Terminology: Denoted and Expressed Values

- A ***denoted value*** is the meaning of a variable
- An ***expressed value*** is the result of an expression

The set of denoted and expressed values can be different

Terminology: Denoted and Expressed Values

- First-order functions
  - denoted values: numbers and functions
  - expressed values: numbers
- Higher-order functions
  - denoted values: numbers and functions
  - expressed values: numbers and functons

Procedure Expressions: Concrete Syntax

```
<prog>   ::=  <expr>
<expr>   ::=  proc (<id>*(,)) <expr>
         ::=  (<expr> <expr>*)
```

```
let identity = proc(x) x
 in (identity 5)
```

Procedure Expressions: Abstract Syntax

```
<prog>   ::=  (a-program <expr>)
<expr>   ::=  (proc-exp (list <id>*) <expr>)
         ::=  (app-exp <expr> (list <expr>*))
<val>    ::=  <num>   |   <proc>
<proc>   ::=  (closure (list <id>*) <expr> <env>)
```

```
(a-program
  (let-exp (list 'identity)
        (list (proc-exp (list 'x) (var-exp 'x)))
        (app-exp (var-exp 'identity) (list-exp 5))))
```

Implementing Procedures

(implementation in DrScheme)

New representation of environments:

```
(define-datatype environment environment?
  (empty-env-record)
  (extended-env-record
    (syms (list-of symbol?))
    (vals (list-of denval?))
    (env environment?)))
```

Recursion

Suppose we try to write the **fact** function using only **let**

let fact = **proc**(**n**) **if n then** *(**n**, (**fact** -(**n**, 1))) **else** 1
in (**fact** 10)

The above doesn't work, because **fact** is not bound in the local function

We'll add **letrec**, but first we'll see how to implement **fact** without it...

Recursion with Let

- Problem: **fact** can't see itself
- Note: anyone calling **fact** can see **fact**
- Idea: have the caller supply **fact** to **fact** (along with a number)

let fact = **proc**(**n**, **f**) **if n then** *(**n**, (**f** -(**n**, 1) **f**)) **else** 1
in (**fact** 10 **fact**)

this works!

What Happened?

- The key insight is delaying some work to the caller
- We can exploit this idea to implement **letrec**, but in a slightly different way
- **letrec** requires a *closure* that refers to itself
- We can delay the actual construction of the closure until it is extracted from the environment

Recursive Environments for Recursive Functions

This isn't going to work

```
let fact = proc(n) if n then *(n, (fact -(n, 1))) else 1
 in (fact  10)
```

Recursive Environments for Recursive Functions

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let fact = proc(n) if n then *(n, (fact -(n, 1))) else 1
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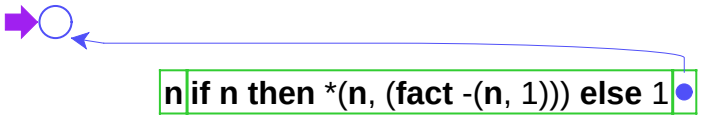
Recursive Environments for Recursive Functions

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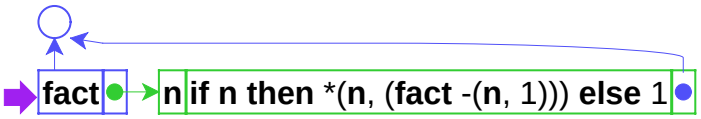
Recursive Environments for Recursive Functions

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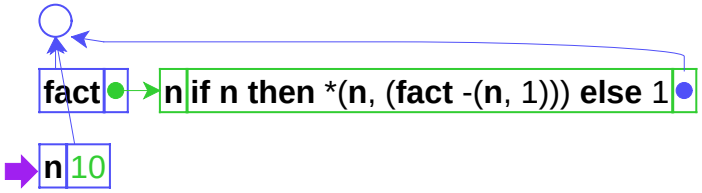
Recursive Environments for Recursive Functions



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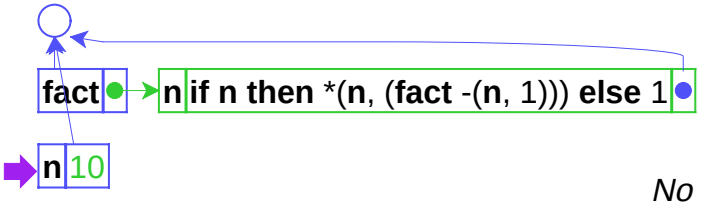
Recursive Environments for Recursive Functions



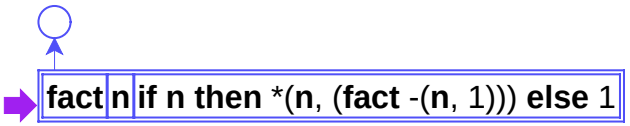
No binding for **fact**

let fact = **proc**(**n**) **if n then** *(**n**, (**fact** -(**n**, 1))) **else** 1
in (**fact** 10)

Recursive Environments for Recursive Functions

letrec fact = **proc**(**n**) **if n then** *(**n**, (**fact** -(**n**, 1))) **else** 1
in (**fact** 10)

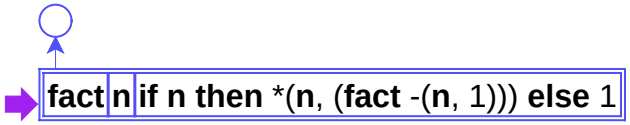
Recursive Environments for Recursive Functions



double box means a recursively extended environment

letrec fact = **proc**(**n**) **if n then** *(**n**, (**fact** -(**n**, 1))) **else** 1
in (**fact** 10)

Recursive Environments for Recursive Functions



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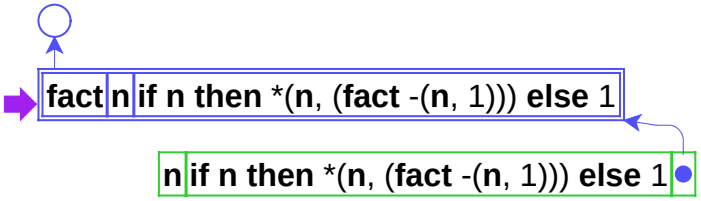
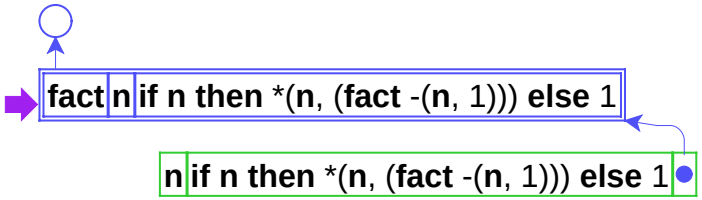
Recursive Environments for Recursive Functions



every lookup of **fact**
generates a closure

letrec fact = **proc**(**n**) **if n then** *(**n**, (**fact** -(**n**, 1))) **else** 1
in (**fact** 10)

Recursive Environments for Recursive Functions



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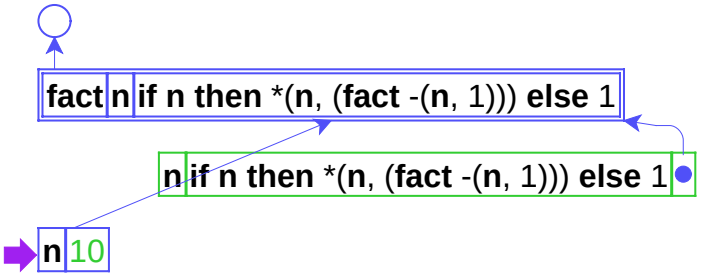
Recursive Environments for Recursive Functions



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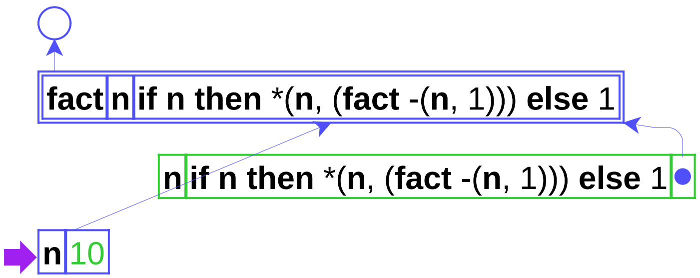

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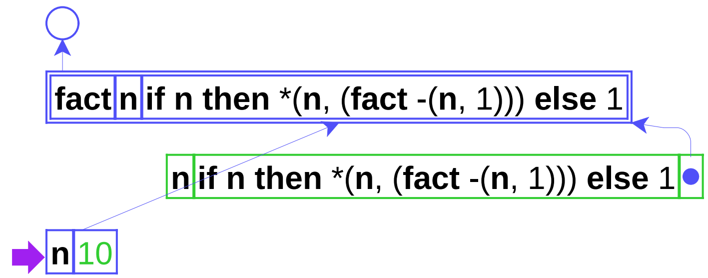
Recursive Environments for Recursive Functions



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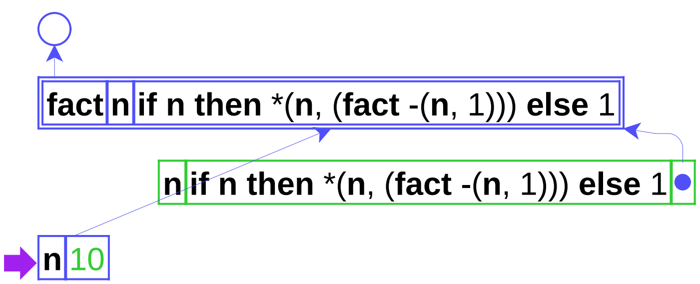
Recursive Environments for Recursive Functions



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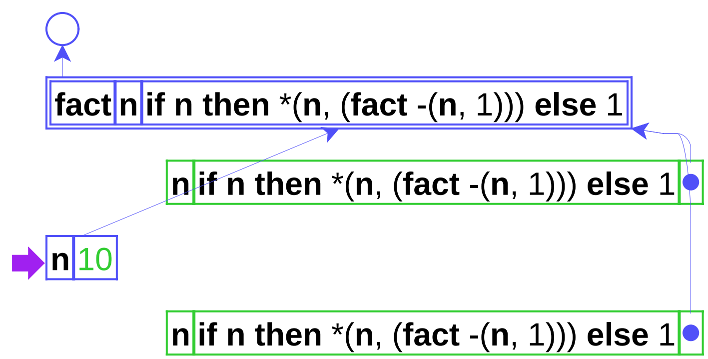
Recursive Environments for Recursive Functions



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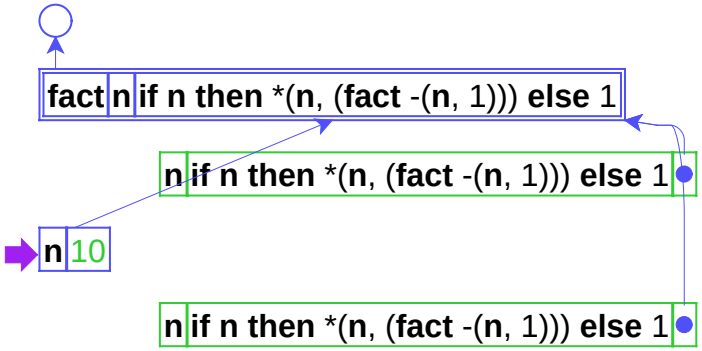
Recursive Environments for Recursive Functions



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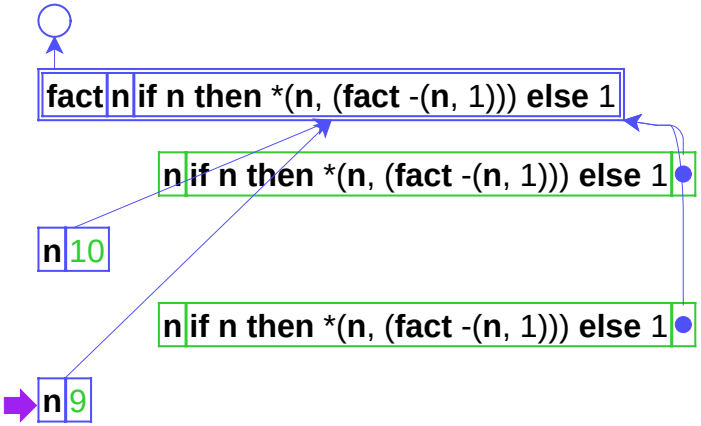
Recursive Environments for Recursive Functions



letrec fact = **proc**(**n**) **if n then** *(**n**, (**fact** -(**n**, 1))) **else** 1
in (**fact** 10)

Implementing Recursively Extended Envs

```
(define-datatype environment environment?
  (empty-env-record)
  (extended-env-record
    (syms (list-of symbol?))
    (vec vector?)
    (env environment?))
  (recursively-extended-env-record
    (proc-names (list-of symbol?))
    (idss (list-of (list-of symbol?)))
    (bodies (list-of expression?))
    (env environment?)))
```

Implementing letrec

(implement in DrScheme)

Back to Recursion with Let...

- Allowing functions to be values is a powerful idea
- As it turns out, we don't even need **let** !

let <id>$_1$ = <expr>$_1$ **...** <id>$_n$ = <expr>$_n$ **in** <expr>

is the same as

(**proc**(<id>$_1$, **...** <id>$_n$) <expr> <expr>$_1$ **...** <expr>$_n$)

Back to Recursion with Let...

- Allowing functions to be values is a powerful idea
- As it turns out, we don't even need **let** !

(**let** ([<id>$_1$ <expr>$_1$] **...** [<id>$_n$ = <expr>$_n$]) <expr>)

is the same as

((**lambda** (<id>$_1$ **...** <id>$_n$) <expr>) <expr>$_1$ **...** <expr>$_n$)

The Lambda Calculus

- We don't even need functions of multiple arguments...

((**lambda** (<id>$_1$ **...** <id>$_n$) <expr>)
 <expr>$_1$ **...** <expr>$_n$)

is the same as

(((**lambda** (<id>$_1$) **...** (**lambda** (<id>$_n$) <expr>))
 <expr>$_1$) **...**
 <expr>$_n$)

Passing multiple arguments one-at-a-time is called ***currying***

The ***lambda calculus*** has only single-argument **lambda** and single-argument function calls, and it's computationally complete