

Environments in Picture Form

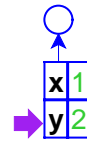


*top purple arrow points to
the current environment*

*purple in bottom area hilites
the current expression*

let x = 1 y = 2
in +(x, y)

Environments in Picture Form



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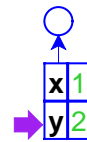
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Environments in Picture Form



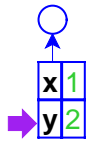
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in (f y)

Environments in Picture Form



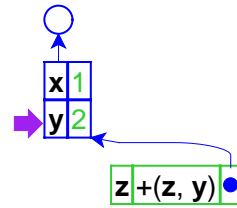
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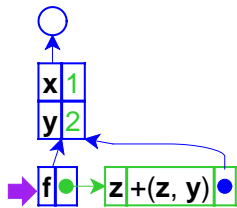
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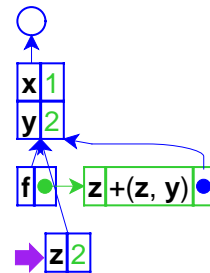
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Environments in Picture Form



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Environments in Picture Form



let x = 1 y = 2
 in let f = proc (z) +(z, y)
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The Need for Recursive Environments



```
let fact = proc(n) if n then *(n, (fact -(n, 1))) else 1
in (fact 10)
```

The Need for Recursive Environments



```
let fact = proc(n) if n then *(n, (fact -(n, 1))) else 1
in (fact 10)
```

The Need for Recursive Environments



```
n if n then *(n, (fact -(n, 1))) else 1
```

```
let fact = proc(n) if n then *(n, (fact -(n, 1))) else 1
in (fact 10)
```

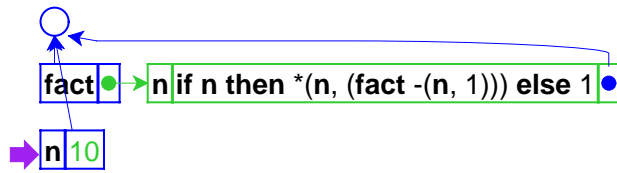
The Need for Recursive Environments



```
fact n if n then *(n, (fact -(n, 1))) else 1
```

```
let fact = proc(n) if n then *(n, (fact -(n, 1))) else 1
in (fact 10)
```

The Need for Recursive Environments



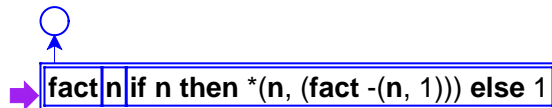
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let fact = proc(n) if n then *(n, (fact -(n, 1))) else 1
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The Need for Recursive Environments



```
letrec fact = proc(n) if n then *(n, (fact -(n, 1))) else 1
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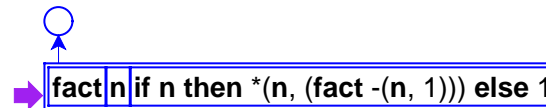
The Need for Recursive Environments



double box means a recursively extended environment

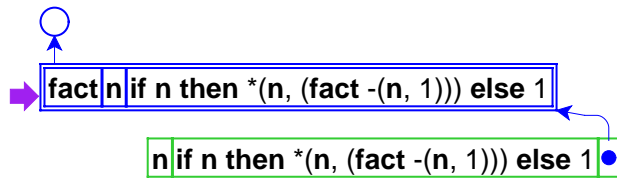
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The Need for Recursive Environments



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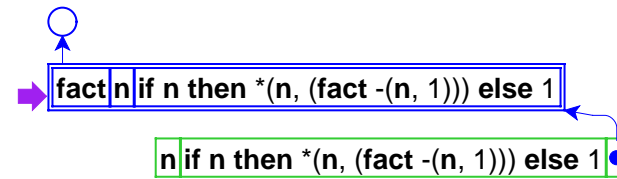
The Need for Recursive Environments



*every lookup of fact
generates a closure*

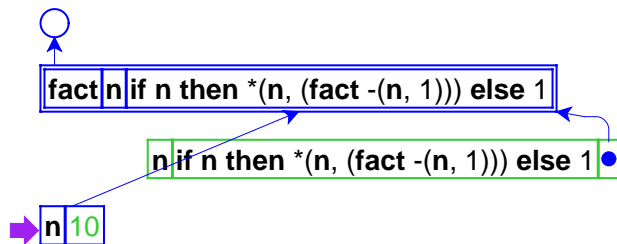
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letrec fact = proc(n) if n then *(n, (fact -(n, 1))) else 1  
in (fact 10)
```

The Need for Recursive Environments



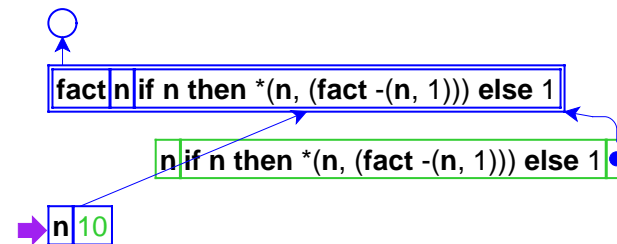
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```

The Need for Recursive Environments



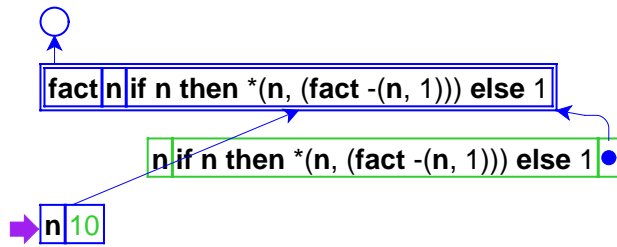
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The Need for Recursive Environments



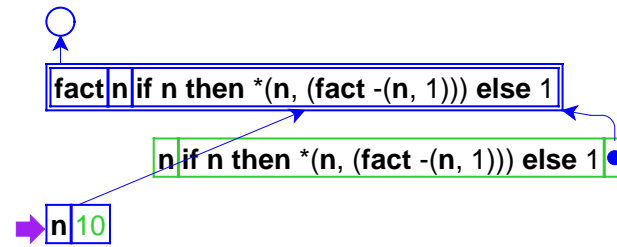
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in (fact 10)
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The Need for Recursive Environments



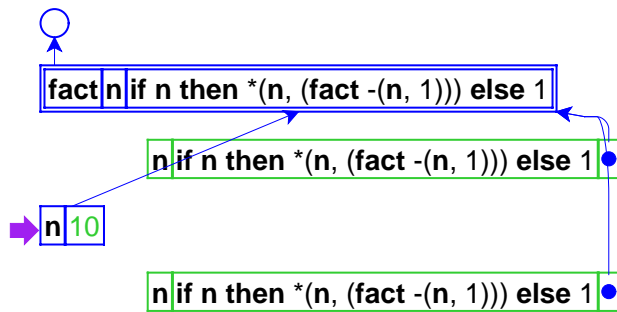
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The Need for Recursive Environments



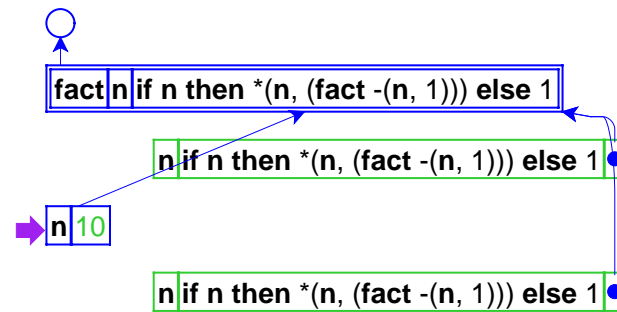
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The Need for Recursive Environments



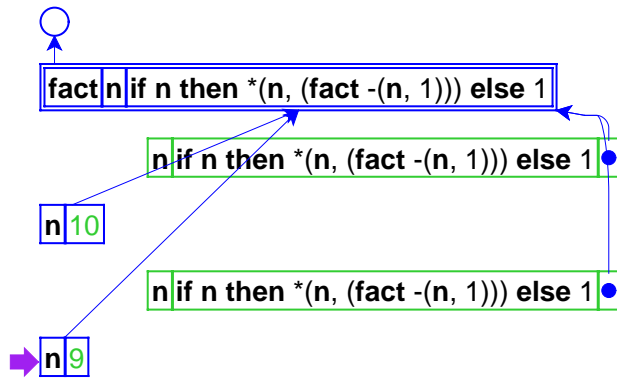
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The Need for Recursive Environments



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letrec fact = proc(n) if n then *(n, (fact -(n, 1))) else 1
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```

Implementing Recursively Extended Envs

(implement in DrScheme)

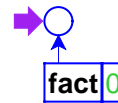
Another Approach to Recursive Closures



alternate approach...

```
letrec fact = proc(n) if n then *(n, (fact -(n, 1))) else 1
in (fact 10)
```

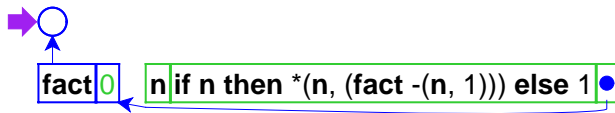
Another Approach to Recursive Closures



*create an environment
with a dummy value...*

```
letrec fact = proc(n) if n then *(n, (fact -(n, 1))) else 1
in (fact 10)
```

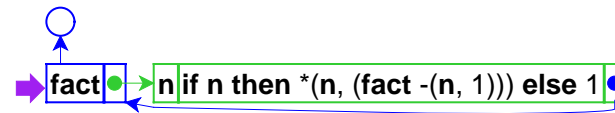
Another Approach to Recursive Closures



create the closure using the environment...

```
letrec fact = proc(n) if n then *(n, (fact -(n, 1))) else 1
in (fact 10)
```

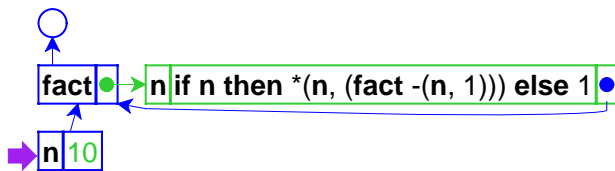
Another Approach to Recursive Closures



then **modify** the environment to fix it up

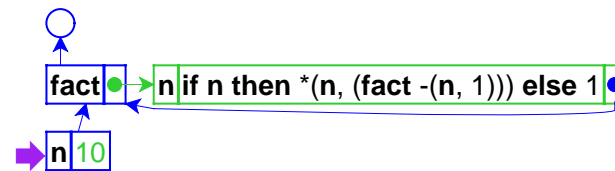
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```

Another Approach to Recursive Closures



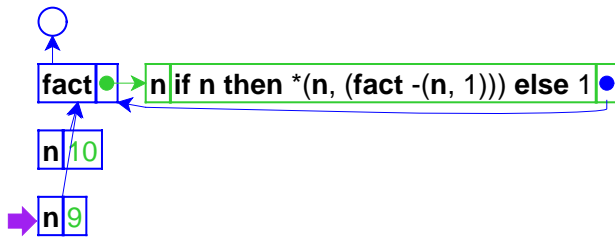
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Another Approach to Recursive Closures



```
letrec fact = proc(n) if n then *(n, (fact -(n, 1))) else 1
in (fact 10)
```


Another Approach to Recursive Closures



*an advantage: closure
is only created once*

```
letrec fact = proc(n) if n then *(n, (fact -(n, 1))) else 1
in (fact 10)
```

Modifying Environments

- The part of the environment that we need to modify is a value in a vector
- So we need evaluation rules to support vector update

Evaluation of Vector Expressions

- Unlike **cons**, **vector** does not create a value
- Instead, it's treated like local functions used to be

```
...
(let ([v (vector 1 2 3)]) (vector-ref v 0))
→
... (define vec1743 (vector 1 2 3))
(let ([v vec1743]) (vector-ref v 0))
→
... (define vec1743 (vector 1 2 3))
(vector-ref vec1743 0)
→
... (define vec1743 (vector 1 2 3))
1
```

Evaluation of Vector Expressions

- The reason for this definition of **vector** is to enable **vector-set!**
- ```
...
(let ([v (vector 1 2 3)]) (begin (vector-set! v 0 5) (vector-ref v 0)))
→
... (define vec1743 (vector 1 2 3))
(let ([v vec1743]) (begin (vector-set! v 0 5) (vector-ref v 0)))
→
... (define vec1743 (vector 1 2 3))
(begin (vector-set! vec1743 0 5) (vector-ref vec1743 0))
→
... (define vec1743 (vector 5 2 3))
(vector-ref vec1743 0)
→
... (define vec1743 (vector 5 2 3))
5
```

## Begin Expressions

- **begin** evaluates a sequence of expressions, in order
- **lambda** and **let** always supply an implicit **begin**

$(\text{let } (\dots) \langle \text{expr} \rangle_1 \dots \langle \text{expr} \rangle_n)$   
=  $(\text{let } (\dots) (\text{begin } \langle \text{expr} \rangle_1 \langle \text{expr} \rangle_n))$

$(\text{lambda } (\dots) \langle \text{expr} \rangle_1 \dots \langle \text{expr} \rangle_n)$   
=  $(\text{lambda } (\dots) (\text{begin } \langle \text{expr} \rangle_1 \langle \text{expr} \rangle_n))$

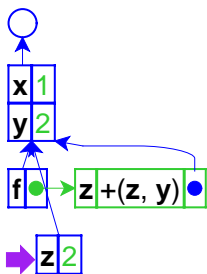
## Changing Recursive Environment Extension

Now we can change `extend-env-recursively` to use **vector-set!**

Go back to just two datatype variants

(implement in DrScheme)

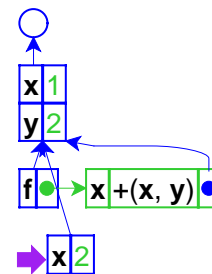
## Back to Lexical Scope



What if we change **z** to **x** ?

```
let x = 1 y = 2
in let f = proc (z) +(z, y)
in (f y)
```

## Back to Lexical Scope

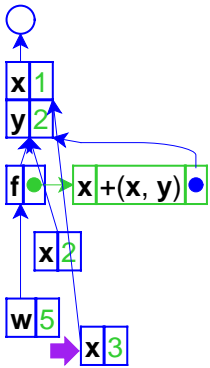


Shape of the environment and location of the argument is unchanged

- argument is always first in first frame
- **y** is always second in second frame

```
let x = 1 y = 2
in let f = proc (x) +(x, y)
in (f y)
```

## Back to Lexical Scope



Still true if **f** is called from a more complex environment

```
let x = 1 y = 2
in let f = proc (x) +(x, y)
in +((f y), let w = 5 in (f 3))
```

## Compilation

So why waste time searching the environment on every variable access?

A compiler can determine the **lexical offset** for each variable statically

Terminology:

- A **compiler** translates a program from language *X* to language *Y*
- An **interpreter** executes a program in language *X*

## Compilation of Variable Accesses

- We'll write a compiler that transforms

```
let x = 1 y = 2
in let f = proc (x) +(x, y)
in (f x)
```

to

```
let = 1 = 2
in let = proc (x) +(<0,0>, <1,1>)
in (<0,0> <1,0>)
```

- We'll also need an interpreter for the new language