

L21 -- Markov Chain
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Graph $G = (E, V)$

$V =$ vertices $\{a, b, c, d, e, f, g, h\}$

$E =$ edges $\{(a, b), (a, c), (a, d), (b, d), (c, d), (c, e), (e, f), (e, g), (f, g), (f, h)\}$
unordered pairs

Draw graph:

	a	b	c	d	e	f	g	h
a	0	1	1	1	0	0	0	0
b	1	0	0	1	0	0	0	0
c	1	0	0	1	1	0	0	0
d	1	1	1	0	0	0	0	0
e	0	0	1	0	0	1	1	0
f	0	0	0	0	1	0	1	1
g	0	0	0	0	1	1	0	0
h	0	0	0	0	0	1	0	0

****adjacency matrix****

Each v in V is a state.

If at b , represent state as

$q = [0 \ 1 \ 0 \ 0 \ 0 \ 0 \ 0 \ 0 \ 0]^T$

Can "think of" fractional state

$q = [1/2 \ 0 \ 0 \ 1/2 \ 0 \ 0 \ 0 \ 0 \ 0]^T$

1/2 at a and 1/2 at d

probability of being in each state:

each $q[i] \geq 0$ and $\sum_i q[i] = 1$

Transition matrix $P =$ normalized adjacency matrix

	a	b	c	d	e	f	g	h
a	0	1/2	1/3	1/3	0	0	0	0
b	1/3	0	0	1/3	0	0	0	0
c	1/3	0	0	1/3	1/3	0	0	0
d	1/3	1/2	1/3	0	0	0	0	0
e	0	0	1/3	0	0	1/3	1/2	0
f	0	0	0	0	1/3	0	1/2	1
g	0	0	0	0	1/3	1/3	0	0
h	0	0	0	0	0	1/3	0	0

then given a state q , we can "transition" to the next state by

$$q_1 = P * q$$

This one "step" of a "Markov Chain".

"Markov" means that each state only depends on previous state.

next step

$$q_2 = P * q_1 \quad \text{or}$$

$$= P * P * q \quad \text{or}$$

$$= P^2 * q$$

$$q_n = P^n * q$$

where $P^n = P * P * P * \dots n \text{ times} \dots * P$

Can think of as a randomized random walk.

+ start state $q = q_0$.

+ each step, takes one path at random

+ q_n is probability distribution of state after n steps

+ thus each column of P^n positive, sums to 1 for all n

Markov Chain is **ergodic** if

exists some t such that for all $n \geq t$ then

each entry in P^n is positive.

--> for any q , then

$$q_n = P^n q$$

is positive in all elements

--> after t steps, always have *some* probability of being anywhere.

When is a chain not ergodic?

+ cyclic

$$P = \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix}$$

$$\begin{bmatrix} 1 & 0 \end{bmatrix}$$

always alternates states in even/odd states

--> can be larger and more irregular, uncommon in practice

+ has absorbing + transient states

P based on *directed* graph

$$P = \begin{bmatrix} 0 & 1/2 & 1/2 & 0 \\ 1/2 & 0 & 1/2 & 1 \\ 1/2 & 1/2 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{bmatrix}$$

$$\begin{bmatrix} 1/2 & 0 & 1/2 & 1 \end{bmatrix}$$

$$\begin{bmatrix} 1/2 & 1/2 & 0 & 0 \end{bmatrix}$$

$$\begin{bmatrix} 0 & 0 & 0 & 0 \end{bmatrix}$$

state d always goes to b , but can never return to d .

also...

$$P = \begin{bmatrix} 0 & 1/2 & 1/2 & 0 \\ 1/2 & 0 & 1/2 & 1/2 \\ 1/2 & 1/2 & 0 & 0 \\ 0 & 0 & 0 & 1/2 \end{bmatrix}$$

may stay at d (w.p. 1/2) but state "seeps" from d to b (and then a,c)

(a,b,c) = absorbing, d = transient

+ not connected

$$P = \begin{bmatrix} 1/2 & 1/2 & 0 & 0 \\ 1/2 & 1/2 & 0 & 0 \\ 0 & 0 & 2/3 & 1/2 \\ 0 & 0 & 1/3 & 1/2 \end{bmatrix}$$

(a,b) cannot reach (c,d) and vice-versa

Consider an ergodic Markov Chain (P,q)

****AMAZING**** property

let $P^{*} = P^{n}$ as $n \rightarrow \infty$

then $q_{*} = P^{*} q$

is ****NOT**** dependent on q

--> That is, for all starting states q, the final state is q_{*}

--> as we do a random walk, we will eventually be in the same expected state.

Note that $q_{*} = P^{*} q = P^{*+1} q$

so $q_{*} = P q_{*}$

--> If state distribution is initially q_{*} , then already in final distribution.

q_{*} second eigenvector of P

second eigenvalue determines rate of convergence

--> smaller <-> faster convergence

Metropolis Algorithm (MCMC)

Metropolis, Rosenbluth, Rosenbluth, Teller, Teller 1953

(Boltzman dist, Manhattan project)

Hastings 1970

(more general)

each state v in V has weight associated with it

$$w(v) \quad \sum_{\{v \text{ in } V\}} w(v) = W$$

Want to land in state v w.p. $w(v)/W$

--> V might be very large, and W unknown.

--> V can be "continuous"

"probe-only" can only measure $w(v)$ at any one state

Strategy: design special Markov Chain so $q_*[v] = w(v)/W$

Start v_0 in V ($q = [0 \ 0 \ 0 \ \dots \ 1 \ \dots \ 0 \ 0]^T$)

choose neighbor u (proportional to $K(v,u)$)

if ($w(u) \geq w(v_i)$) --> $v_{\{i+1\}} = u$

else w.p. $w(u)/w(v)$ --> $v_{\{i+1\}} = u$

else --> $v_{\{i+1\}} = v_i$

if ergodic:

there exists some t s.t. for $i \geq t$

$\Pr[v_i = v] = w(v)/W$

NOTE: not in limit, but for some finite t (even for continuous) V
through AMAZING "coupling from past"

But t is hard to find.

Often goal is to create many samples:

formal: run for t_+ steps, take sample, ...

run for another t_+ steps, take sample, ... repeat

in practice: run for 1000 steps (burn in),
take next 5000 steps as random samples

has "auto-correlation" but eventually more time efficient than tN steps for N
samples

and t unknown.

"inherently sequential" makes very hard to parallelize

Applies even if V is continuous