

cs5460/6460: Operating Systems

Lecture: System boot

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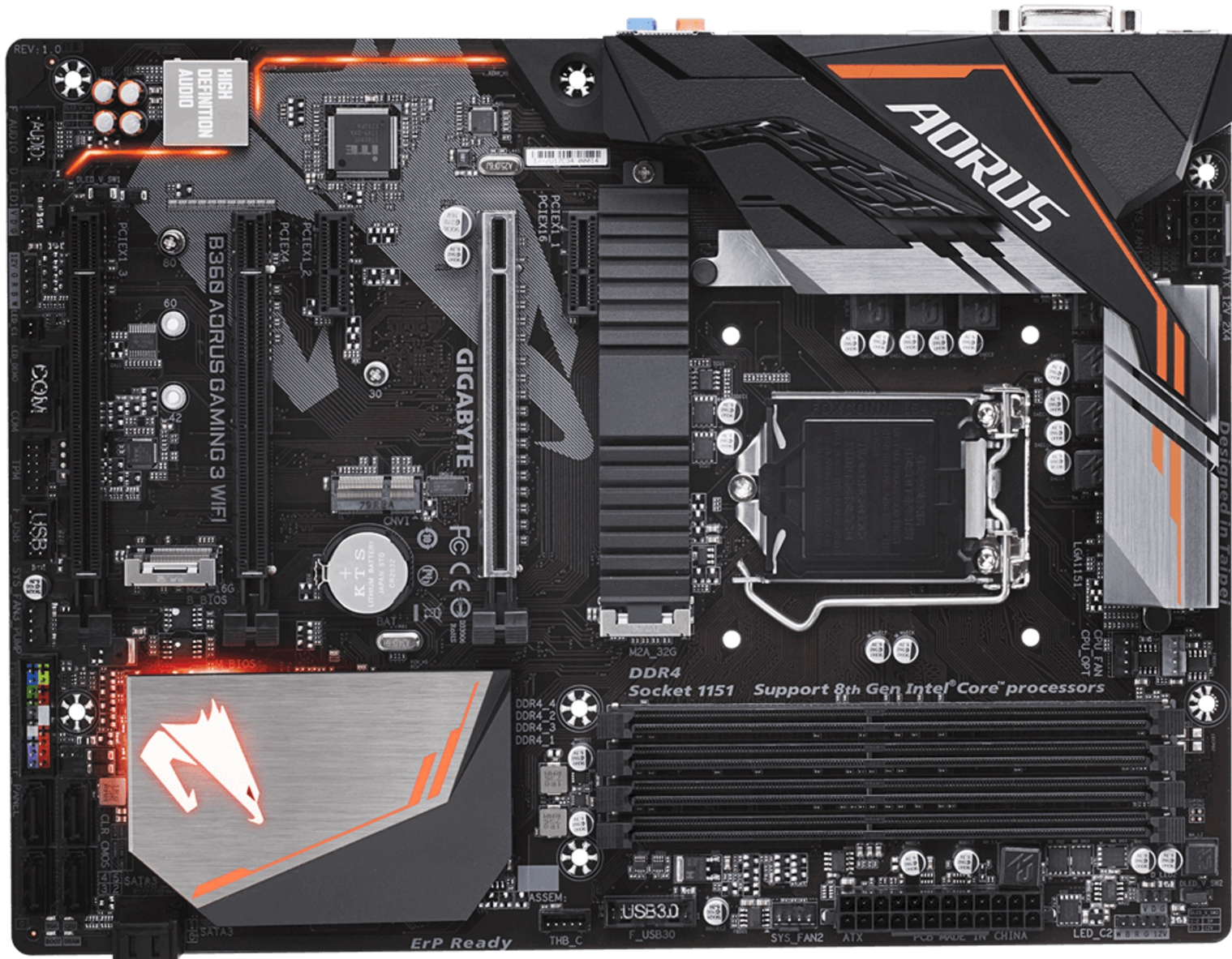
February, 2024

What happens when we turn on the power?

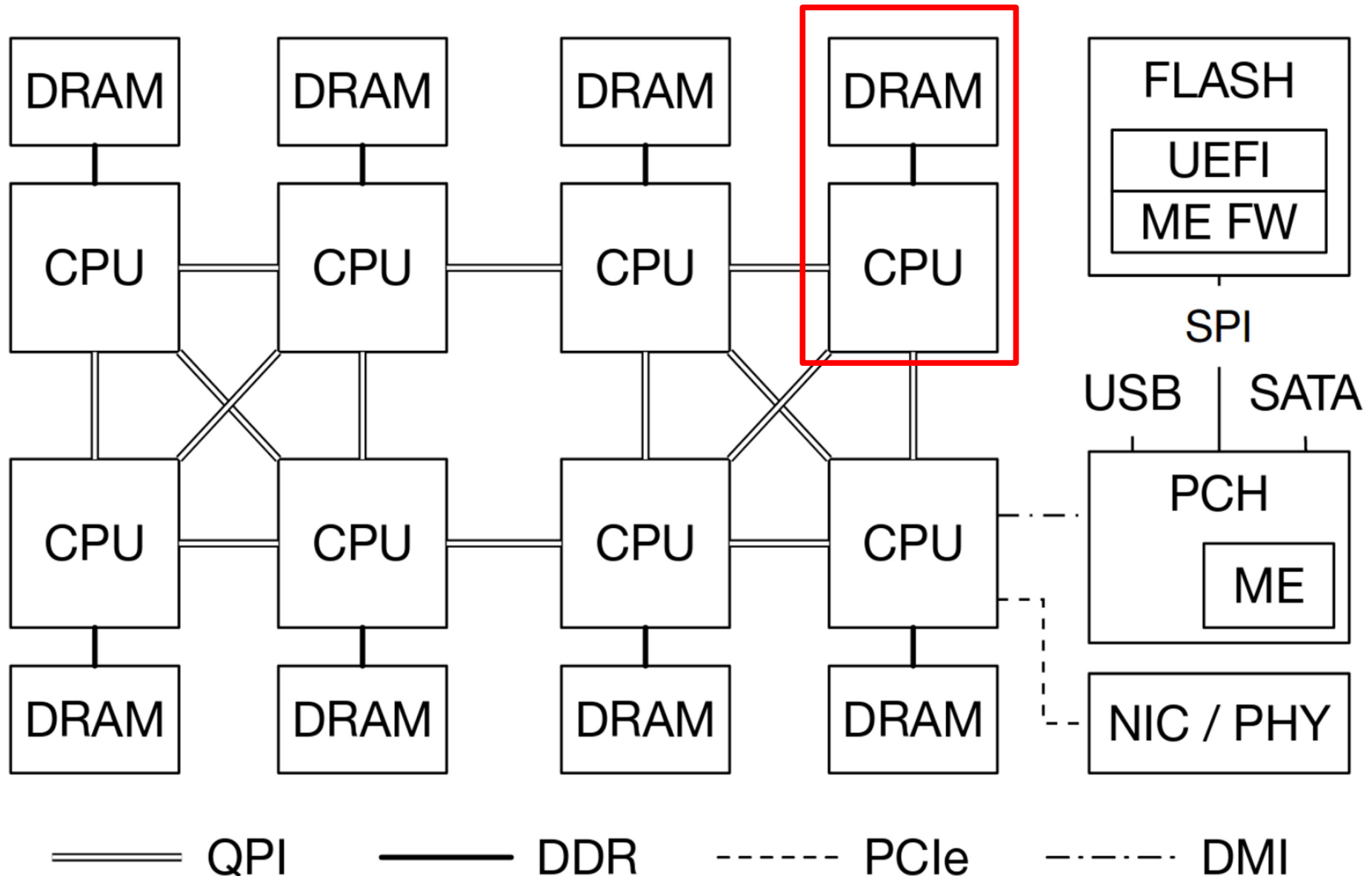
- Well it's complicated
 - Intel SGX Explained is a good start (Section 2.13 [1])
- At a high-level a sequence of software pieces initializes the platform
 - Management engine (ME), microcode, firmware (BIOS), bootloader

- The most important thing: the OS is not the only software running on the machine
 - And not the most privileged
- Today, at least two layers sit underneath the OS/hypervisor
- System Management Mode (SMM) (ring -2)
 - Runs below the hypervisor/OS
- Intel Management Engine and Intel Innovation Engine (ring -3)

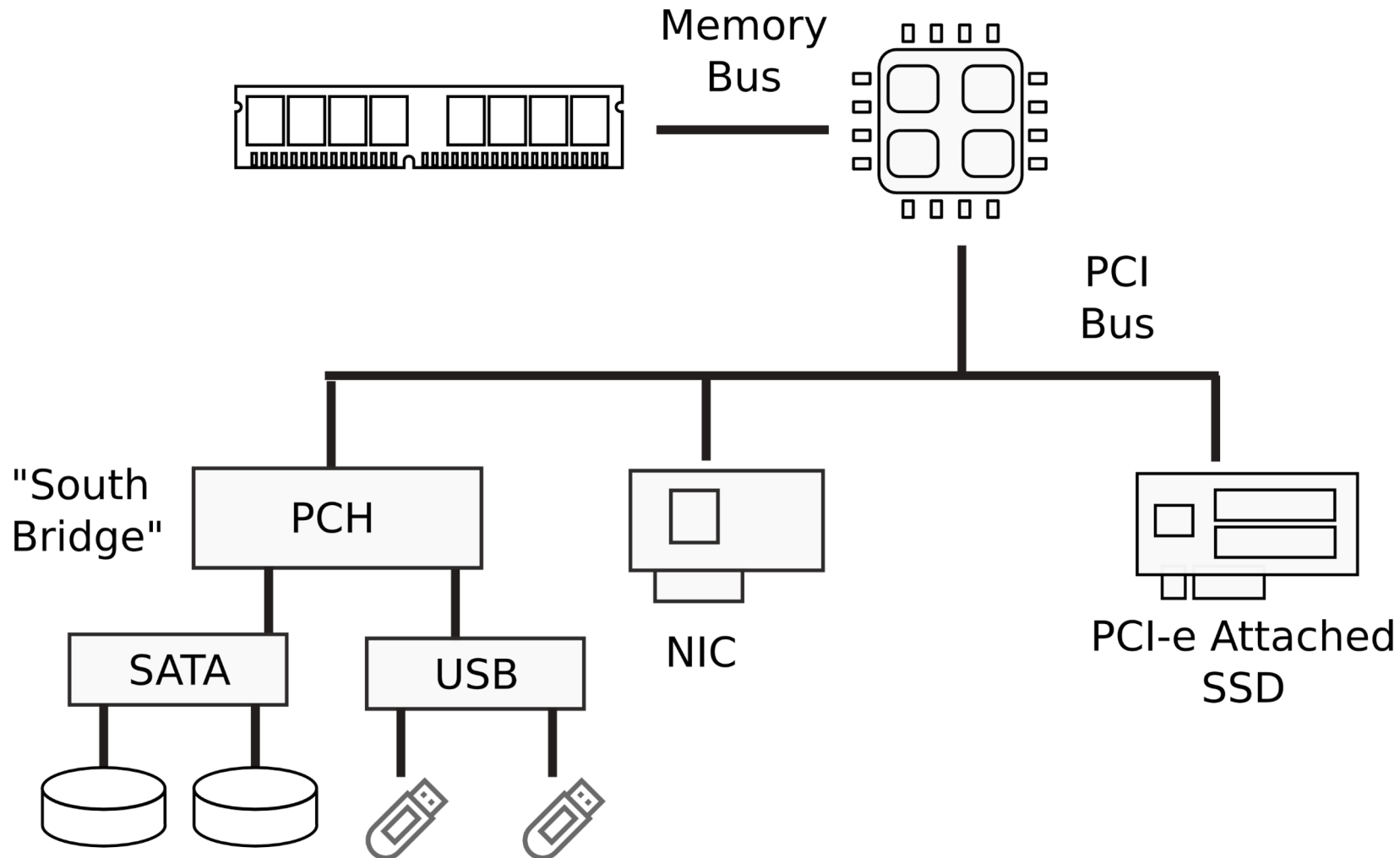
B360 AORUS Motherboard



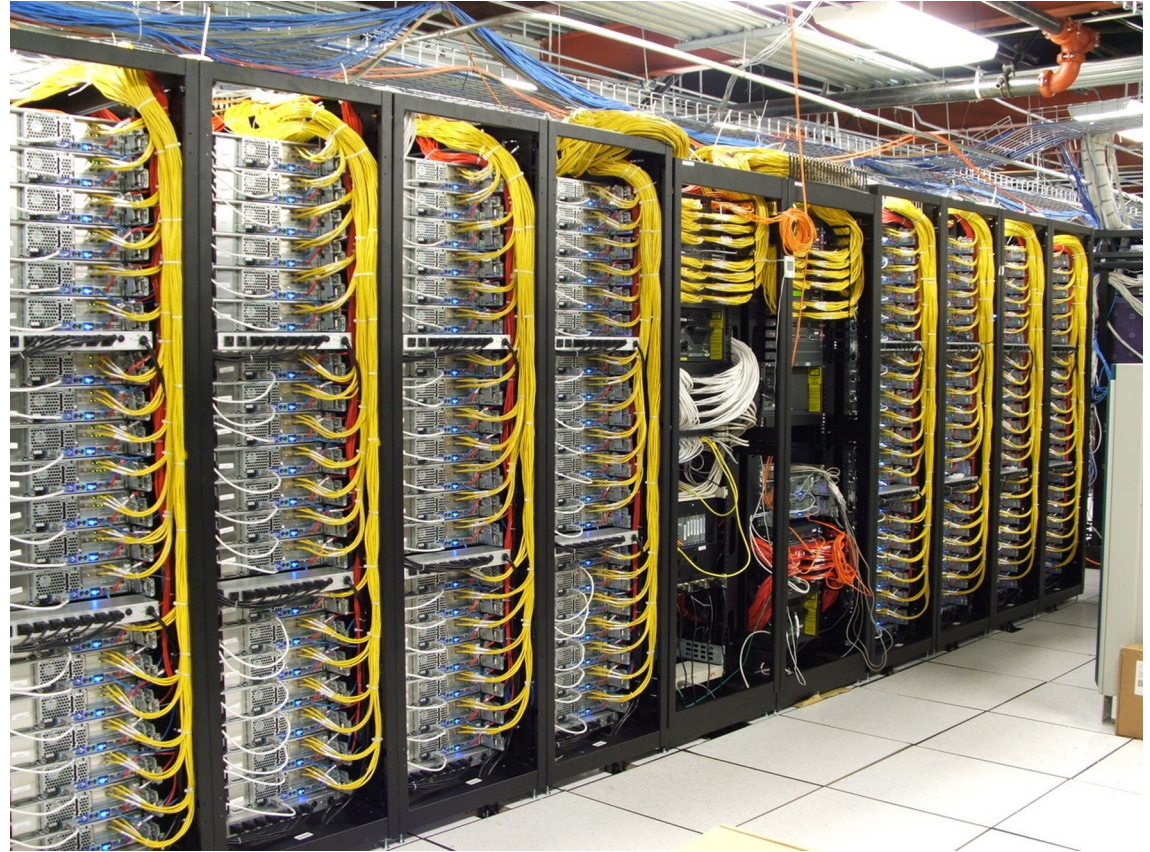
PC motherboard components



I/O Devices



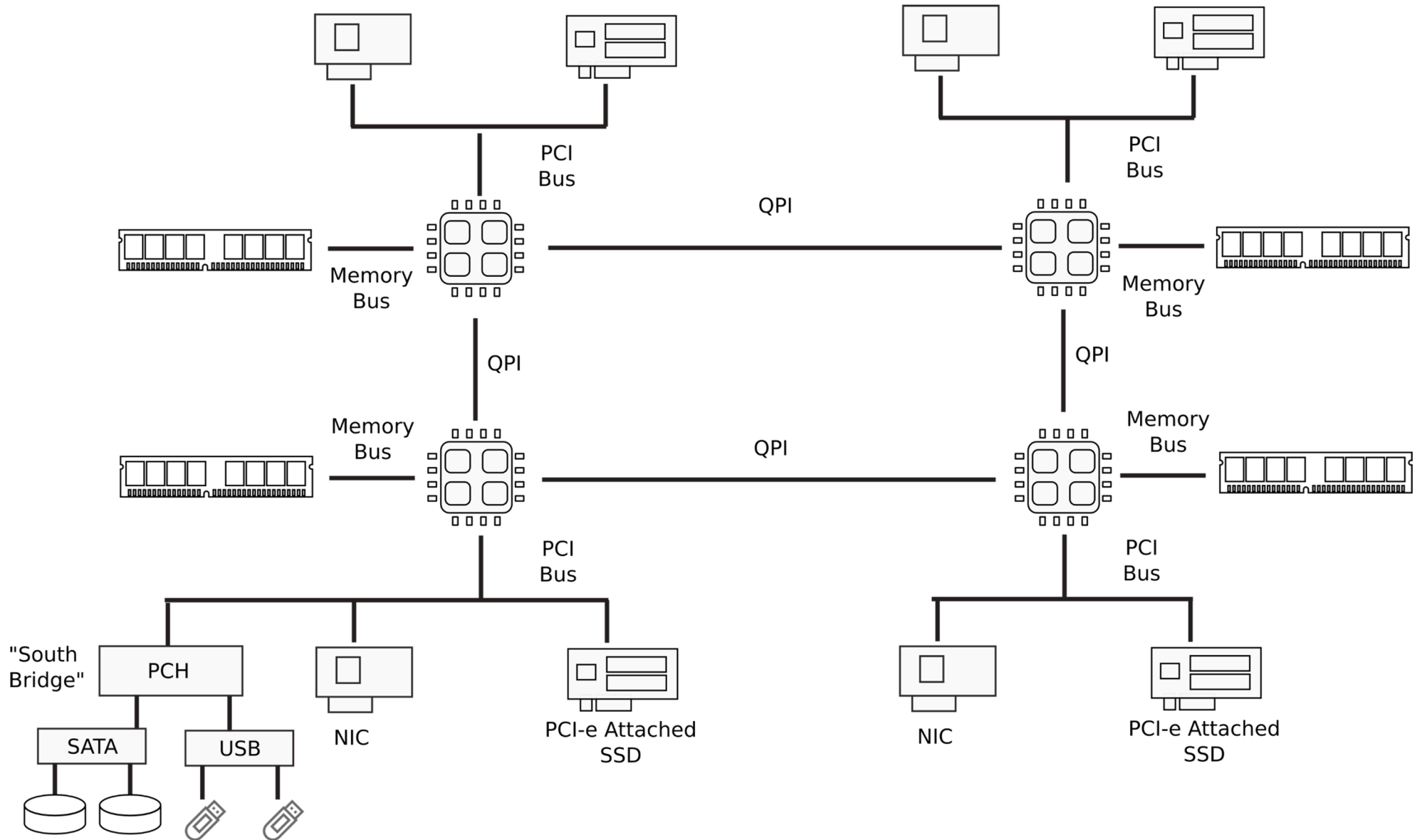
Dell R830 4-socket server



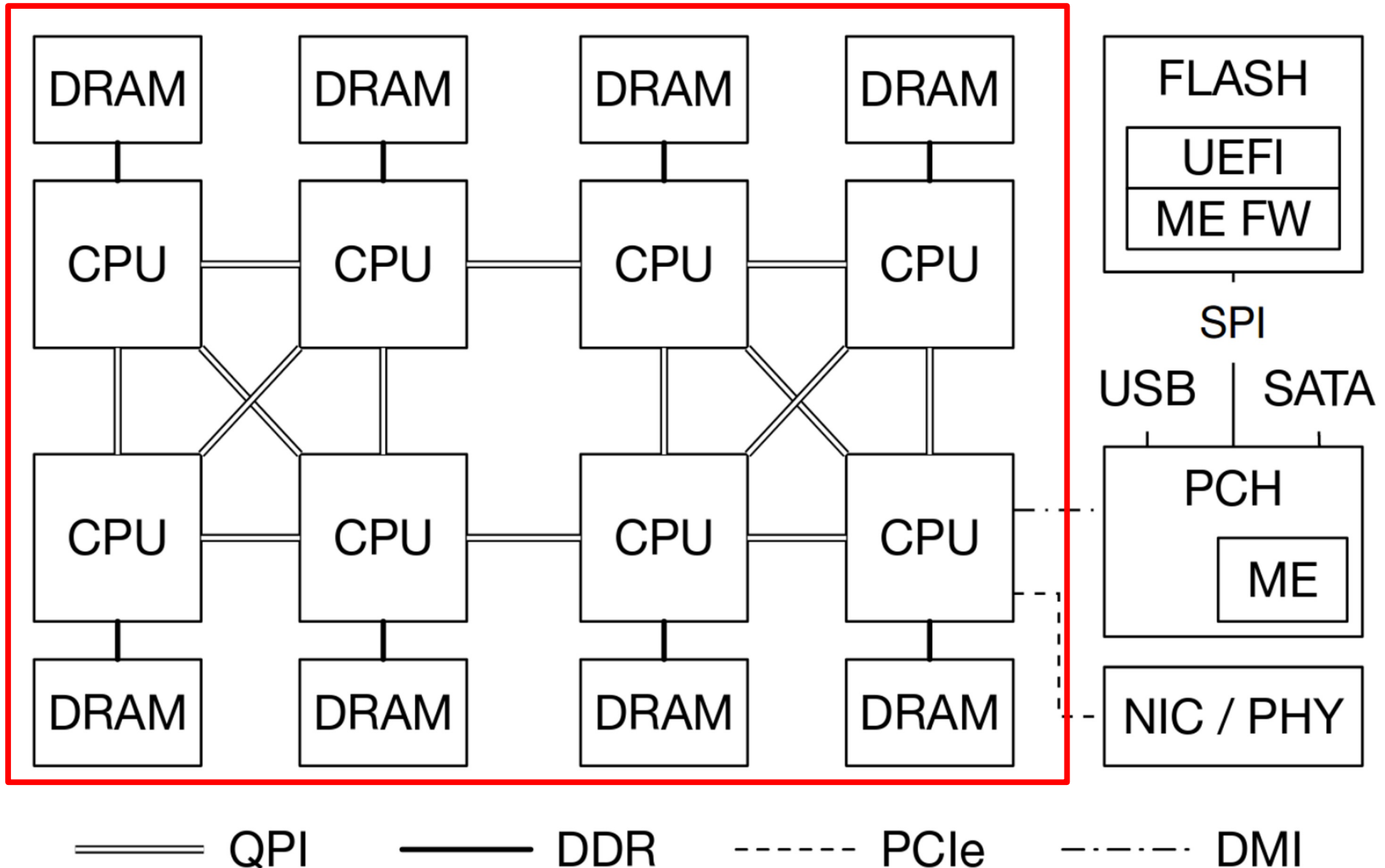
Dell Poweredge R830 System Server with 2 sockets on the main floor and 2 sockets on the expansion

http://www.dell.com/support/manuals/us/en/19/poweredge-r830/r830_om/supported-configurations-for-the-poweredge-r830-system?guid=guid-01303b2b-f884-4435-b4e2-57bec2ce225a&lang=en-us

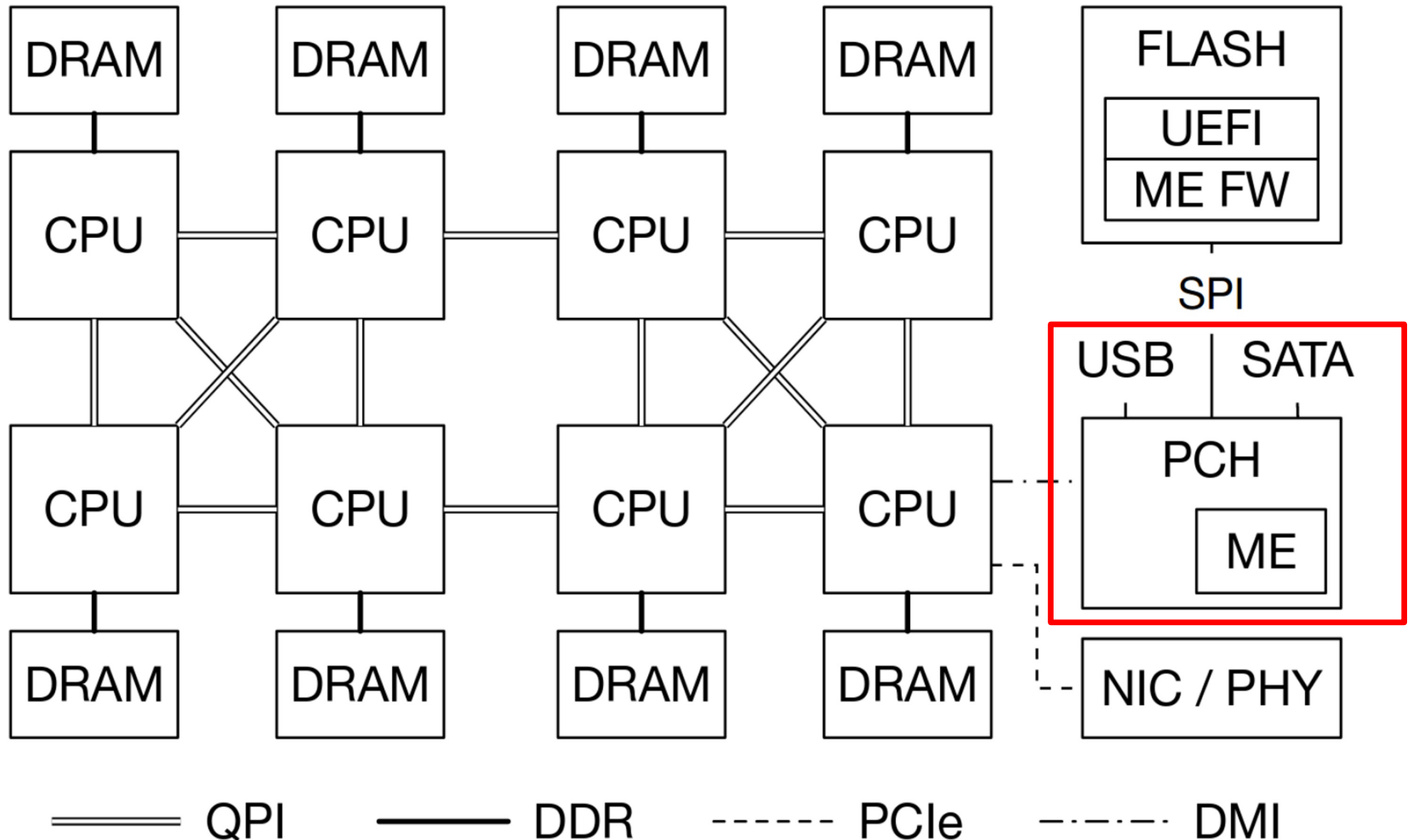
Multi-socket machines



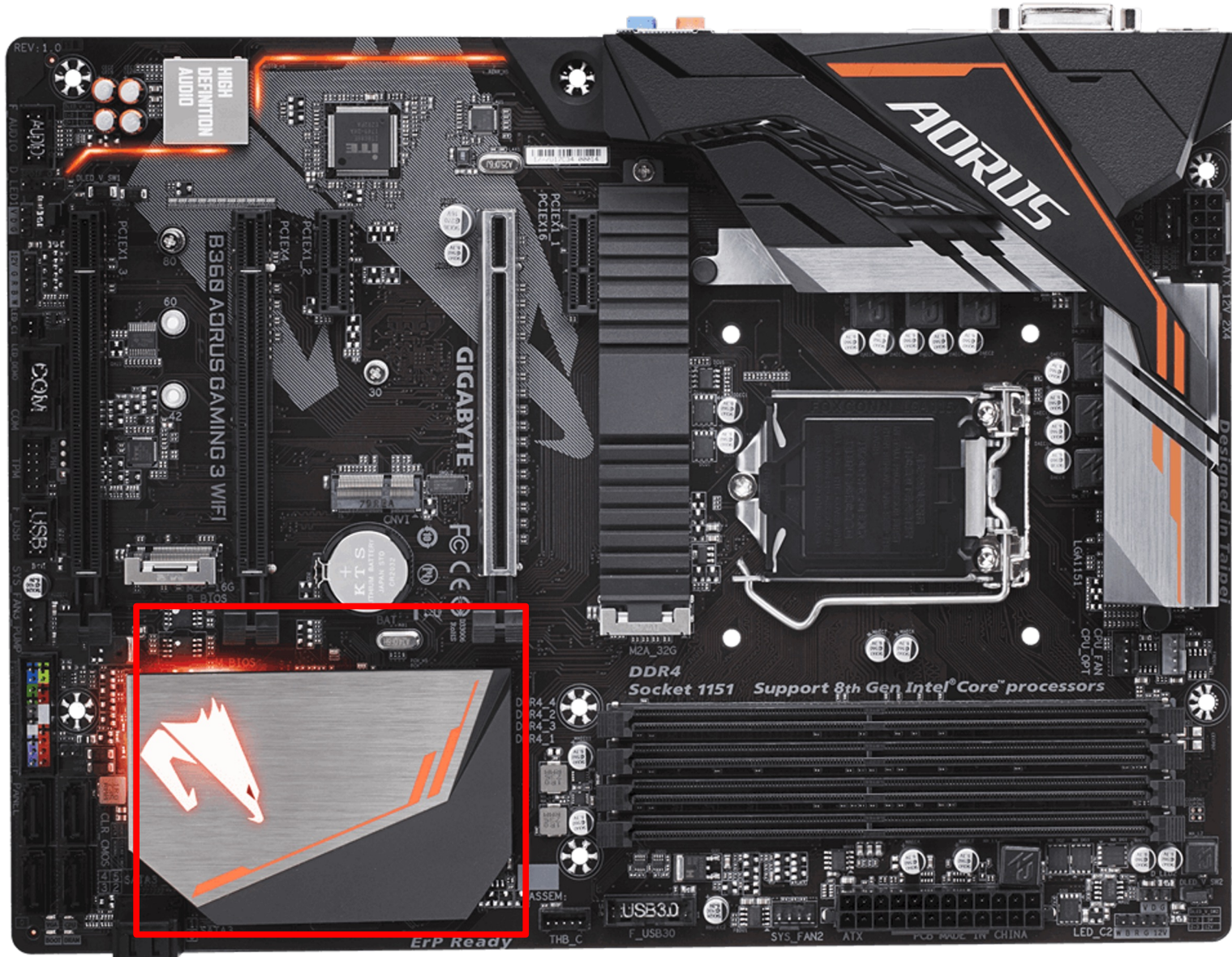
PC motherboard components



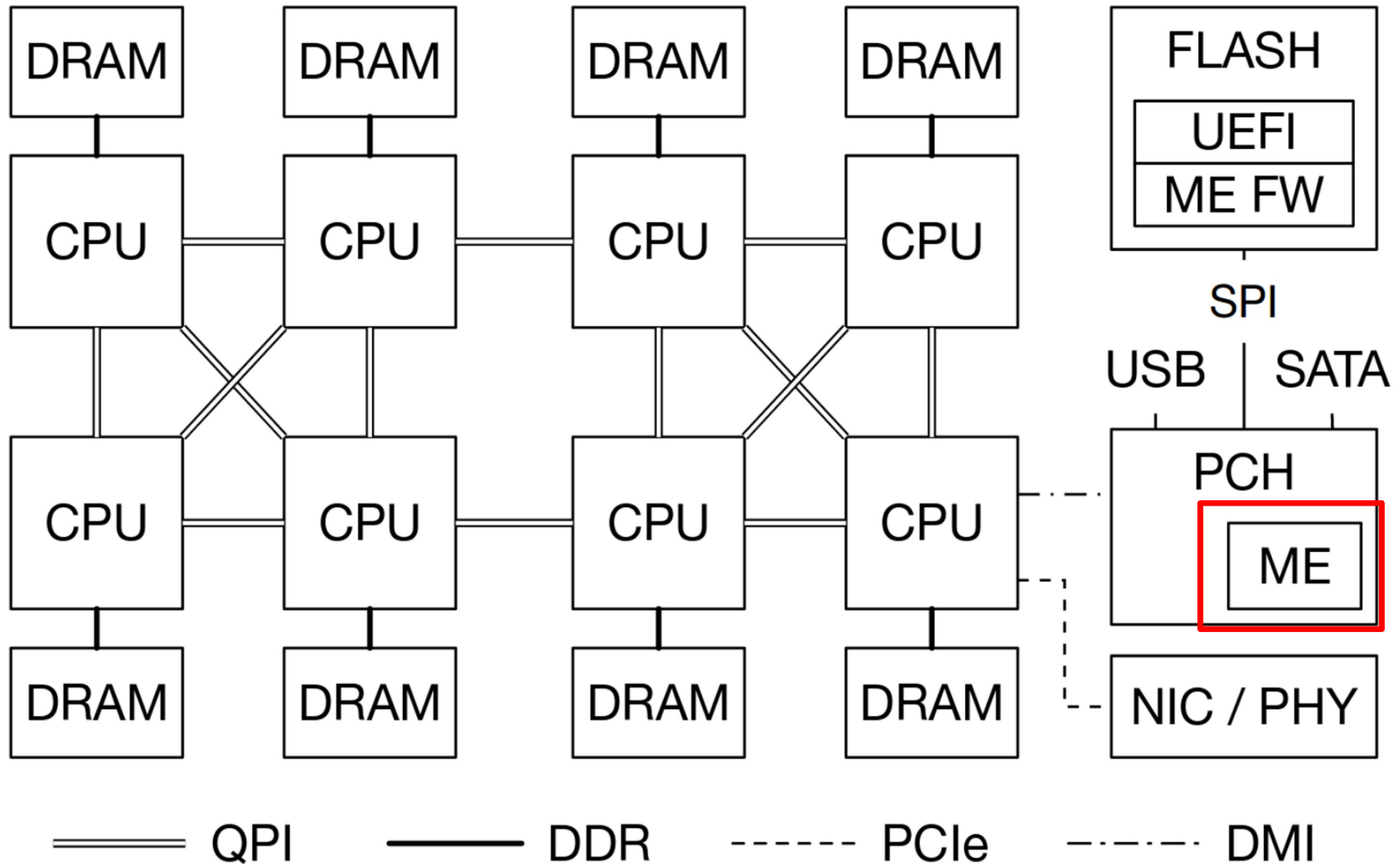
PCH – Platform Controller Hub



B360 AORUS Motherboard



ME gets power before CPUs



Intel Management Engine (ME)

- Full-featured computer
 - Intel Quark x86-based 32-bit CPU
 - Internal RAM (1.7MB)
- Can access all DRAM via DMA
- Can control boot chain
- Can access network interface (NIC) on the motherboard
 - Has it's own MAC and IP address
 - Via System Management Bus (SMBus)
 - Or an ATM compatible NIC
- Connected to the power supply
 - Stays on as long as power is provided to power supply

ME: Theft prevention use-case

- In S5 (computer off) ME cannot access DRAM
 - DRAM is off
- But ME can use its internal memory
 - ME can disable a stolen laptop equipped with cellular modem remotely
- As long as power is connected
- And cell network has signal

Intel Management Engine (ME)

- All modern motherboard chips contain ME
- Part of Active Management Technology (AMT)
- Convenient way for administrators to fix your machine remotely
 - Obviously a huge opportunity for an attack

What's running there?

Do you ever read “Modern Operating Systems”?

POSITIVE TECHNOLOGIES

```
> strings vfs
...
..\..\src\os\servers\vfs\misc.c
FS: bogus child for forking
FS: forking on top of in-use child
...
```

MINIX3 by Andrew Tanenbaum

```
Directory of minix3-master\servers\vfs

14.03.2010  23:52    14'978  main.c
14.03.2010  23:52         741  Makefile
14.03.2010  23:52    17'653  misc.c
14.03.2010  23:52         677  mmap.c
14.03.2010  23:52    15'650  mount.c
...
```

"FS: bogus child for forking"

All

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6 results (0.34 seconds)

[misc.c in minix-filesystem | source code search engine - Searchcode](#)

<https://searchcode.com/codesearch/view/55926734/>

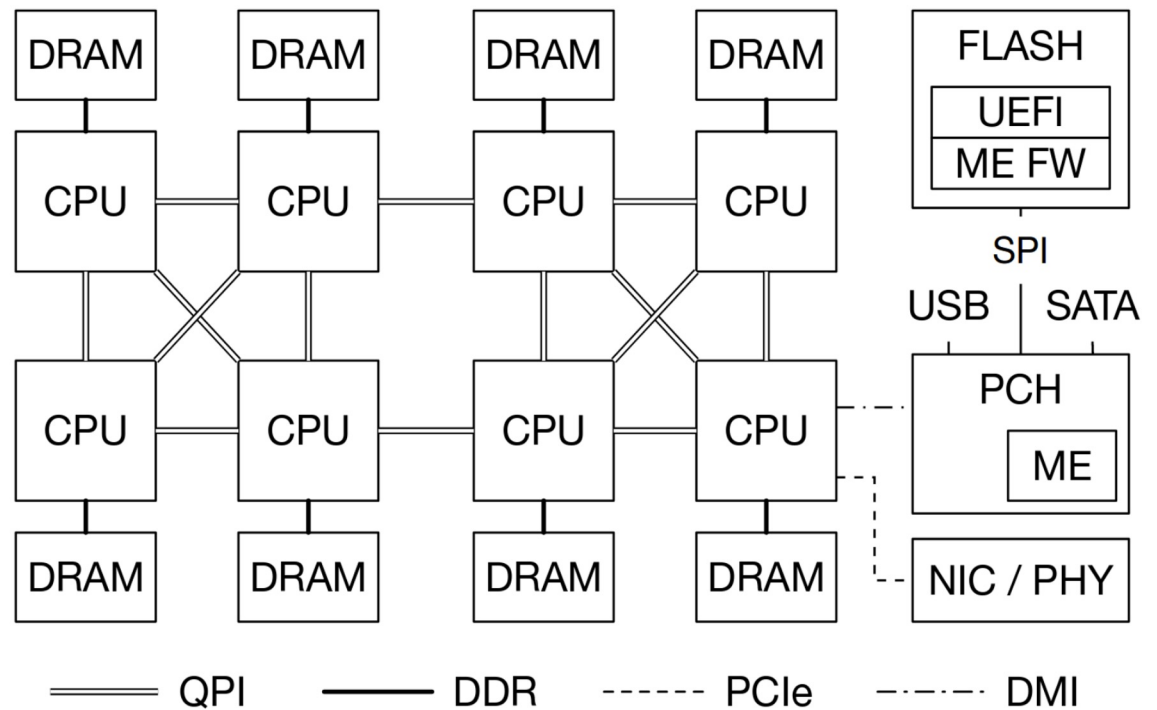
childno = _ENDPOINT_P(m_in.child_endpt); if(childno < 0 || childno >= NR_PROCS) panic(__FILE__, "FS: bogus child for forking", m_in.child_endpt); ...

```
/* PM gives child endpoint, which implies process slot information.
 * Don't call isokendpt, because that will verify if the endpoint
 * number is correct in fproc, which it won't be.
 */
```

```
childno = _ENDPOINT_P(m_in.child_endpt);
if(childno < 0 || childno >= NR_PROCS)
    panic(__FILE__, "FS: bogus child for forking", m_in.child_endpt);
if(fproc[childno].fp_pid != PID_FREE)
    panic(__FILE__, "FS: forking on top of in-use child", childno);
```

ME starts first

- Reads its initialization code from the BIOS chip
- Via the SPI bus



Bootstrap processor (BSP)

- One of the logical processors is chosen as **bootstrap processor** (BSP)
 - Will start initialization
- Others become “**application processors**” (AP)
 - Waiting for a special interrupt from the BSP

BSP starts reading BIOS

- Executes instructions stored in the BIOS chip
- An interesting detail is that BSP starts with DRAM disabled
 - Hence there is no stack to call functions
 - What can be done?

BSP starts without DRAM

- Custom-written assembly code that uses no stack
- Or a ROMCC compiler
 - Generates code from C that uses no stack
 - Used in the coreboot project

Cache-as-RAM

- Use CPU caches as temporary replacement for RAM
 - Initialize DRAM
 - Copy BIOS firmware into DRAM and continue

BIOS firmware

Initialize

- Interrupt controllers
- Devices, e.g., network interfaces
- If one of PCI devices contains “option ROM” load and execute it
 - Network cards may contain iPXE ROM
 - Implement boot from the network host

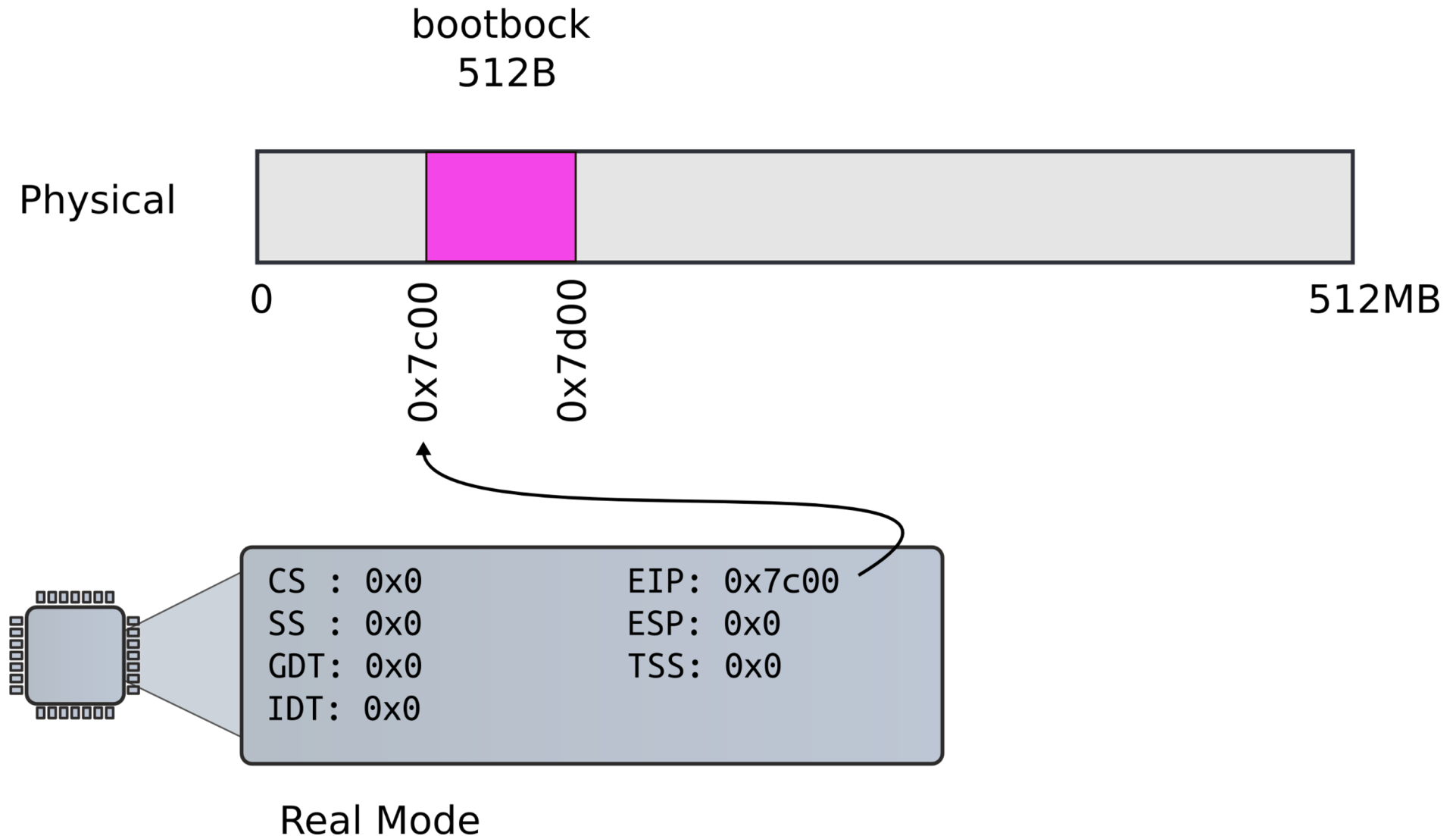
System Management Mode

- Another compartment that runs underneath your OS or a hypervisor
- Initialized by BIOS
- Protected with hardware memory mechanisms
- OS cannot access this region of memory
- Runs under your OS and the hypervisor
- Receives interrupts periodically, can take over the entire system any time
- No way to disable

BIOS loads the boot loader

- BIOS ends by loading a **boot loader**
- Modern BIOSes can load the boot loader from a variety of sources (hard disks, USB drives, optical disks)
- Default way is to load the **first sector** (512 bytes) from disk into the memory location at **0x7c00**
- BIOS then starts executing instructions at the address 0x7c00
 - This is exactly what we see when we run xv6 under QEMU
 - QEMU emulates hardware: runs BIOS, follows the same protocol

BIOS loads bootloader



Outline of the boot sequence

Outline of the boot sequence

- Setup segments (data and code)
- Switch to protected mode
- Load GDT (segmentation is on)
- Setup stack (to call C functions)
- Load kernel from disk
- Setup first page table
 - 2 entries [0 : 4MB] and [2GB : (2GB + 4MB)]
- Setup high-address stack
- Jump to main()

Bootloader starts

```
9111 start:
9112     cli # BIOS enabled interrupts; disable
9113
9114     # Zero data segment registers DS,ES,and
9115                                     SS.
9115     xorw %ax,%ax # Set %ax to zero
9116     movw %ax,%ds # -> Data Segment
9117     movw %ax,%es # -> Extra Segment
9118     movw %ax,%ss # -> Stack Segment
```

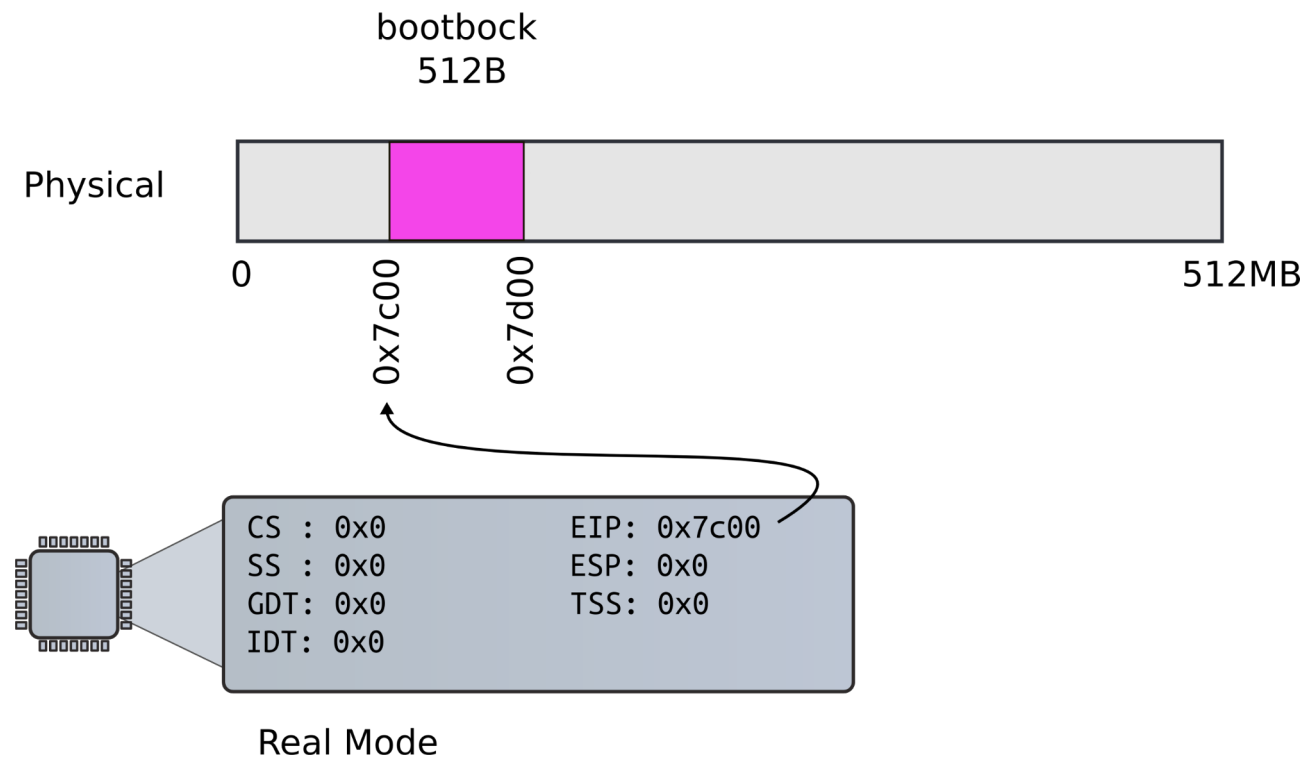
xv6/bootasm.S [bootloader]

Why start happens to be 0x7c00?

9111 start:

9112 cli # BIOS enabled interrupts; disable

9113



xv6/bootasm.S [bootloader]

Linker is instructed to link the boot block code in the Makefile

```
9111 start:
```

```
9112     cli # BIOS enabled interrupts; disable
```

```
9113
```

```
bootblock: bootasm.S bootmain.c
```

```
    $(CC) $(CFLAGS) -fno-pic -O -nostdinc -I. -c bootmain.c
```

```
    $(CC) $(CFLAGS) -fno-pic -nostdinc -I. -c bootasm.S
```

```
    $(LD) $(LDFLAGS) -N -e start -Ttext 0x7C00 -o bootblock.o bootasm.o  
bootmain.o
```

```
    $(OBJDUMP) -S bootblock.o > bootblock.asm
```

```
    $(OBJCOPY) -S -O binary -j .text bootblock.o bootblock
```

```
    ./sign.pl bootblock
```

xv6/Makefile

Switch to protected mode

- Switch from real to protected mode
- Use a bootstrap GDT that makes virtual addresses map directly to physical addresses so that the effective memory map doesn't change during the transition.

```
9141 lgdt gdtdesc
```

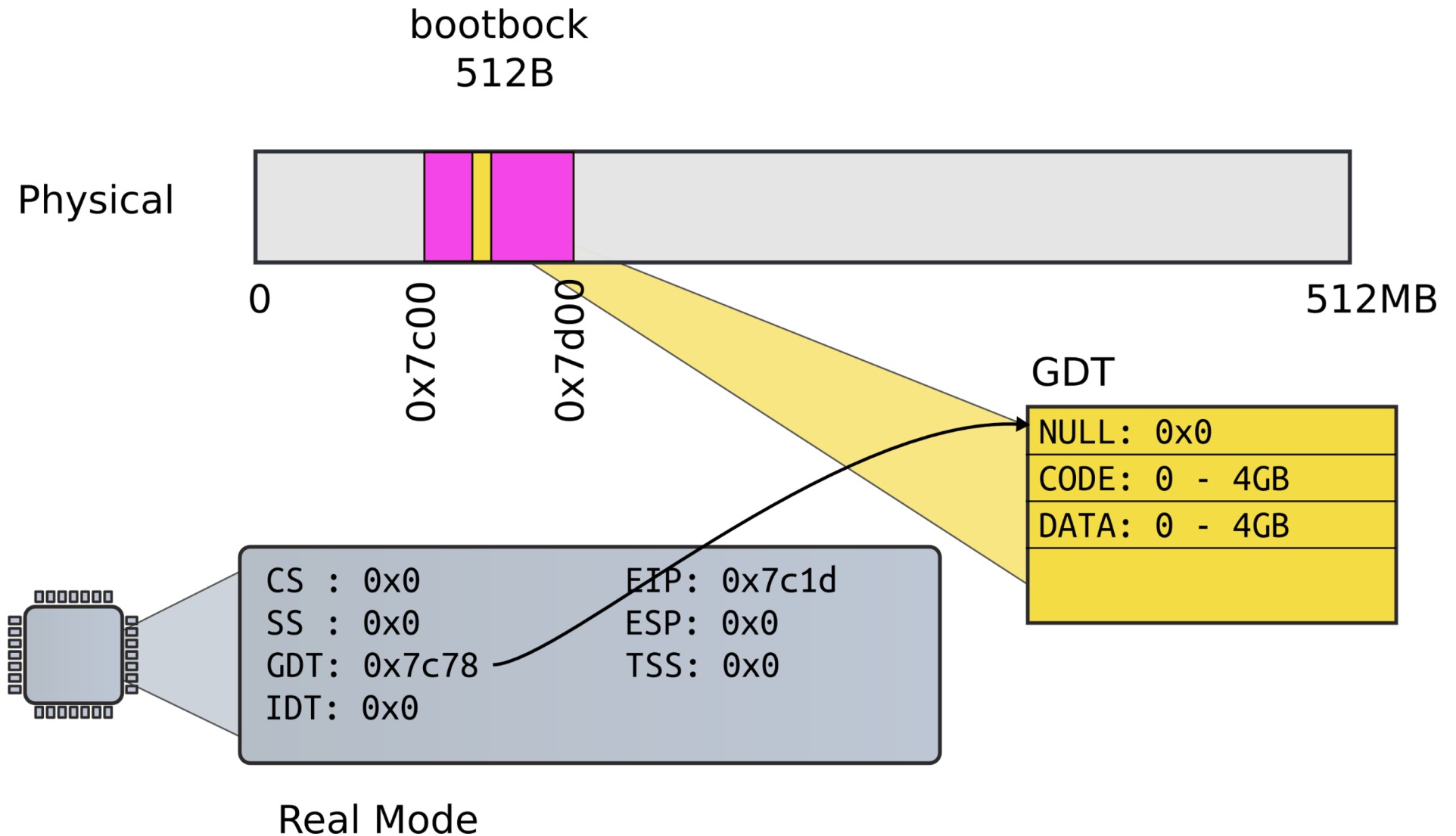
```
9142 movl %cr0, %eax
```

```
9143 orl $CR0_PE, %eax
```

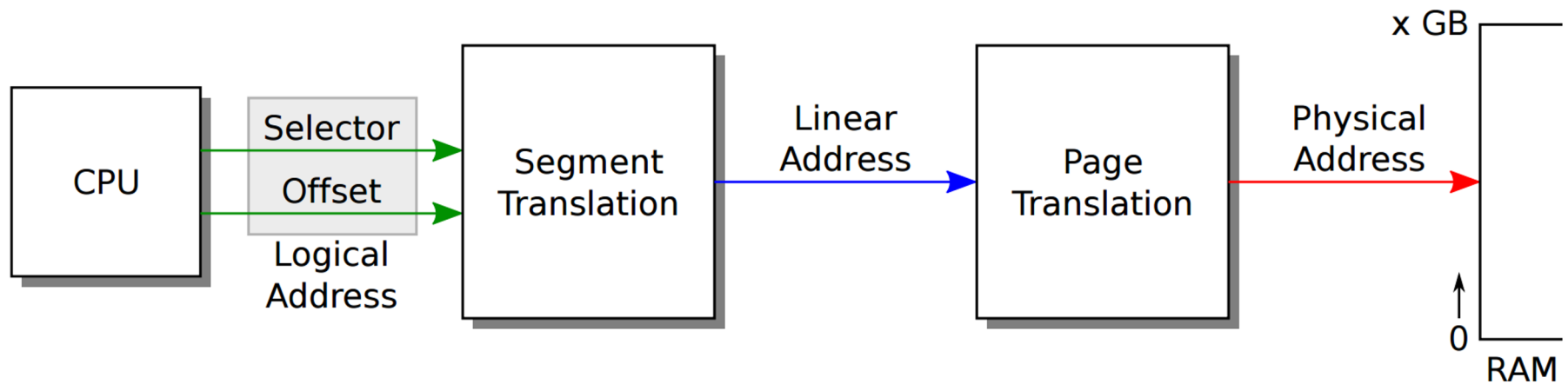
```
9144 movl %eax, %cr0
```

```
xv6/bootasm.S [bootloader]
```

Load GDT



Recap: complete address translation



How GDT is defined

```
9180 # Bootstrap GDT
9181 .p2align 2 # force 4 byte alignment
9182 gdt:
9183     SEG_NULLASM # null seg
9184     SEG_ASM(STA_X|STA_R, 0x0, 0xffffffff) # code seg
9185     SEG_ASM(STA_W, 0x0, 0xffffffff) # data seg
9186
9187 gdtdesc:
9188     .word (gdtdesc - gdt - 1) # sizeof(gdt) - 1
9189     .long gdt
```

xv6/bootasm.S [bootloader]

How GDT is defined

```
9180 # Bootstrap GDT
9181 .p2align 2 # force 4 byte alignment
9182 gdt:
9183     SEG_NULLASM # null seg
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9185     SEG_ASM(STA_W, 0x0, 0xffffffff) # data seg
9186
9187 gdtdesc:
9188     .word (gdtdesc - gdt - 1) # sizeof(gdt) - 1
9189     .long gdt
```

xv6/bootasm.S [bootloader]

Actual switch

- Use long jump to change code segment

```
9153  ljmp $(SEG_KCODE<<3), $start32
```

- Explicitly specify code segment, and address
- Segment is 0b1000 (0x8)

Why CS is 0x8, not 0x1?

- Segment selector:



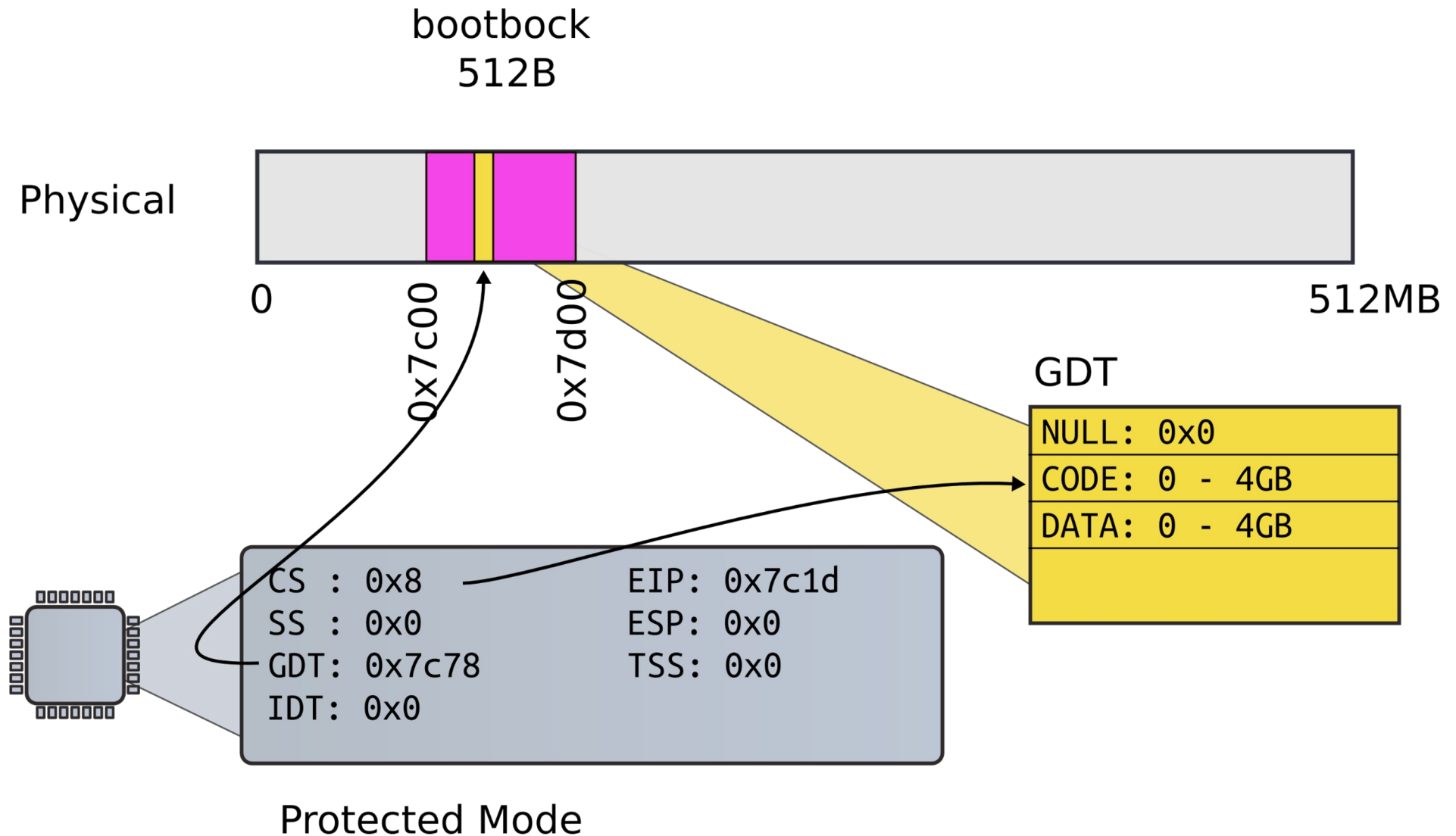
Table Indicator

0 = GDT

1 = LDT

Requested Privilege Level (RPL)

Long jump

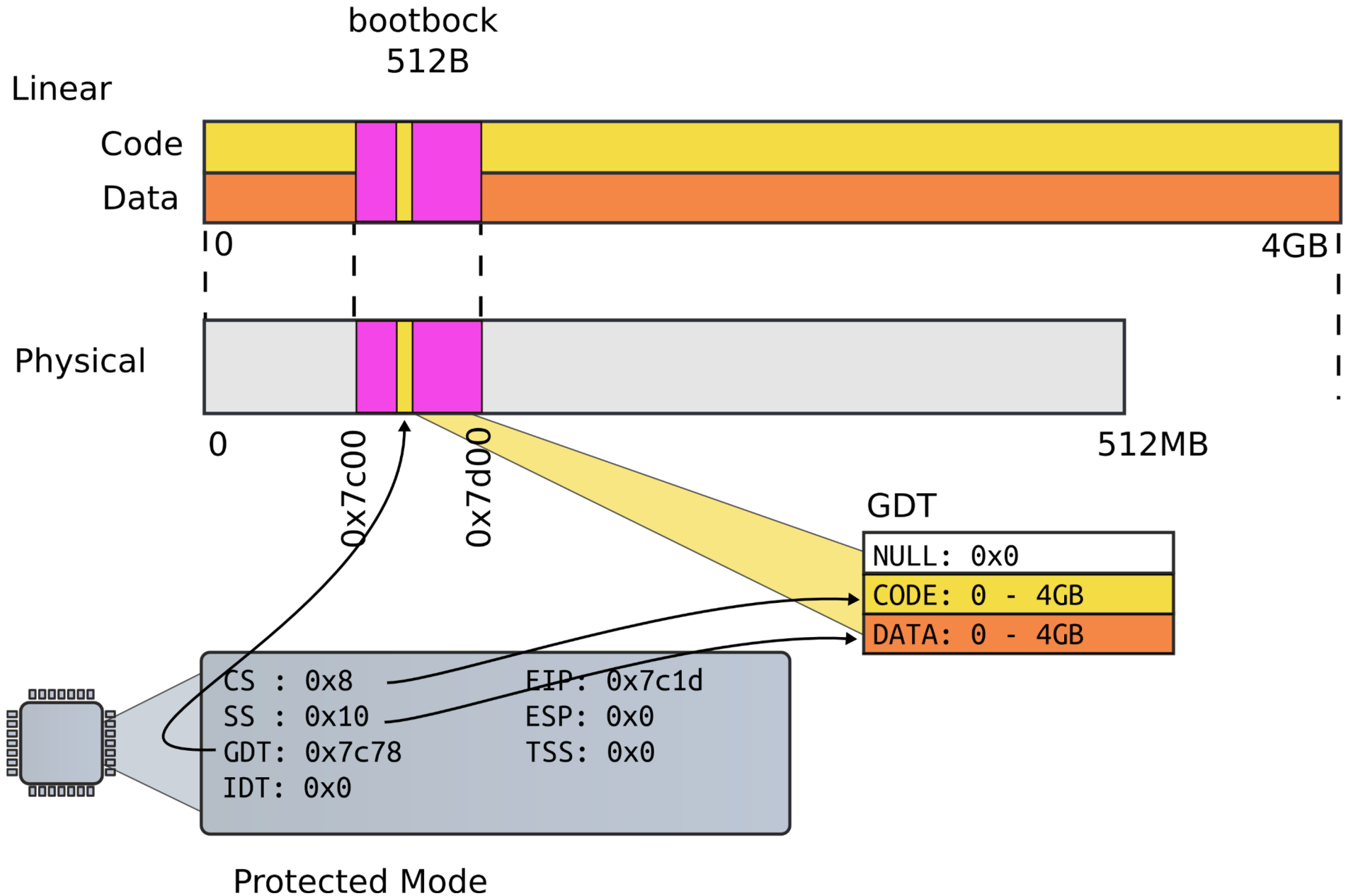


Segments

```
9155 .code32 # Tell assembler to generate 32-bit code now.
9156 start32:
9157     # Set up the protected-mode data segment registers
9158     movw $(SEG_KDATA<<3), %ax # Our data segment selector
9159     movw %ax, %ds # -> DS: Data Segment
9160     movw %ax, %es # -> ES: Extra Segment
9161     movw %ax, %ss # -> SS: Stack Segment
9162     movw $0, %ax # Zero segments not ready for use
9163     movw %ax, %fs # -> FS
9164     movw %ax, %gs # -> GS
```

xv6/bootasm.S [bootloader]

Segments



Setup stack

- Why do we need a stack?

```
9166 movl $start, %esp
```

```
9167 call bootmain
```

```
xv6/bootasm.S [bootloader]
```

Setup stack

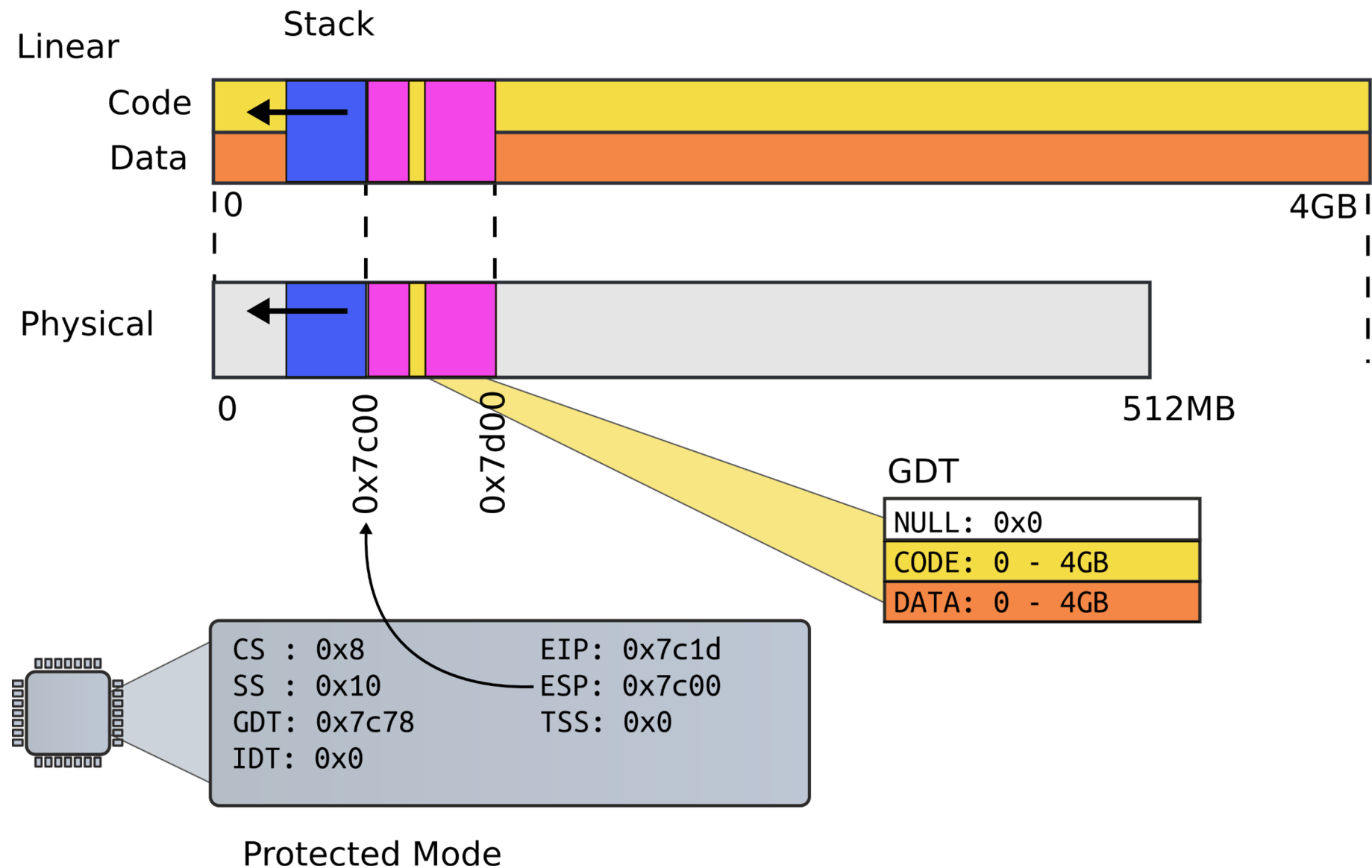
- Need stack to use C
- Function invocations
- Note, there were no stack instructions before that

```
9166 movl $start, %esp
```

```
9167 call bootmain
```

```
xv6/bootasm.S [bootloader]
```

First stack



Invoke first C function

```
9166 movl $start, %esp
```

```
9167 call bootmain
```

xv6/bootasm.S [bootloader]

bootmain(): read kernel from disk

```
9216 void
9217 bootmain(void)
9218 {
9219     struct elfhdr *elf;
9220     struct proghdr *ph, *eph;
9221     void (*entry)(void);
9222     uchar* pa;
9223
9224     elf = (struct elfhdr*)0x10000; // scratch space
9225
9226     // Read 1st page off disk
9227     readseg((uchar*)elf, 4096, 0);
9228
9229     // Is this an ELF executable?
9230     if(elf->magic != ELF_MAGIC)
9231         return; // let bootasm.S handle error
9232
```

xv6/bootmain.c [bootloader]

```
9232
9233     // Load each program segment (ignores ph flags).
9234     ph = (struct proghdr*)((uchar*)elf + elf->phoff);
9235     eph = ph + elf->phnum;
9236     for(; ph < eph; ph++){
9237         pa = (uchar*)ph->paddr;
9238         readseg(pa, ph->filesz, ph->off);
9239         if(ph->memsz > ph->filesz)
9240             stosb(pa + ph->filesz, 0, ph->memsz - ph->filesz);
9241     }
9242
9243     // Call the entry point from the ELF header.
9244     // Does not return!
9245     entry = (void(*)(void))(elf->entry);
9246     entry();
9247 }
```

bootmain(): read kernel

xv6/bootmain.c [bootloader] from disk

How do we read disk?

```
9257
9258 // Read a single sector at offset into dst.
9259 void
9260 readsect(void *dst, uint offset)
9261 {
9262     // Issue command.
9263     waitdisk();
9264     outb(0x1F2, 1); // count = 1
9265     outb(0x1F3, offset);
9266     outb(0x1F4, offset >> 8);
9267     outb(0x1F5, offset >> 16);
9268     outb(0x1F6, (offset >> 24) | 0xE0);
9269     outb(0x1F7, 0x20); // cmd 0x20 - read sectors
9270
9271     // Read data.
9272     waitdisk();
9273     insl(0x1F0, dst, SECTSIZE/4);
9274 }
```

xv6/bootmain.c [bootloader]

How do we read disk (cont)?

```
9250 void
9251 waitdisk(void)
9252 {
9253     // Wait for disk ready.
9254     while((inb(0x1F7) & 0xC0) != 0x40)
9255         ;
9256 }
9257
```

How do we read disk?

```
9257
9258 // Read a single sector at offset into dst.
9259 void
9260 readsect(void *dst, uint offset)
9261 {
9262     // Issue command.
9263     waitdisk();
9264     outb(0x1F2, 1); // count = 1
9265     outb(0x1F3, offset);
9266     outb(0x1F4, offset >> 8);
9267     outb(0x1F5, offset >> 16);
9268     outb(0x1F6, (offset >> 24) | 0xE0);
9269     outb(0x1F7, 0x20); // cmd 0x20 - read sectors
9270
9271     // Read data.
9272     waitdisk();
9273     insl(0x1F0, dst, SECTSIZE/4);
9274 }
```

xv6/bootmain.c [bootloader]

```
9232
9233     // Load each program segment (ignores ph flags).
9234     ph = (struct proghdr*)((uchar*)elf + elf->phoff);
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9236     for(; ph < eph; ph++){
9237         pa = (uchar*)ph->paddr;
9238         readseg(pa, ph->filesz, ph->off);
9239         if(ph->memsz > ph->filesz)
9240             stosb(pa + ph->filesz, 0, ph->memsz - ph->filesz);
9241     }
9242
9243     // Call the entry point from the ELF header.
9244     // Does not return!
9245     entry = (void(*)(void))(elf->entry);
9246     entry();
9247 }
```

Call kernel entry

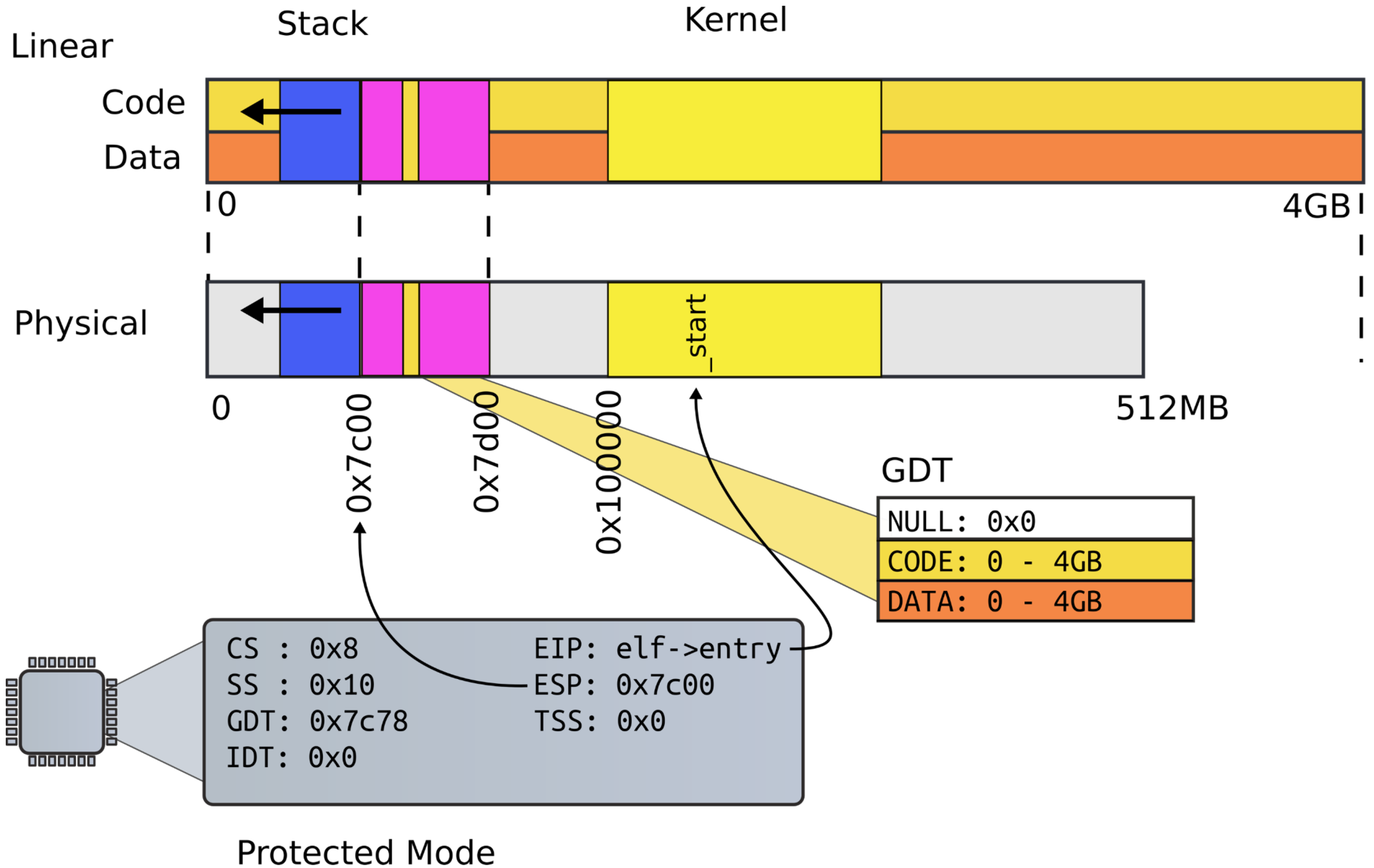
xv6/bootmain.c [bootloader]

```
1039 .globl entry
1136 # By convention, the _start symbol specifies the ELF entry point.
1137 # Since we haven't set up virtual memory yet, our entry point is
1138 # the physical address of 'entry'.
1139 .globl _start
1140 _start = V2P_W0(entry)
1141
1142 # Entering xv6 on boot processor, with paging off.
1143 .globl entry
1144 entry:
1145 # Turn on page size extension for 4Mbyte pages
1146     movl %cr4, %eax
1147     orl $(CR4_PSE), %eax
1148     movl %eax, %cr4
```

entry(): kernel ELF entry

xv6/entry.S [kernel]

Kernel



entry(): kernel ELF entry

```
1039 .globl entry
```

```
1136 # By convention, the _start symbol specifies the ELF entry  
point.
```

```
1137 # Since we haven't set up virtual memory yet, our entry point  
is
```

```
1138 # the physical address of 'entry'.
```

```
1139 .globl _start
```

```
1140 _start = V2P_W0(entry)
```

```
1141
```

```
1142 # Entering xv6 on boot processor, with paging off.
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```
1143 .globl entry
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```
1144 entry:
```

```
1145 # Turn on page size extension for 4Mbyte pages
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```
1146     movl %cr4, %eax
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1147     orl $(CR4_PSE), %eax
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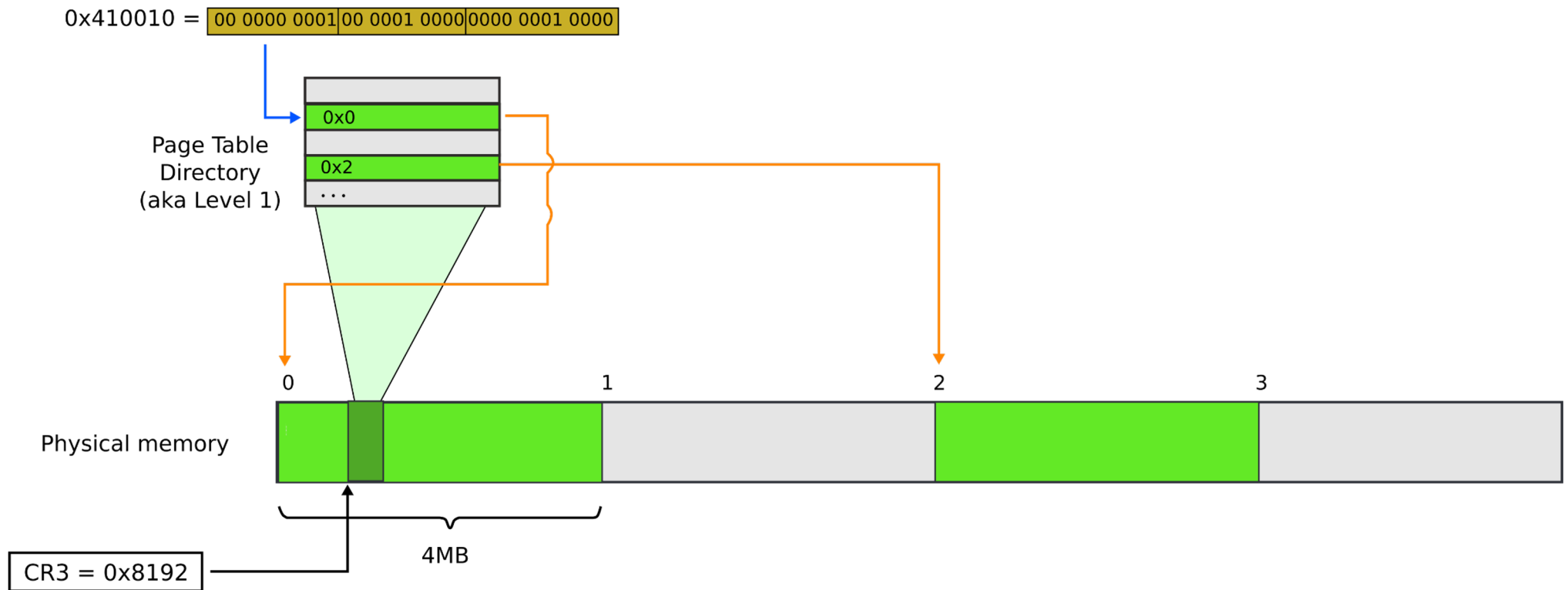
```
1148     movl %eax, %cr4
```

xv6/entry.S [kernel]

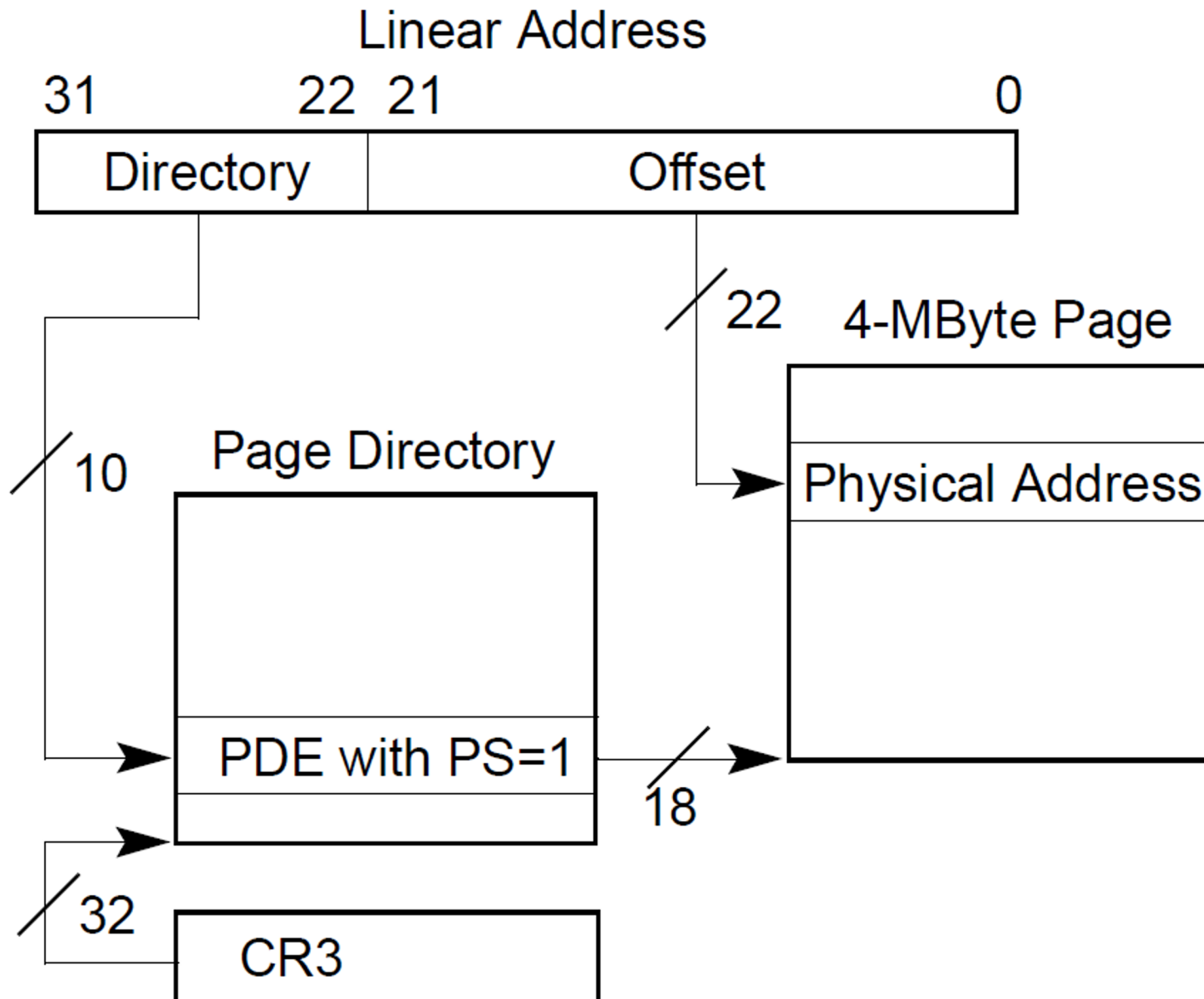
32bit x86 supports two page sizes

- 4KB pages
- 4MB pages

Page translation for 4MB pages



Page translation for 4MB pages

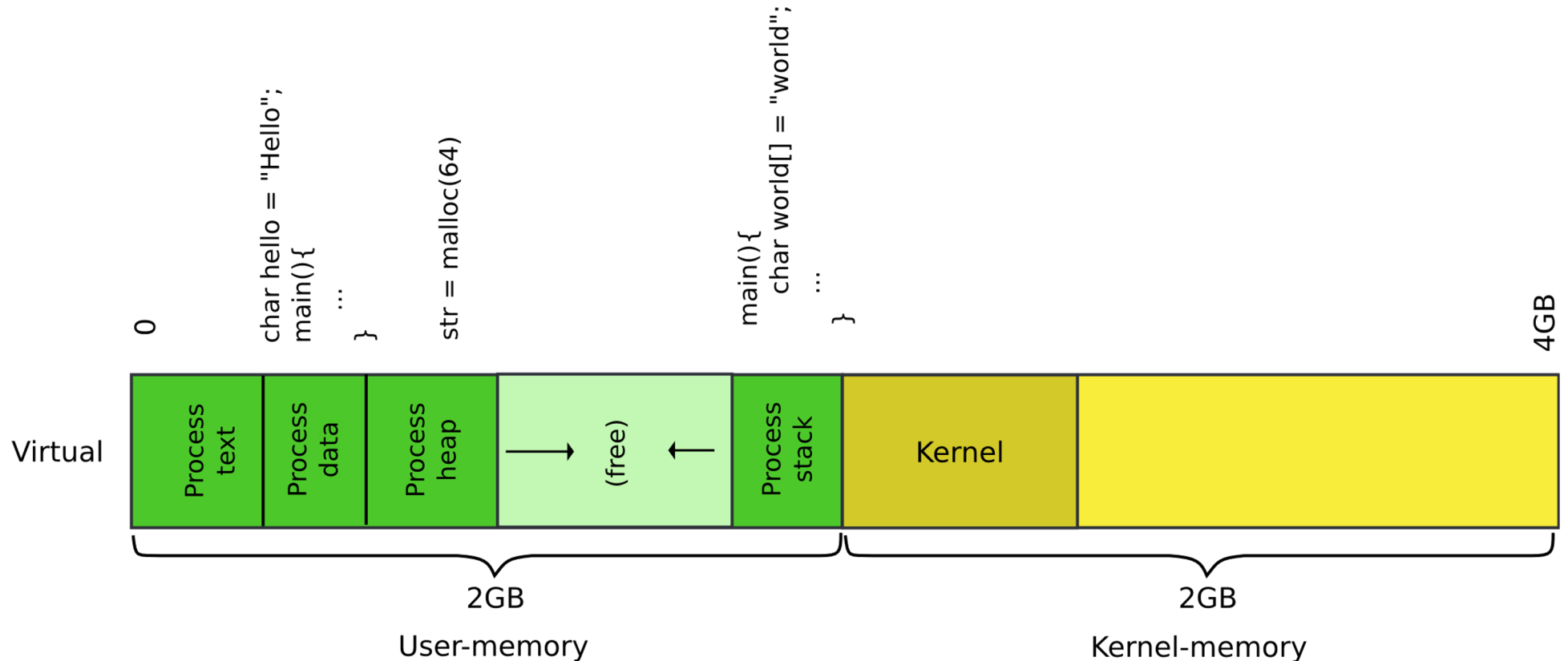


Set up page directory

```
1149 # Set page directory
1150 movl $(V2P_W0(entrypgdir)), %eax
1151 movl %eax, %cr3
```

xv6/entry.S [kernel]

Our goal: 2GB/2GB address space



First page table

- Two 4MB entries (large pages)
- Entry #0
 - 0x0 – 4MB → 0x0:0x400000
- Entry #512
 - 0x0 – 4MB → 0x80000000:0x80400000


```
1406 // The boot page table used in entry.S and entryother.S.
1407 // Page directories (and page tables) must start on page
1408 // boundaries,
1409 // hence the __aligned__ attribute.
1410 // PTE_PS in a page directory entry enables 4Mbyte pages.
```

```
1411 __attribute__((__aligned__(PGSIZE)))
1412 pde_t entrypgdir[NPDENTRIES] = {
1413     // Map VA's [0, 4MB) to PA's [0, 4MB)
1414     [0] = (0) | PTE_P | PTE_W | PTE_PS,
1415     // Map VA's [KERNBASE, KERNBASE+4MB) to PA's [0, 4MB)
1416     [KERNBASE>>PDXSHIFT] = (0) | PTE_P | PTE_W | PTE_PS,
1417 };
```

First page table

```
1406 // The boot page table used in entry.S and entryother.S.
1407 // Page directories (and page tables) must start on page
1408 // boundaries,
1409 // hence the __aligned__ attribute.
1410 // PTE_PS in a page directory entry enables 4Mbyte pages.
1411 __attribute__((__aligned__(PGSIZE)))
1412 pde_t entrypgdir[NPDENTRIES] = {
1413     // Map VA's [0, 4MB) to PA's [0, 4MB)
1414     [0] = (0) | PTE_P | PTE_W | PTE_PS,
1415     // Map VA's [KERNBASE, KERNBASE+4MB) to PA's [0, 4MB)
1416     [KERNBASE>>PDXSHIFT] = (0) | PTE_P | PTE_W | PTE_PS,
1417 };
```

First page table

First page table (cont)

```
0870 // Page directory and page table constants.  
0871 #define NPDENTRIES 1024
```

```
1406 // The boot page table used in entry.S and entryother.S.
1407 // Page directories (and page tables) must start on page
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1416     [KERNBASE>>PDXSHIFT] = (0) | PTE_P | PTE_W | PTE_PS,
1417 };
```

First page table

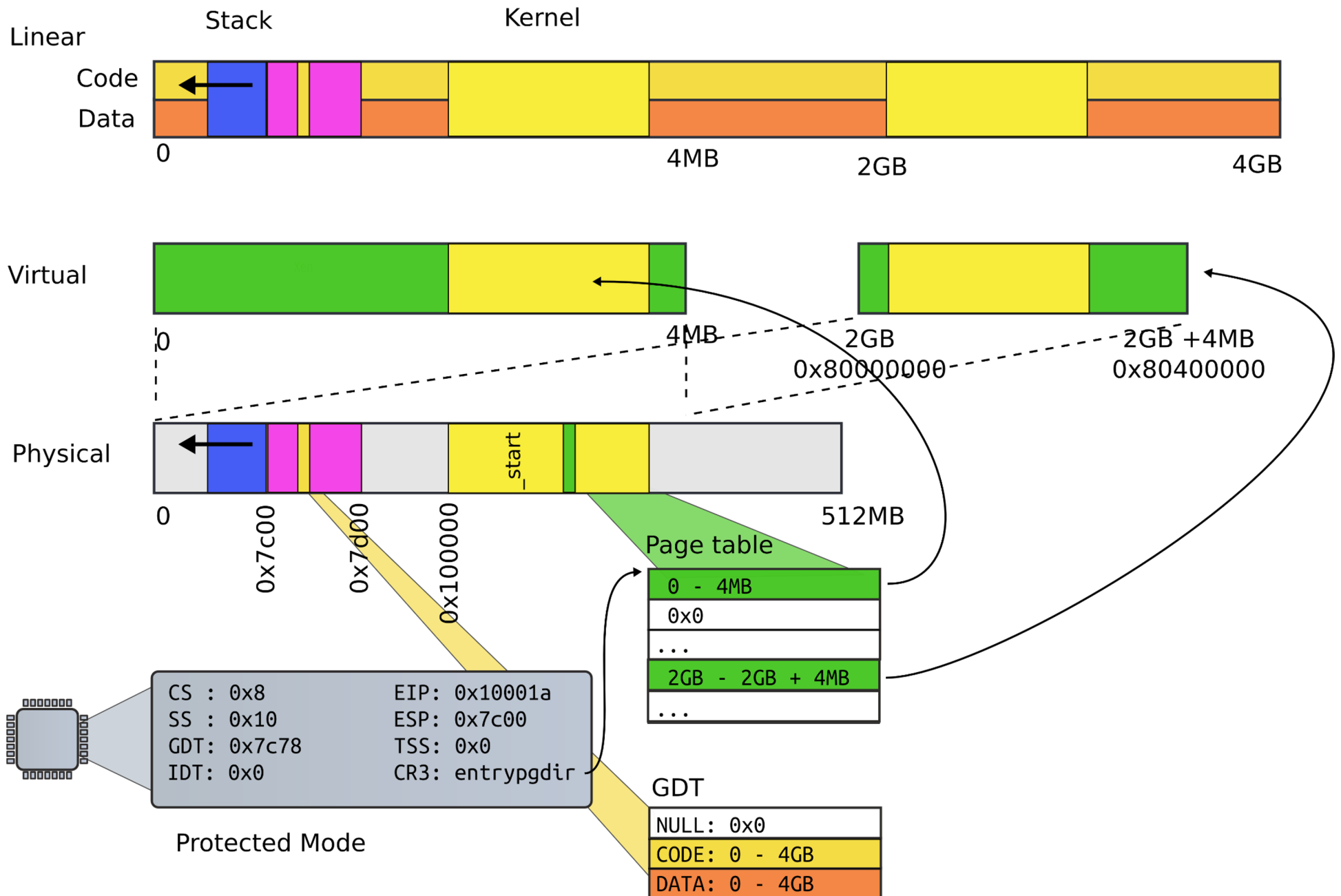
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```

First page table

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1413     // Map VA's [0, 4MB) to PA's [0, 4MB)
1414     [0] = (0) | PTE_P | PTE_W | PTE_PS,
1415     // Map VA's [KERNBASE, KERNBASE+4MB) to PA's [0, 4MB)
1416     [KERNBASE>>PDXSHIFT] = (0) | PTE_P | PTE_W | PTE_PS,
1417 };
```

First page table

First page table



Turn on paging

```
1152 # Turn on paging.  
1153 movl %cr0, %eax  
1154 orl $(CR0_PG|CR0_WP), %eax  
1155 movl %eax, %cr0
```

xv6/entry.S [kernel]

High address stack (4K)

```
1157 # Set up the stack pointer.
```

```
1158 movl $(stack + KSTACKSIZE), %esp
```

```
1159
```

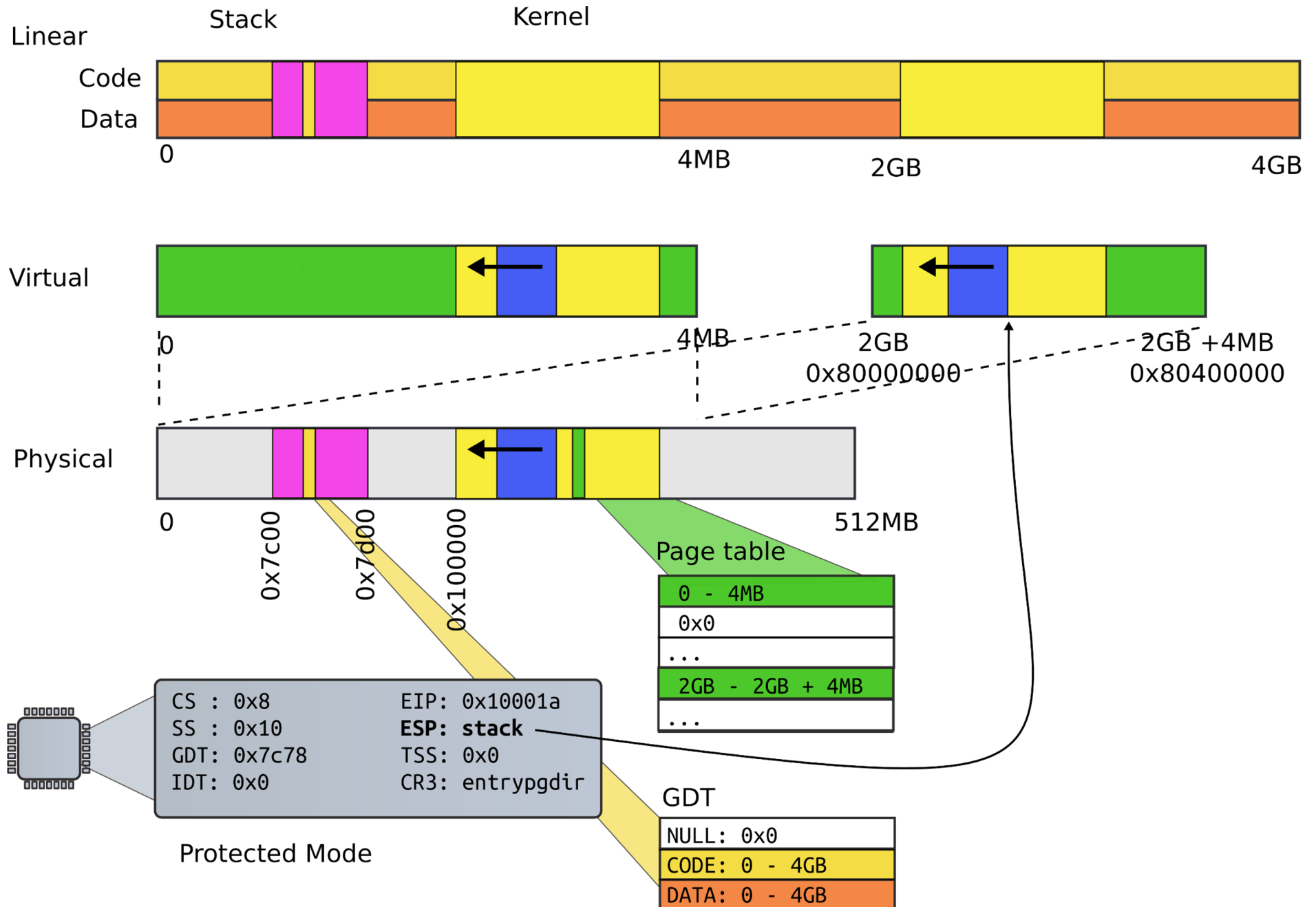
```
...
```

```
1167 .comm stack, KSTACKSIZE
```

```
0151 #define KSTACKSIZE 4096 // size of  
                                per-process kernel stack
```

xv6/entry.S [kernel]

High address stack (4K)



Jump to main()

```
1160 # Jump to main(), and switch to executing at
1161 # high addresses. The indirect call is
      needed because
1162 # the assembler produces a PC-relative
      instruction
1163 # for a direct jump.
1164 mov $main, %eax
1165 jmp *%eax
1166
```

xv6/entry.S [kernel]

Running in main()

```
1313 // Bootstrap processor starts running C code here.
1314 // Allocate a real stack and switch to it, first
1315 // doing some setup required for memory allocator to work.
1316 int
1317 main(void)
1318 {
1319     kinit1(end, P2V(4*1024*1024)); // phys page allocator
1320     kvmalloc(); // kernel page table
1321     mpinit(); // detect other processors
1322     lapicinit(); // interrupt controller
1323     seginit(); // segment descriptors
1324     cprintf("\ncpu%d: starting xv6\n\n", cpunum());
1325     ...
1340 }
```

xv6/main.c [kernel]

Recap of the boot sequence

- Setup segments (data and code)
- Switched to protected mode
- Loaded GDT (segmentation is on)
- Setup stack (to call C functions)
- Loaded kernel from disk
- Setup first page table
 - 2 entries [0 : 4MB] and [2GB : (2GB + 4MB)]
- Setup high-address stack
- Jumped to main()

Conclusion

- We've booted
- We're running in main()

Thank you!

References

- [1] Costan, Victor, and Srinivas Devadas. "Intel SGX Explained." IACR Cryptology ePrint Archive 2016 (2016): 86.
<https://eprint.iacr.org/2016/086.pdf>


```
1. #include <stdio.h>

2. void func_a(void){
3.     printf("func_a\n");
4.     return;
5. }

6. void func_b(void) {
7.     printf("func_b\n");
8.     return;
9. }

10. int main(int ac, char **av)
11. {
12.     void (*fp)(void);

13.     fp = func_b;
14.     fp();
15.     return;
16. }
```

Function pointers

08048432 <func_b>:

8048432:	55	push	%ebp
8048433:	89 e5	mov	%esp,%ebp
8048435:	83 ec 18	sub	\$0x18,%esp
8048438:	c7 04 24 07 85 04 08	movl	\$0x8048507,(%esp)
804843f:	e8 ac fe ff ff	call	80482f0 <puts@plt>
8048444:	90	nop	
8048445:	c9	leave	
8048446:	c3	ret	

08048447 <main>:

8048447:	55	push	%ebp
8048448:	89 e5	mov	%esp,%ebp
804844a:	83 e4 f0	and	\$0xfffffffff0,%esp
804844d:	83 ec 10	sub	\$0x10,%esp

Load pointer to func_p on

the stack

8048450:	c7 44 24 0c 32 84 04	movl	
\$0x8048432,0xc(%esp)			
8048457:	08		
8048458:	8b 44 24 0c	mov	0xc(%esp),%eax
804845c:	ff d0	call	*%eax
804845e:	90	nop	
804845f:	c9	leave	
8048460:	c3	ret	

Function pointers

