

#### iBTune: Individualized Buffer Tuning for Large-scale Cloud Databases

Jian Tan, Tieying Zhang, Feifei Li, Jie Chen, Qixing Zheng, Ping Zhang, Honglin Qiao, Yue Shi, Wei Cao, Rui Zhang Alibaba

VLDB 2019



# Outline

**Background and Motivation** 

#### **Algorithm and System**

#### **Deployment and Result**

## Background





#### Table 1: Usage of different memory pools

Memory	7	buffer	$\overline{\ }$	insert	log	join	key	read	sort
Pool	[	$\operatorname{pool}$		buffer	buffer	buffer	buffer	buffer	buffer
Avg. Size	2	9609.98N	1	8.00M	$200.00 \mathrm{M}$	0.13M	8.00M	0.13M	$1.25 \mathrm{M}$
Percent		99.27%	/	0.03%	0.67%	0.00%	0.03%	0.00%	0.00%

The memory uses at Alibaba product environment

#### Buffer pool is the largest memory consumer

### Motivation



Reduce memory (buffer pool) while guaranteeing SLA (response time).

- DBA manually uses a small number of BP sizes (10 configurations in our case).

- Each instance's BP size might be different as the query workload is different.

- Manual tuning is not scalable for large cloud databases since each instance has different BP size.



#### **iBTune:** Individualized Buffer Tuning for Largescale Cloud Databases

# iBTune - Preliminary Attempt



Buffer pool (BP) size is sensitive to miss ratio: BP size is reduced from 188G to 80G when it's hit ratio is from 99.968% to 99.950%



#### Intuition:



- Challenge: Heuristic method (such as shrinking 10% each time) does not work, since we have to try many times, which makes the system unstable and is unacceptable for mission-critical applications.
  - Calculate BP based on hit ratio (miss ratio) to avoid restarting system multiple times
  - Confirm whether the BP size meets the requirement of SLA



# Outline

#### **Background and Motivation**

**Algorithm and System** 

**Deployment and Result** 



#### Finding BP=f(miss ratio)



 A number of empirical measurements on real systems have shown power law popularity distributions and follow that:

 $\frac{\log\left(mr_{target}\right) - \log\left(mr_{cur}\right)}{\log\left(bp_{bp_{target}}\right) - \log\left(bp_{cur}\right)} \approx -\alpha_i.$ 

 Parameter α is obtained from the workload, which is 1.2 in our case. A large class of heavy-tailed requests with popularities following a power law distribution fits in our formulation.



#### Calculating tolerable miss ratio

- K-nearest-neighbors (DB instances)
- Find the nearest neighbors
  - Such as 6 in our case
  - Neighbor distance is calculated by similarity
- Calculate tolerable miss ratio
  - The weighted mean of the miss ratios of the k-nearest-neighbors

How to obtain k-nearest-neighbors?

 $N(i) \leftarrow \text{connected nodes of } i \text{ on } G;$ for each instance j in N(i) do  $\begin{vmatrix} \text{ if } \underline{mr_{(cur,j)} > mr_{(cur,i)}} \text{ then} \\ | w_{ij} = \exp((E_j - E_i)^2/(2\sigma^2)); \\ \text{else} \\ | w_{ij} = 0; \\ \text{end} \\ \end{vmatrix}$ end  $mr_{(tolerable,i)} \leftarrow \frac{\sum_{j \in N(i)} w_{ij}mr_{(cur,j)}}{\sum_{i \in N(i)} w_{ii}};$ 



達摩院

## Calculating similarity

- Features
  - RT
  - QPS
  - miss ratio
  - CPU usage
  - logical read
  - io read

#### The last three metrics are divided by QPS



Pearson correlation coefficients



Till now, we can get the k-nearest-neighbors and tolerable miss ratio and calculate the new BP size

## Predicting RT



Pairwise DNN: Predict the respond time (RT) based on the tolerable miss ratio

• Training



The predicted RT is compared with the safe SLA (RT)

## Predicting RT



Pairwise DNN: Predict the respond time (RT) based on the tolerable miss ratio

• Training

•



• The granularity of metric is a day

The predicted RT is compared with the safe SLA (RT)

## Safe SLA (RT)





#### **Response Time Distribution**

Determine the safe RT for different applications

Group all the instances into different applications, and find the 95% percentile of the response times in each group as the corresponding safe limit for that application.



# Outline

#### **Background and Motivation**

### **Algorithm and System**

**Deployment and Result** 

#### System halt avoidance



#### Based on X-Paxos: high availability protocol implementation at Alibaba



#### **Evaluation**



- All results are from our product environment
- X-Engine: MySQL compatible database based on LSM-Tree storage engine
- With high performance Paxos implementation
- Pairwise DNN: 100K data samples

	Taobao	Tmall	Youku	Fliggy	Others
Select	5245/s	2815/s	1017/s	22930/s	120/s
Insert	4222/s	0	1/s	1520/s	10/s
Update	0	30/s	315/s	4/s	980/s
Delete	2708/s	0	10/s	0	0



1,000 sample instances scattered across different applications

iBTune has been deployed on 10,000 database instances Memory saving : ~17%

## Single instance

Performance before and after BP adjustment (45G->21G) during holidays and workdays:

- Red line is the time when BP size is adjusted
- Green lines show the holiday which is 7-days
- Predicted RT: only 3 points exceeded which is acceptable

The IO read metric is the real IO, since all our DB instances turn on direct IO



#### Multiple instances



10 representative instances. The memory saving ranges from 50% to 10%, which strongly supports that a single number does not fit all. Instance 1 has a large increase in RT after the adjustment. We find that there is one query that consumes 99.97% of the total response time. The lookup value in WHERE condition changes for this query.

#### **Conclusion & Future Work**



- iBTune has been deployed on 10,000 database instances with memory saving : ~17%
- Future work
  - Cache preload
    - Backup node needs run SQLs to load data into cache after BP adjustment
    - Perform switching after preload
  - Buffer increase
    - Currently reply on DBA
    - Automatic increase buffer
  - Multiple parameters tuning
    - DBMS configure file



# Thanks