## CS 4400

## Computer Systems

## LECTURE 8

Array allocation and access

## Arrays in C

- Array declaration $T$ A $[N]$;
- allocates a contiguous region of $L \cdot N$ bytes, $L$ is the size of $T$
- introduces A as a constant pointer to the beginning of the array
- Let $\mathrm{x}_{\mathrm{A}}$ be address stored in A , element $i$ is stored at $\mathrm{x}_{\mathrm{A}}+L \cdot i$.
- With IA32's flexible addressing modes, translation to assembly code is straightforward.
- suppose E is array of int's with its address in \%edx and i in \%ecx

$$
\text { movl }(\% e d x, \% e c x, 4), \% e a x \text { stores E[i] in \%eax }
$$

- optimizing compilers are particularly good at simplifying address computations, which may make assembly code hard to read


## Pointer Arithmetic

- Computed value is scaled according to size of data type.
- for int* $p$, expression $p+k$ has value $x_{p}+4 \cdot k$
- for char* str, what is the value of expression $s t r+j$ ?
- Array subscripting operation can be applied to array names and other pointers.
- $A[i]$ equivalent to * $(A+i)$
- Examples (\%edx: address of $E$, \%ecx: value of $i$, \%eax: result):
- E[2] movl 8 (\%edx), \%eax
- E+i-1 leal -4 (\%edx, \%ecx,4), \%eax
- $*(\& E[i]+i) \quad ? ?$

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## Exercise: Pointer Arithmetic

- Let the address of short $S[]$ be in \%edx and index i be in \%ecx.
- Put a pointer result in \%eax, and a short result in \%ax.

|  | type | value | assembly code |
| :---: | :---: | :---: | :---: |
| S+1 | short* | $\mathrm{x}_{\mathrm{s}}+2$ | leal 2 (\%edx), \%eax |
| S [3] |  |  |  |
| \&S [i] |  |  |  |
| $S[4 * i+1]$ |  |  |  |
| S+i-5 |  |  |  |
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## Clicker Question

If you have ResponseCard clicker, channel is 41.
If you are using ResponseWare, session id is CS1400U.
Suppose we have declared int arr [N]. Which of the following is equivalent to the reference arr[i]?
A. *(arr + 4 * i)
B. *(\&arr[0] + i)
C. *((int*) ((char*)arr + 4 * i))
D. exactly 2 of the above
E. all of A-C
F. none of A-C

## Clicker Question

Suppose we have declared char* arr [N]. Which of the following correctly puts arr [i] in \%eax? (Suppose that arr in \%edx and i in \%eax.)

$$
\begin{aligned}
& \text { A. leal (\%edx, \%eax), \%eax } \\
& \text { B. leal (\%edx, \%eax, 4), \%eax } \\
& \text { C. movl (\%edx, \%eax), \%eax } \\
& \text { D. movl (\%edx, \%eax, 4), \%eax } \\
& \text { E. none of the above }
\end{aligned}
$$

## Arrays and Loops

- Array references in loops often have very regular patterns.

```
for(i = 0, val = 0; i < 5; i++)
    val = (10 * val) + x[i];
```

- For efficiency, optimizing compilers exploit these patterns.

```
        xorl %eax,%eax
        leal 16(%ecx),%ebx
.L12:
    addl $4,%ecx ;x++
    jbe .L12
```

    leal (\%eax, \%eax,4), \%edx ;compute 5*val
    movl (\%ecx) , \%eax ; compute *x
    leal (\%eax, \%edx, 2 ), \%eax ;compute *x+2* (5*val)
    cmpl \%ebx, \%ecx ;compare \(x\)-xend
    ```
; val=0
; xend=x+4
;if \(x<=x e n d, ~ g o t o l o o p\)
```

- Uses pointer arithmetic instead of loop index i.

```
int* xend \(=x+4\);
do \{
    val \(=(10\) * val) + *x;
\} while(++x <= xend);
```


## Nested Arrays

- The same principles hold for arrays of arrays.
- int A [4] [3]; is an array of four 3-integer arrays ("rows")
- arrays are linearized in memory in row-major order
- A[i][j] is at memory address $x_{A}+L(C \cdot i+j)$.
- Example (\%eax: address of A, \%edx: value of $i$, \%ecx: value of $j$ )

```
sall $2,%ecx
; j*4
leal (%edx,%edx,2),%edx ;i*3
leal (%ecx,%edx,4),%edx ;j*4 + i*12
movl (%eax,%edx),%eax ;read A[i][j]
```

- Exercise: Compute the address of the second row.

```
#define N 16
typedef int fix_matrix[N][N]; fixed-size array
int fix_prod(fix_matrix A, fix_matrix B, int i, int k) {
    int j, result;
    for(j = 0, result =0; j < N; j++)
        result += A[i][j] * B[j][k];
    return result;
}
```

```
int fix_prod(fix_matrix A, fix_matrix B, int i, int k) {
    int *Aptr, *Bptr, cnt, result;
    Aptr = &A[i][0];
    Bptr = &B[0][k];
    cnt = N-1;
    result = 0;
compiler optimizations
    do {
        result += (*Aptr) * (*Bptr);
        Aptr++;
        Bptr += N;
        cnt--;
    } while(cnt >= 0);
    return result;
}
```

Aptr is in \%edx
Bptr is in \%ecx
result is in \%esi
cnt is in \%ebx
.L23:
movl (\%edx), \%eax imull (\%ecx), \%eax addl \%eax,\%esi
addl \$64,\%ecx
addl \$4,\%edx decl \%ebx jns.L23

## Exercise: Nested Arrays

```
#define M ??
#define N ??
int mat1[M][N];
int mat2[N][M];
int sum_element(int i, int j) {
    return mat1[i][j] + mat2[j][i];
}
```

```
movl 8(%ebp),%ecx
```

movl 8(%ebp),%ecx
movl 12(%ebp),%eax
movl 12(%ebp),%eax
leal 0(,%eax,4),%ebx
leal 0(,%eax,4),%ebx
leal 0(,%ecx,8),%edx
leal 0(,%ecx,8),%edx
subl %ecx,%edx
subl %ecx,%edx
addl %ebx,%eax
addl %ebx,%eax
sall \$2,%eax
sall \$2,%eax
movl mat2(%eax,%ecx,4),%eax
movl mat2(%eax,%ecx,4),%eax
movl mat1 (%ebx,%edx,4),%eax

```
movl mat1 (%ebx,%edx,4),%eax
```


## Clicker Question

The following will compile (gcc) without error or warning.

```
#define N 100
int foo(int arr[][N], int i, int j) {
    return arr[i][j];
}
```

A. true
B. false

## New to C?: Dynamic Memory Alloc

- For allocation of memory at run time, library routine malloc is used.
- arguments specify number of bytes to be allocated
- return value is a pointer to the allocated memory or NULL
- malloc allocates one contiguous block (of specified size).

$$
\begin{aligned}
& \text { NODE* head }=\text { malloc(sizeof(NODE)); // implicit } \\
& \text { head->next }=\text { malloc(sizeof(NODE)); // cast }
\end{aligned}
$$

- To release dynamically-allocated memory, the library routine free is used.
- argument is the pointer to the block of memory to be released
free(ptr);


## Clicker Question

Suppose we have
short* arr = malloc(user_input*sizeof(short));
Which of the following references the second element?
A. arr[1]
B. *(arr+1)
C. *(arr+2)
D. exactly 2 of the above
E. all of A-C
F. none of A-C

## Clicker Question

Suppose we have
short* matrix = malloc(N*N*sizeof(short));

Which of the following references the element in the second row and second column?
A. arr[1]
B. $\operatorname{arr}[\mathrm{N}]$
C. arr $[\mathrm{N}+1]$
D. arr[1][1]
E. none of the above

## Exercise: Compiler Optimizations

```
#define N 16
typedef int fix_matrix[N][N];
    int i;
    for (i = 0; i < N; i++)
        A[i][i] = val;
}
```

void fix_set_diag(fix_matrix A, int val) \{

Write a function
fix_set_diag_opt that uses optimizations similar to those in the assembly code. Do not assume that N is 16 .
movl 8 (\%ebp), \%ecx movl 12 (\%ebp), \%edx movl \$0, \%eax .L14:
movl \%edx, (\%ecx, \%eax)
addl \$68,\%eax
cmpl \$1088,\%eax
jne. .L14

