## Writing Functions in Scheme

- Suppose we want a function ct which takes a list of symbols and returns the number of symbols in the list

$$
\begin{gathered}
\left(\text { ct }^{\prime}(\mathbf{a b c} \mathbf{b}) \rightarrow \rightarrow 3\right. \\
\left(\text { ct }^{\prime}()\right) \rightarrow \rightarrow 0 \\
\left(\text { ct }^{\prime}\left(\mathrm{x}_{\mathrm{y}} \mathbf{z w t} \mathbf{~ w ~}\right) \rightarrow \rightarrow 5\right.
\end{gathered}
$$

How can we write this function?

## Writing Functions in Scheme

- Answer \#1: Have the instructor write it

```
;; ct : <list-of-sym> -> <num>
;;(ct '()) }->->
;;(ct '(a b c)) }->->
(define (ct I)
(cond
    [(null? I) 0]
    [else (+1 (ct (cdr I)))]))
```


## Checking My Answer: Empty List

| (define (ct l) | $\rightarrow \quad$ (define (ct I) |
| :---: | :---: |
| (cond | (cond |
| [(null? I) 0] | [(null? I) 0] |
| [else (+ 1 (ct (cdr I) ) )] ) | [else (+ 1 (ct (cdr I) ) )] ) |
| (cond | (cond |
| [(null? '()) 0] | [\#t 0] |
| [else (+ 1 (ct (cdr $\left.{ }^{\prime}()\right)$ ) $]$ ) | [else (+ 1 (ct (cdr $\left.{ }^{\prime}()\right)$ ) $)$ ]) |


| ```(define (ct l) (cond [(null? I) 0] [else (+ 1 (ct (cdr I))]))``` | $\begin{aligned} \rightarrow \quad & (\text { define }(\text { ct I) }) \\ & (\text { cond } \\ & {[(\text { null? I) } 0]} \\ & {[\text { else }(+1 \text { (ct (cdr I)) })])) } \end{aligned}$ |
| :---: | :---: |
| (cond <br> [\#t 0] <br> $[$ else $(+1$ (ct (cdr '())))]) | 0 |


| ```(define (ct l) (cond [(null? I) 0] [else (+ 1 (ct (cdr I)))])``` | $\begin{aligned} \rightarrow \quad & (\text { define }(\text { ct I) }) \\ & (\text { cond } \\ & {[(\text { null? I) } 0]} \\ & {[\text { else }(+1(\text { ct }(\text { cdr I I)) })])) } \end{aligned}$ |
| :---: | :---: |
| (ct ' ${ }^{(a b c}$ ) | $\begin{aligned} & \text { (cond } \\ & {[(\text { null? ' }(\mathbf{a b c} \mathbf{~ b})) 0]} \\ & {\left[\text { else }\left(+1\left(\mathbf{c t}\left(\mathbf{c d r}{ }^{\prime}(\mathbf{a b c} \mathbf{~ c})\right)\right)\right]\right)} \end{aligned}$ |

Checking My Answer: List of 3 Symbols

| (define (ct I) | $\rightarrow \quad$ (define (ct I) |
| :---: | :---: |
| (cond | (cond |
| [(null? I) 0] | [(null? I) 0] |
| [else (+ 1 (ct (cdr I) ) $]$ )) | $[$ else (+ 1 (ct (cdr I) ) )] ) |
| (cond | (cond |
| [(null? '(a b c)) 0] | [\#f 0] |
| [else (+ 1 (ct (cdr '(a b c) )) )]) | [else (+1 (ct (cdr ' ${ }^{\text {( } \mathbf{b} \mathbf{~ c ~ c ~})}$ ) ) $)$ ]) |

## Checking My Answer: List of 3 Symbols

| (define (ct l) | $\rightarrow \quad$ (define (ct I) |
| :---: | :---: |
| (cond | (cond |
| [(null? I) 0] | [(null? I) 0] |
| [else (+ 1 (ct (cdr I) ) )] ) | [else (+ 1 (ct (cdr I) ) )] ) |
| (cond | $\left(+1\left(\mathbf{c t ~}\left(\mathbf{c d r}{ }^{\prime}(\mathbf{a b c} \mathbf{~ c})\right)\right.\right.$ ) |
| [\#f 0] |  |
|  |  |

Checking My Answer: List of 3 Symbols

```
(define (ct l)
    (cond
    [(null? I) 0]
    [else (+ 1 (ct (cdr I))]])
(+ 1 (ct (cdr '(a b c))))
(define (ct I)
    (cond
        [(null? I) 0]
        [else (+ 1 (ct (cdr I))]])
(+1
    (ct '(b c)))
```


## Checking My Answer: List of 3 Symbols

```
(define (ct I) }->\underset{\mathrm{ (cond (define (ct I)}}{\mathrm{ (cond }
    [(null? I) 0]
    [else (+ 1 (ct (cdr I)))]))
(+1
    (ct '(b c)))
```


## Checking My Answer: List of 3 Symbols

(define (ct I)
(cond
[(null? I) 0] [else (+ 1 (ct (cdr I)) )]))
(+ 1

## (cond

[(null? '(b c)) 0]
$\left[\right.$ else (+ $\left.\left.\left.\left.1\left(\mathbf{c t ~ ( c d r}{ }^{\prime}(\mathbf{b} \mathbf{c})\right)\right)\right]\right)\right)$
$\rightarrow \quad$ (define (ct I)
(cond
[(null? I) 0]
$[$ else $(+1($ ct $($ cdr I) $))]))$
(+ 1
(cond
[\#f 0]
$[$ else $(+1(c t(c d r \quad$ ' $\mathbf{b ~ c} \mathbf{c})))]))$

## Checking My Answer: List of 3 Symbols

| ```(define (ct I) (cond [(null? I) 0] [else (+ 1 (ct (cdr I)))]))``` | $\begin{aligned} \rightarrow \quad & (\text { define }(\text { ct I) } \\ & (\text { cond } \\ & {[(\text { null? I) } 0]} \\ & {[\text { else }(+1(\text { ct }(\text { cdr I) }))])) } \end{aligned}$ |
| :---: | :---: |
| (+1 | (+ 1 |
| (cond | (+1 |
| [\#f 0] | (ct (cdr ' $\left.{ }^{(b \mathbf{c}}\right)$ )) ) |
| [else (+1 (ct (cdr ' $\left.{ }^{\text {(b c c })}\right)$ ) $\left.)\right]$ ) |  |

Checking My Answer: List of 3 Symbols

```
(+ 1
```

(+ 1

```
(+ 1
    (+1
    (+1
    (+1
    (ct (cdr '(b c)))))
```

    (ct (cdr '(b c)))))
    ```
    (ct (cdr '(b c)))))
```


## $\rightarrow \quad$ (define (ct I)

``` (cond
[(null? I) 0] [else (+ 1 (ct (cdr I)))]))
(+ 
    (+1
    (ct '(c))))
```


## Checking My Answer: List of 3 Symbols

```
(define (ct I)
    (cond
    [(null? I) 0]
    [else (+ 1 (ct (cdr I)))]))
```

```
++1
```

++1
(+1
(+1
(ct '(c))))

```
    (ct '(c))))
```

$\rightarrow \quad$ (define (ct I) (cond [(null? I) 0] [else (+ 1 (ct (cdr I)))]))
(+ 1
(+ 1
(cond
[(null? '(c)) 0] $\left[\right.$ else (+ $1\left(\right.$ ct $\left.\left.\left.\left.\left.\left(\mathbf{c d r}{ }^{\prime}(\mathbf{c})\right)\right)\right)\right]\right)\right)$

## Checking My Answer: List of 3 Symbols

## Checking My Answer: List of 3 Symbols

```
    (cond
    [(null? I) 0]
    [else (+ 1 (ct (cdr I)))]))
(+1
    (+1
        (cond
        [#f 0]
        | (define (ct I)
        (cond
        [(null? I) 0]
        [else (+ 1 (ct (cdr I)))]))
        + }
        (+1
        (+1
        [else (+ 1 (ct (cdr '(c))))])))
            (ct (cdr '(c))))))
```

Checking My Answer: List of 3 Symbols

| ```(define (ct I) (cond [(null? I) 0] [else (+ 1 (ct (cdr I)))]))``` | $\begin{aligned} \rightarrow \quad & \text { (define }(\text { ct I) } \\ & (\text { cond } \\ & {[(\text { null? I) } 0]} \\ & {[\text { else }(+1(\text { ct }(\text { cdr I) }))])) } \end{aligned}$ |
| :---: | :---: |
| (+1 | (+1 |
| (+1 | (+1 |
| (+1 | (+1 |
| (ct (cdr '(c) )) )) | (ct '()) )) ) |

## Checking My Answer: List of 3 Symbols

```
\(\begin{array}{ll}\text { (define (ct I) } & \rightarrow \underset{(\text { define }(c t ~ I)}{(\text { cond }}\end{array}\)
    [(null? I) 0]
    [else (+ 1 (ct (cdr I)))]))
+ +
    (+
        (+ 1
            (ct '()))))
```


## Checking My Answer: List of 3 Symbols

## Checking My Answer: List of 3 Symbols



Checking My Answer: List of 3 Symbols

| (define $($ ct I) | $\rightarrow$ |
| :--- | :---: |
| $($ define $($ ct I) $)$ |  |
| $[($ null? I) 0$]$ | $[($ null? I) 0$]$ |
| $[$ else $(+1($ ct $($ cdr II) $)]))$ | $[$ else $(+1($ ct $($ cdr I) $))]))$ |
| $(+1$ | $(+1$ |
| $(+1$ | $(+1$ |
| $(+1$ | $1))$ |
| $0)))$ |  |

## Checking My Answer: List of 3 Symbols

| define $($ ct I) <br> $($ cond <br> $[($ null? I) 0$]$ | $($ define $($ ct I) |
| :--- | :--- |
| $[$ else $(+1($ ct $($ cdr II) $)]))$ | $[($ null? I) 0$]$ |
|  | $[$ else $(+1($ ct $($ cdr I) $))]))$ |
| $(+1$ | $(+1$ |
| $(+1$ | $2)$ |

1))

## Writing Functions in Scheme: Answer \#2

Answer \#2: Use the general design recipe

- Locate or write a data definition
- Write a contract
- Write examples
- Create a template that follows the shape of the data definition
- Convert the template to the final function
- Run examples as tests


## Writing Functions in Scheme: Answer \#2

Answer \#2: Use the general design recipe

- Locate or write a data definition
- Write a contract
- Write examples
- Create a template that follows the shape of the data definition
- Convert the template to the final function
- Run examples as tests


## works $90 \%$ of the time

## Contracts

A contract is a comment that identifies set of input values and output values
;; ct: <list-of-sym> -> <num>

- All mentioned data sets should have a data definition somewhere


## Data Definitions

What is a "list of symbols"?

```
<list-of-sym> ::= '()
    ::= (cons <symbol> <list-of-sym>)
```

- Sometimes the data definition is given, somtimes you have to create it
- Usually include it in your code as a comment


## Examples

Examples (usually in comments at first) help clarify the purpose of the function

$$
\begin{aligned}
& \because\left(\text { ct }^{\prime}()\right) \rightarrow \rightarrow 0 \\
& ; \because\left(\text { ct }^{\prime}(\mathbf{a} \mathbf{b} \mathbf{c})\right) \rightarrow \rightarrow 3
\end{aligned}
$$

- Make sure that every case in the data definition is covered at least once


## Template

A template reflects the structure of the input according to the data definition

```
<list-of-sym> ::= '()
    ::= (cons <symbol> <list-of-sym>)
```

(define (ct l)
(cond
[(null? 1) ...]
[(pair? 1) ...(car 1)...(ct (cdr 1))...]))

## Template

A template reflects the structure of the input according to the data definition

```
    <list-of-sym> ::= '()
    ::= (cons <symbol> <list-of-sym>)
```

```
(define (ct l)
    (cond
        [(null? 1) ...]
    [(pair? l) ...(car l)...(ct (cdr l))...]))
```

- Corresponding predicate for each data case


## Template

A template reflects the structure of the input according to the data definition

```
<list-of-sym> ::= '()
    ::= (cons <symbol> <list-of-sym>)
```

```
(define (ct l)
```

    (cond
        [(null? 1) ...]
        [(pair? 1) ...(car l)...(ct (cdr 1))...]))
    - Two cases in data definition implies cond with two cond-lines


## Template

A template reflects the structure of the input according to the data definition

```
    <list-of-sym> ::= '()
    ::= (cons <symbol> <list-of-sym>)
```

(define (ct 1)
(cond
[(null? 1) ...]
[(pair? l) ...(car l)...(ct (cdr l))...]))

- Extract parts in cases with meta-variables


## Template

A template reflects the structure of the input according to the data definition

```
<list-of-sym> ::= '()
    ::= (cons <symbol> <list-of-sym>)
```

(define (ct l)
(cond
[(null? 1) ...]
[(pair? 1) ...(car 1)...(ct (cdr 1))...]))

- Recursive call for self-references in data definition


## Template to Function

Transform template to function line-by-line

```
(define (ct l)
    (cond
        [(null? 1) ...]
        [(pair? l) ...(car l)...(ct (cdr l))...]))
```


## Template

A template reflects the structure of the input according to the data definition

```
<list-of-sym> ::= '()
    ::= (cons <symbol> <list-of-sym>)
```

(define (ct l)
(cond
[(null? 1) ...]
[(pair? 1) ...(car 1)...(ct (cdr 1))...]))

- A template depends only on the input data; it ignores the function's purpose
(Nevertheless, generating a template, which is fairly automatic, usually provides most of the function)


## Template to Function

Transform template to function line-by-line

```
(define (ct l)
    (cond
        [(null? 1) 0]
        [(pair? l) ...(car l)...(ct (cdr l))...]))
```


## Template to Function

Transform template to function line-by-line

```
(define (ct l)
    (cond
        [(null? 1) 0]
        [(pair? l) (+ 1 (ct (cdr l)) )]))
```

- Sometimes, a part of the template isn't needed


## Reminder: Template Steps

- Create a cond expression with one line for each case in the data definition
- Write down a predicate for each case
- For the answer, extract parts in cases with meta-variables
- For each self-reference in the data definition, add a recursive call

Shape of template shape $==$ Shape of data definition

## Reminder: Recipe

- Locate or write a data definition
- Write a contract
- Write examples
- Create a template that follows the shape of the data definition
- Convert the template to the final function
- Run examples as tests


## More Examples

(more examples in class)

## Generalized Recipe

- Locate or write data definitions
- Write contracts
- Write examples
- Create a template that follows the shape of the data definition, one for each data definition
- Convert the templates to the final functions
- Run examples as tests

