

## Outline

- ▶ Data Definitions and Templates
- ▶ Syntax and Semantics
- ▶ Defensive Programming

## Data Definitions

### Question 1:

Are both of the following data definitions ok?

```
; A w-grade is either
; - num
; - posn
; - empty

with ; A posn is
      ; (make-posn num num)

; A z-grade is either
; - num
; - (make-posn num num)
; - empty
```

Yes.

## Data Definitions

### Question 2:

Do **w-grade** and **z-grade** identify the same set of values?

```
; A w-grade is either
; - num
; - posn
; - empty

with ; A posn is
      ; (make-posn num num)

; A z-grade is either
; - num
; - (make-posn num num)
; - empty
```

Yes, every **w-grade** is a **w-grade**,  
and every **z-grade** is a **w-grade**

## Data Definitions

### Question 3:

Are **w-grade** and **w-grade** the same data definition?

```
; A w-grade is either
; - num
; - posn
; - empty

with ; A posn is
      ; (make-posn num num)

; A z-grade is either
; - num
; - (make-posn num num)
; - empty
```

No, in the sense that they generate different templates

## Data Definitions and Templates

The template depends on the *static, textual* content of a data definition, only

```
; A w-grade is either      (define (func-for-w-grade w)
; - num                    (cond
; - posn                   [(number? w) ...]
; - empty                  [(posn? w) ... (func-for-posn w) ...]
; A posn is                [(empty? w) ...])
; (make-posn num num)      (define (func-for-posn p)
;                           ... (posn-x p) ... (posn-y p) ...)

; A z-grade is either      (define (func-for-z-grade z)
; - num                    (cond
; - (make-posn num num)    [(number? z) ...]
; - empty                  [(posn? z) ... (posn-x z) ... (posn-y z) ...]
;                           [(empty? z) ...])
```

## Data Definitions and Templates

Why we treat the data definition statically to generate a template:

- Provides well-defined, simple rules for generating a template
  - "Dynamic" coverage is difficult in general
  - Recall 3520 anecdote: thinking in terms of dynamic coverage  $\Rightarrow$  broken programs
- Similar to the way that data choices affect modularity
  - Details of modularity are beyond the scope of this class, but we want to build the right instincts

## Outline

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## Execution in DrScheme

Suppose that DrScheme's definition window contains

```
(define (f x)
  (/ x 2))
(f 10)
```

What's the result of clicking **Execute** ?

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## Execution in DrScheme

Suppose that DrScheme's definition window contains

```
(define (f x)
  (/ x 0))
(f 10)
```

What's the result of clicking **Execute** ?

*/: divide by 0*

## Execution in DrScheme

Suppose that DrScheme's definition window contains

```
(define (f x)
  (/ x 0))
```

What's the result of clicking **Execute** ?

Nothing (although `f` would produce an error if it were used)

## Execution in DrScheme

Suppose that DrScheme's definition window contains

```
(define (f x)
  (/ x (0)))
```

What's the result of clicking **Execute** ?

*expected a name after an open parenthesis,  
found a number* — even without using `f`

## Execution in DrScheme

Suppose that DrScheme's definition window contains

```
(define (f x)
  (cond x))
```

What's the result of clicking **Execute** ?

*cond: expected a question--answer clause* — even without using `f`

## Execution in DrScheme

Suppose that DrScheme's definition window contains

```
(define (f x)
  (cond
    [false x]))
```

What's the result of clicking **Execute** ?

Nothing

## Execution in DrScheme

Suppose that DrScheme's definition window contains

```
(define (f x)
  (cond
    [false x]))
(f 10)
```

What's the result of clicking **Execute** ?

*cond: all questions were false*

## Errors in DrScheme

DrScheme complains about a function body

- sometimes before the function is used
- sometimes only when the function is called

Why?

Because some errors are ***syntax errors*** and some errors are ***run-time errors***

## Syntax Errors

A ***syntax error*** is like a question that isn't a well-formed sentence

- `f (x) = x + 0`
  - DrScheme doesn't understand this notation, just like...
- "Parlez vous Francais?"
  - English-only speaker doesn't understand this notation
- `(define (f x) (/ x (0)))`
  - Prens around a zero make no sense to DrScheme, just like...
- "Does rain dog cat?"
  - Not enough verbs for this to make sense in English

When DrScheme sees a syntax error, it refuses to evaluate

## Run-Time Errors

A *run-time error* is like a well-formed question with no answer

- `(/ 12 0)`
  - A clear request to DrScheme, but no answer, just like...
- "Why are you wearing a green hat?"
  - There's no answer if I'm wearing a blue hat
- `(cond [false 10])`
  - There's no reasonable choice for DrScheme, just like...
- "If you can't understand me, what's your name?"
  - No one who understands the question should answer

DrScheme evaluates around run-time errors until forced to answer

## The Difference between DrScheme and English

In a (good) programming language, all errors are well-defined, and the rules are relatively simple

- DrScheme has a simple, well-defined grammar, and deviations from the grammar are syntax errors
- The reduction rules for each construct and primitive operation are well-defined, producing either a value or an error

## Beginner Scheme Grammar

A `<var>` is a name, a `<con>` is a constant, and a `<prm>` is an operator name

A `<defn>` is one of

```
(define (<var> <var> ... <var>) <exp>)  
(define <var> <exp>)  
(define-struct <var> (<var> ... <var>))
```

A `<exp>` is one of

```
<var>  
<con>  
(<prm> <exp> ... <exp>)  
(<var> <exp> ... <exp>)  
(cond [<exp> <exp>] ... [<exp> <exp>])  
(cond [<exp> <exp>] ... [else <exp>])  
(and <exp> ... <exp>)  
(or <exp> ... <exp>)
```

## Evaluation Rules: and/or

```
(and true) → true  
(and true question ... question)  
           → (and question ... question)  
(and false question ... question) → false  
  
(or false) → false  
(or false question ... question)  
           → (or question ... question)  
(or true question ... question) → true
```

Note that

```
(and 7 false)
```

fits the grammar, but has no matching evaluation rule, so it produces a run-time error

## Outline

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## Execution in DrScheme

Suppose that DrScheme's definition window contains

```
; f : num -> num
(define (f x)
  (+ x 2))
(f 'apple)
```

What's the result of clicking **Execute** ?

*+: expects a <number>, given 'apple*

But this is really a contract violation at the call to **f**

The implementor of **f** might want to clarify that this error is someone else's fault, not a bug in **f**

## Defensive Programming

```
; f : num -> num
(define (real-f x)
  (+ x 2))
(define (f x)
  (cond
    [(number? x) (real-f x)]
    [else (error 'f "not a number")]))
(f 'apple)
```

*f: not a number*

The **error** function triggers a run-time error

You don't have to program defensively in this course, but it sometimes helps to defend against your own mistakes!