# Latency Numbers Every Programmer Should Know (Dean 2009)

LI cache reference	<sup>1</sup> / <sub>2</sub> ns		
Branch mispredict	5 ns		
L2 cache reference	7 ns		
Mutex lock/unlock	25 ns		
Main memory reference	100 ns		
Compress I K bytes with Zippy	3,000 ns	3 µs	
Send I K bytes over I Gbps network	10,000 ns	10 µs	
Read 4 K randomly from SSD	150,000 ns	150 μs	
Read I MB sequentially from memory	250,000 ns	250 µs	
Round trip within same datacenter	500,000 ns	500 µs	
Read I MB sequentially from SSD	1,000,000 ns	I,000 µs	l ms
Disk seek	10,000,000 ns	10,000 µs	10 ms
Read I MB sequentially from disk	20,000,000 ns	20,000 µs	20 ms
Send packet CA→Netherlands→CA	150,000,000 ns	Ι 50,000 μs	150 ms











# Propagation Delay



Time to travel along medium

# Propagation Delay



Time to travel along medium, depends on physics

- distance travelled
- speed of link some fraction of the speed of light

 $delay_{propagation} = \frac{scale \times distance}{speed of light}$ 

## Transmission Delay



Time to convert from bytes to medium

## Transmission Delay



Time to convert from bytes to medium, depends on device and data

- device rate R
  - ° fiber, ethernet: I-400 Gbps
  - WiFi: ~10 Mbps
  - dial-up: 56.6 Kbps
- packet size

$$delay_{transmission} = \frac{size}{R}$$

# Processing Delay



Time to inspect bytes and choose next step

# Processing Delay



Time to inspect bytes and choose next step, depends on device speed

Typically a few nanoseconds, so we ignore it

 $delay_{processing} = 0$ 



Time packets bytes are held in a queue



Time packets bytes are held in a queue, depends on traffic





Time packets bytes are held in a queue, depends on traffic

*incoming data rate* = *average packet size* × *incoming packet rate* 

traffic intensity = 
$$\frac{incoming \ data \ rate}{R}$$

*traffic intensity*  $\leq 1 \Rightarrow$  no delay



Time packets bytes are held in a queue, depends on traffic

*incoming data rate* = *average packet size* × *incoming packet rate* 

traffic intensity = 
$$\frac{incoming \ data \ rate}{R}$$

*traffic intensity* 
$$> 1 \Rightarrow$$
 delay growing



Time packets bytes are held in a queue, depends on traffic

 $delay_{queue} = ???$ 

## Total Delay

 $delay = delay_{\text{propagation}} + delay_{\text{transmission}} + delay_{\text{processing}} + delay_{\text{queueing}}$ 

$$delay_{propagation} = \frac{scale \times distance}{speed of light}$$

$$delay_{transmission} = \frac{size}{R}$$

$$delay_{processing} = 0$$

 $delay_{queue} = ???$ 



















```
laptop$ ping www.cs.utah.edu
PING wp.wpenginepowered.com (141.193.213.10): 56 data bytes
64 bytes from 141.193.213.10: icmp_seq=0 ttl=51 time=35.962 ms
64 bytes from 141.193.213.10: icmp_seq=1 ttl=51 time=28.266 ms
64 bytes from 141.193.213.10: icmp seq=2 ttl=51 time=34.257 ms
```

```
64 bytes from 141.193.213.10: icmp_seq=4 ttl=51 time=135.983 ms

^C

--- wp.wpenginepowered.com ping statistics ---

5 packets transmitted, 5 packets received, 0.0% packet loss
```

64 bytes from 141.193.213.10: icmp seq=3 ttl=51 time=37.075 ms

round-trip min/avg/max/stddev = 28.266/54.309/135.983/40.950 ms

## Inferring Delay Components

Suppose that we take 5 pings using a packet of size 100:

	22ms
l2ms	
	45ms
19ms	
	29ms

and for the same destination, 5 pings using a packet of size 200:

19ms		
	] 46ms	
l 4ms		
		81 ms
I 5ms		

Each measurement is a *delay* 

 $delay = delay_{\text{propagation}} + delay_{\text{transmission}} + delay_{\text{processing}} + delay_{\text{queueing}}$ 

## Inferring Delay Components

Suppose that we take 5 pings using a packet of size 100:

	22ms
I2ms	
	45ms
	19ms
	29ms

and for the same destination, 5 pings using a packet of size 200:

		19ms		
			46ms	
		l 4ms		
				81 ms
		] 15ms		
Each me	easi	rement is a <i>delay</i>		



application	Firefox, ping,
transport	TCP, UDP, ICMP,
network	IP
link	ethernet, WiFi,
physical	electrons, photons,



















traceroute uses this trick systematically to explore the network

#### laptop\$ traceroute www.cs.utah.edu

traceroute: Warning: www.cs.utah.edu has multiple addresses; using 141.193.213.10 traceroute to wp.wpenginepowered.com (141.193.213.10), 64 hops max, 52 byte packets

- 1 10.0.0.1 (10.0.0.1) 11.987 ms 4.197 ms 4.602 ms
- 2 100.93.170.195 (100.93.170.195) 15.651 ms 100.93.170.194 (100.93.170.194) 18.858 ms 100.93.170.195 (100.93.170.195) 16.754 ms
- 3 po-333-417-rur501.saltlakecity.ut.utah.comcast.net (96.216.76.73) 16.825 ms po-333-418-rur502.saltlakecity.ut.utah.comcast.net (96.216.76.81) 15.903 ms po-333-417-rur501.saltlakecity.ut.utah.comcast.net (96.216.76.73) 18.208 ms

```
• • • •
```

- 15 50.242.151.238 (50.242.151.238) 33.149 ms 172.69.132.4 (172.69.132.4) 39.945 ms 66.208.229.106 (66.208.229.106) 42.951 ms
- 16 141.193.213.10 (141.193.213.10) 33.762 ms 172.71.156.2 (172.71.156.2) 60.377 ms 141.193.213.10 (141.193.213.10) 32.911 ms

## Latency vs. Throughput

Latency ⇒ how long you have to wait for one small thing

 a time, such as milliseconds
 RTT can help us understand latency

• **Throughput** ⇒ how long you have to wait for everything

a rate, such as bytes per second

Mailing a box of flash drives can have very high throughput, but also high latency

#### Summary

Network delays are largely beyond our control, but we can reason about them

Four kinds of delay add up:

- propagation depends on distance
- **transmission** depends on size and medium
- processing practically 0
- **queuing** random