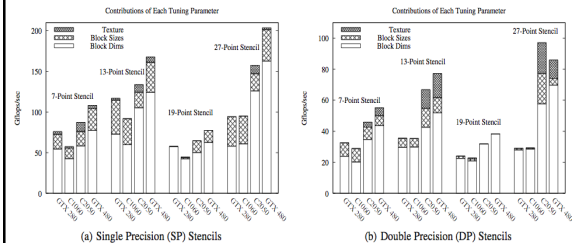


Figure 8: Stencil Tuning Effect Breakups

# GPU Cluster Performance

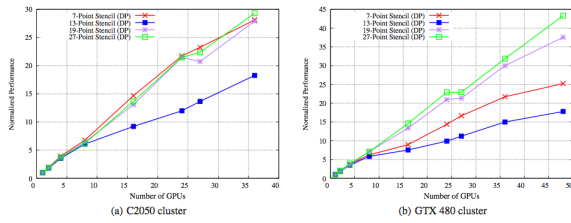


Figure 11: Weak Scaling of DP Stencils on GPU Clusters

