# Big Data Systems

## Today ...

- General course overview
- Expectations from this course
- ► Q&A
- Introduction to Big Data
- Assignment #1

## General Course Information

#### Course Web Page

- http://www.cs.utah.edu/~hari/teaching/fall2015.html
- also on canvas
- Text & References
  - Mining of Massive Datasets, Leskovec, Rajaraman, Ullman
  - will use online resources and papers
- ► CHPC accounts → fill this form asap <u>http://goo.gl/forms/nT8TkFtiJq</u>
- ► TA Vairavan Sivaraman
  - Email <u>vairavan.sivaraman@utah.edu</u>
  - Office hours

## Class Interaction ...

I strongly encourage discussions & interactions

10% for class interaction and online participation

## Academic conduct

Adherence to the CoE and SoC academic guidelines is expected. Please read the following.

- College of Engineering Academic Guidelines from their web page
- School of Computing Misconduct Policy Please Read

<u>TL;DR</u> NO CHEATING

## Exams, Assignments & Projects

- Assignment 0 Due in 1 week
- Assignments 1-2
  - one every week
  - Simple spark problems
- Assignments 3,4
  - Larger problems
  - 2 weeks each
- Final Project  $\rightarrow$  submit proposal before start of fall break
  - Default project
- 2-3 students per group
- Mid-term exam after fall break

## Things you should know

- ► No formal prerequisite
- Good Programming Skills (any language)
  - Python/Scala/Java
  - Make a case if you would like to use any other language
- Be prepared to learn new programming tools & techniques
- Linear Algebra
- Sequential algorithms, complexity analysis

## What will we cover ?

- Hadoop & Spark
- MapReduce
- Parallel Algorithms
- Searching & Sorting in Parallel
- Randomized Algorithms Graph Coloring
- PageRank
- Clustering Data
- Recommendation Systems
- Matrix Factorization
- Social Networks
- Graph Algorithms
- Large-Scale Machine Learning

## Expectations

- The focus will be on understanding Big data algorithms and ensuring scalability for large scale problems
- While we will use Spark/Python, the skills acquired should be transferable to other platforms/frameworks
- While I will show you code examples, this is not a course to take if you want to learn to program in Java/Python/Scala/....
- Programming assignments will be frequent and hard. Additionally, the cluster is unreliable, so do not wait until the last evening to test your code.
- Choose the default project if you

## Where to look for help

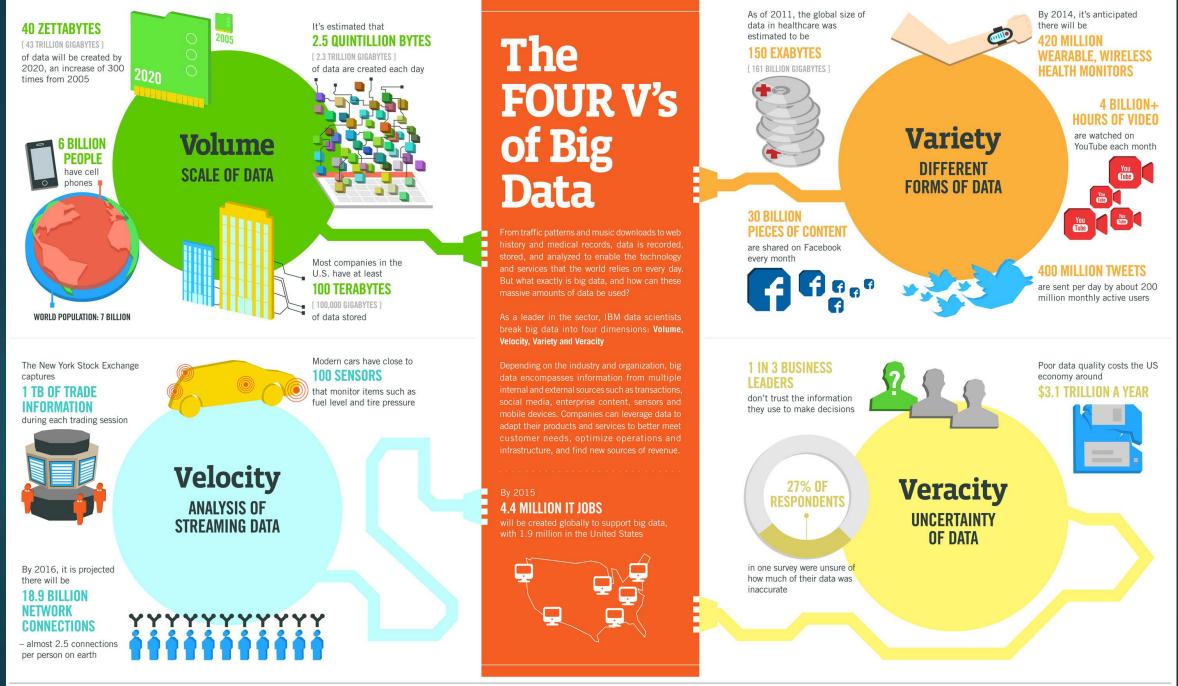
- Course website / Canvas
- Email
- ► TA office hours  $\rightarrow$  MEB 3XXX
- ► My office hours → MEB 3454 → Tue,Thu 2-3pm



CS 5965/6965 - Big Data Systems - Fall 2015

## What is Big Data ?







The recognition that data is at the center of our digital world and that there are big challenges in collecting, storing, processing, analyzing, and making use of such data.

What is **Big** depends on the application domain

## Kinds of Data

- Web data & web data accesses
- Emails, chats, tweet
- ► Telephone data
- Public databases Gene banks, census records, …
- Private databases medical records, credit card transactions, ...
- Sensor data camera surveillance, wearable sensors, seismograms
- Byproduct of computer systems operations power signal, CPU events, ...

## Data is valuable

Google & facebook make money by mining user data

- for revenue from advertisements
- Financial firms analyze financial records, real-time transactions and current events for profitable trades
- Medical records can be processed for better and potentially cheaper – health care
- Roadways are monitored for traffic analysis and control
- Face detection in airports

## Collection of Big Data

- The amount of data available is increasing exponentially
- But, it is still challenging to collect it
  - Difficult to get access
  - Redundancy
  - Noise

## Processing & Analysis of Big Data

- Large datasets do not fit in memory
- Processing large datasets is time consuming
  - ► Parallel processing is necessary → challenging
- Message Passing Interface (MPI) (1991)
  - Low(er) level C/Fortran API for communication
  - Powerful, hard(er) to code, !fault tolerant
- Mapreduce Google (2004)
  - Originated from web data processing
  - ease of programming, fault tolerant
  - limited semantics
- Spark, GraphLab, Storm, .....

## Storage & I/O

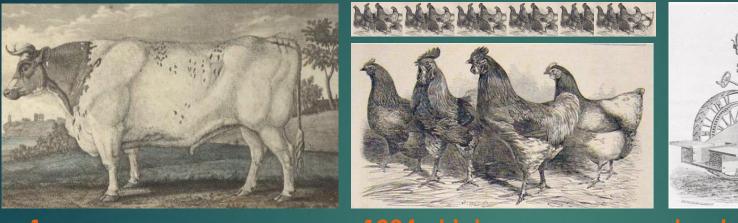
- Storage and I/O are critical for big data performance and reliability
- Hardware: disks, Flash, SSD, nonvolatile memory, 3D memory
- Parallelism: RAID, parallel data storage, DFS
- Data durability and consistency

## Data privacy & protection

- Misuse of big data is a big concern
  - A person's online activities can reveal all aspects of the person's life
- Systems need to provide clear guidelines on data privacy and protection
  - Sensitive clinical information
- Understand how the big data world operates
  - as an user
  - as a developer

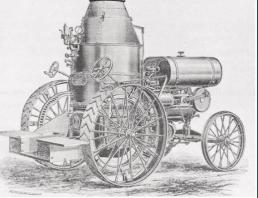
### Parallel Thinking THE MOST IMPORTANT GOAL OF TODAY'S LECTURE

## Parallelism & beyond ...



**1 ox:** single core performance





**tractor:** better algorithms

If you were plowing a field, which would you rather use? Two strong oxen or 1024 chickens? Seymour Cray

Credit: Phillip Stanley-Marbell

Consider an array A with n elements,

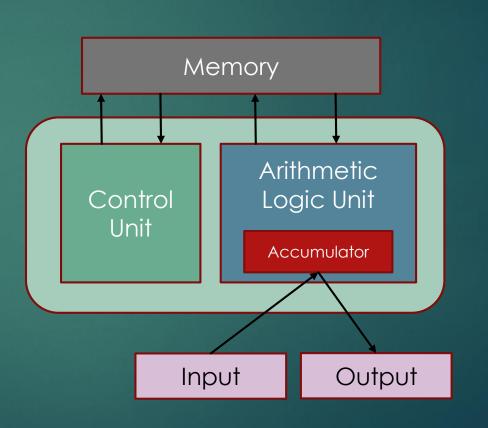
Goal: to compute,

$$x = \sum_{1}^{n} \sqrt{A_i}$$

Machine Model Programming Model Performance analysis

## Von Neumann architecture

- Central Processing Unit (CPU, Core)
- Memory
- Input/Output (I/O)
- One instruction per unit/time
- Sequential



## Characterizing algorithm performance

- ► 0-notation
  - Given an input of size n, let T(n) be the total time, and S(n) the necessary storage
  - Given a problem, is there a way to compute lower bounds on storage and time -> Algorithmic Complexity
  - ► T(n) = O(f(n)) means

 $T(n) \leq cf(n)$ , where c is some unknown positive constant

compare algorithms by comparing f(n).

## Scalability

- $\blacktriangleright Scale Vertically \rightarrow scale-up$ 
  - Add resources to a single node
  - CPU, memory, disks,
- $\blacktriangleright Scale Horizontally \rightarrow scale-out$ 
  - Add more nodes to the system

## Parallel Performance



best sequential time/time on p processors

► Efficiency

speedup/p, (< 1)



## Amdahl's Law

Sequential bottlenecks:

Let s be the percentage of the overall work that is sequential

▶ Then, the speedup is given by

$$S = \frac{1}{s + \frac{1-s}{p}} \le \frac{1}{s}$$

## Gustafson

Sequential part should be independent of the problem size Increase problem size, with increasing number of processors

## Strong & Weak Scalability

Increasing number of cores

**Strong** (fixed-sized) scalability

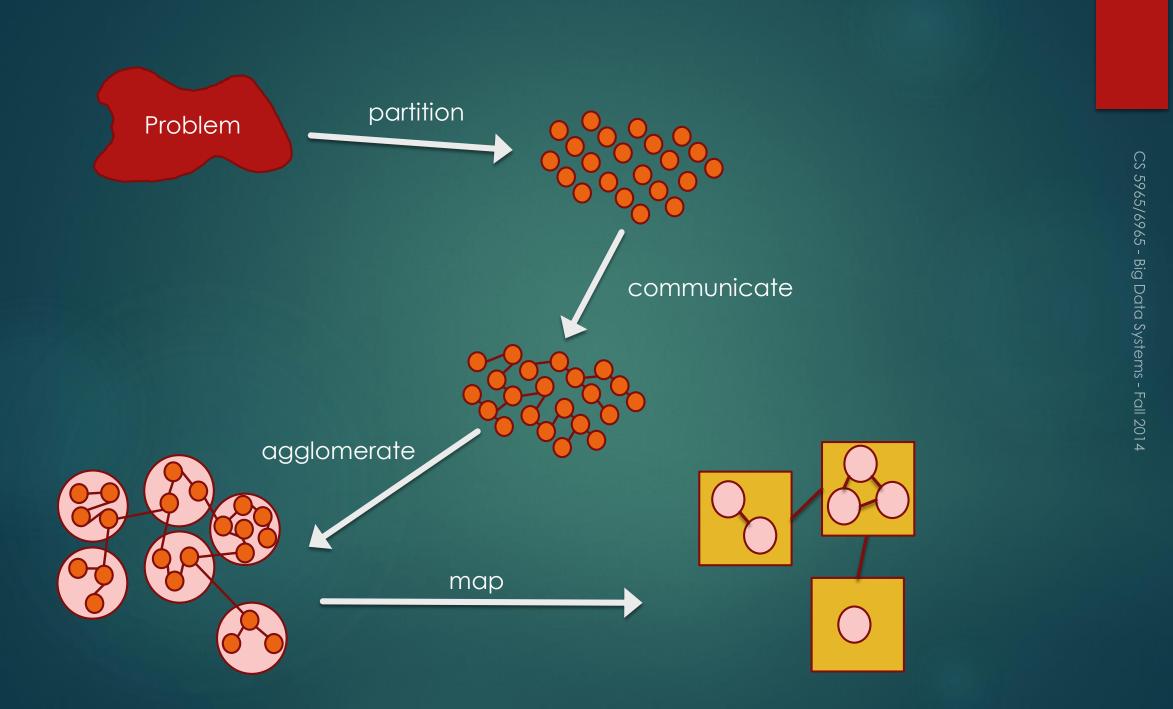
keep problem size fixed

Weak (fixed-sized) scalability

keep problem size/core fixed

## Parallel Programming

- Partition Work Data & Tasks
- Determine Communication
- Agglomeration to number of available processors
- Map to processors
- Tune for architecture



Consider an array A with n elements,

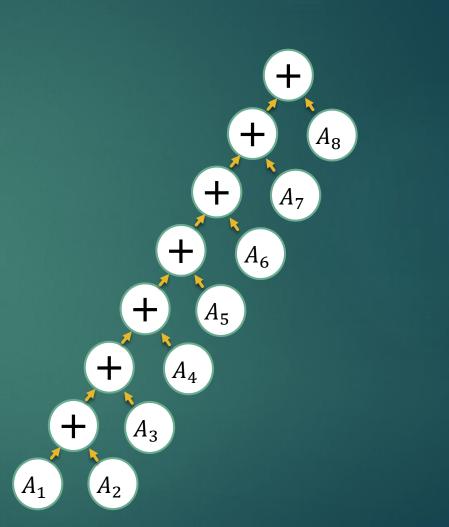
Goal: to compute,

$$x = \sum_{1}^{n} \sqrt{A_i}$$

CS 5965/6965 - Big Data Systems - Fall 2015

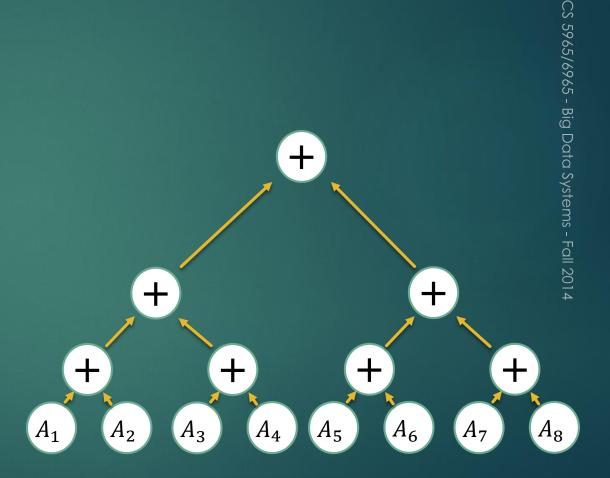
## Work/Depth Models

- Abstract programming model
- Exposes the parallelism
  - Compute work W and depth D
  - D longest chain of dependencies
  - ▶ P = W/D
- Directed Acyclic Graphs
- Concepts
  - parallel for (data decomposition)
  - recursion (divide and conquer)



## Work/Depth Models

- Abstract programming model
- Exposes the parallelism
  - Compute work W and depth D
  - D longest chain of dependencies
  - ▶ P = W/D
- Directed Acyclic Graphs
- Concepts
  - parallel for (data decomposition)
  - recursion (divide and conquer)



## Sequential vs Parallel for

- Dependent statements
  - $\blacktriangleright W = \sum W_i$
  - $\blacktriangleright D = \sum D_i$
- Independent statements
  - $\blacktriangleright W = \sum W_i$
  - $\blacktriangleright D = \max(D_i)$

## Parallel Sum – language model

// Recursive implementation
Algorithm SUM(a, n)
// Input: array a of length  $n = 2^k, k = \log n$ parallel\_for i  $\leftarrow [0, n/2)$  b(i)  $\leftarrow a(2i) + a(2i+1)$ return SUM(b); //  $W\left(\frac{n}{2}\right), D\left(\frac{n}{2}\right)$ 

**Complexity:** 

$$D(n) = D\left(\frac{n}{2}\right) + O(1) = O(\log n)$$
$$W(n) = W\left(\frac{n}{2}\right) + O(n) = O(n)$$

CS 5965/6965 - Big Data Systems - Fall 2015

## Questions?

CS 5965/6965 - Big Data Systems - Fall 2015