

TCP Congestion Control

Congestion control in TCP avoids overwhelming a network

- how *often* to send $\Rightarrow$ timeout
- how *much* to send $\Rightarrow$ window

TCP Timeout

TCP timeout uses an estimate of RTT and its variance:

$$\text{timeout} = \text{RTT}_{est} + 4 \times \text{var}_{est}$$

- RTT_{est} is an exponentially weighted moving average:

$$\text{RTT}_{est} = (1-\alpha) \times \text{RTT}_{est} + \alpha \times \text{RTT}_{measured}$$

TCP uses $\alpha = \frac{1}{8}$

- var_{est} is similarly weighted:

$$\text{var}_{est} = (1-\beta) \times \text{var}_{est} + \beta \times |\text{RTT}_{measured} - \text{RTT}_{est}|$$

TCP uses $\beta = \frac{1}{4}$

TCP Congestion Control

Flow control relies on `rwnd` and `window ≤ rwnd`

TCP Congestion Control

from TCP header

Flow control relies on `rwnd` and `window ≤ rwnd`

TCP Congestion Control

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Congestion control uses `cwnd` and `window ≤ cwnd`

- start `cwnd` at maximum segment size, MSS
- grow until `ssthresh`, which starts large, but can adapt

TCP Congestion Control

Flow control relies on `rwnd` and `window ≤ rwnd`

Congestion control uses `cwnd` and `window ≤ cwnd`

- start `cwnd` at maximum segment size, MSS  about 1.5KB
- grow until `ssthresh`, which starts large, but can adapt

TCP Congestion Control

Flow control relies on `rwnd` and `window ≤ rwnd`

Congestion control uses `cwnd` and `window ≤ cwnd`

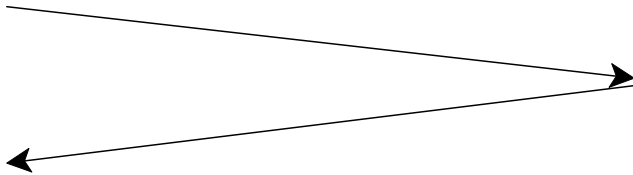
- start `cwnd` at maximum segment size, MSS
- grow until `ssthresh`, which starts large, but can adapt

start at 64KB

TCP Slow Start

sender

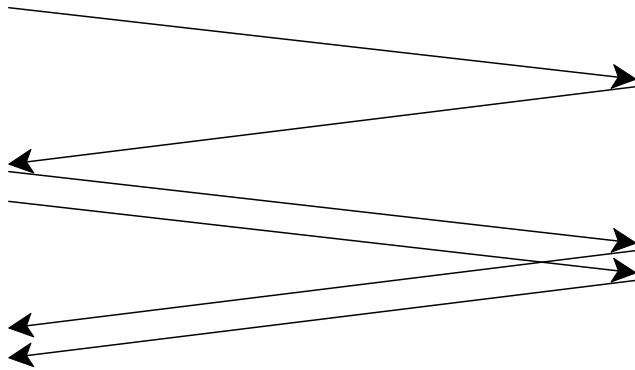
receiver



TCP Slow Start

sender

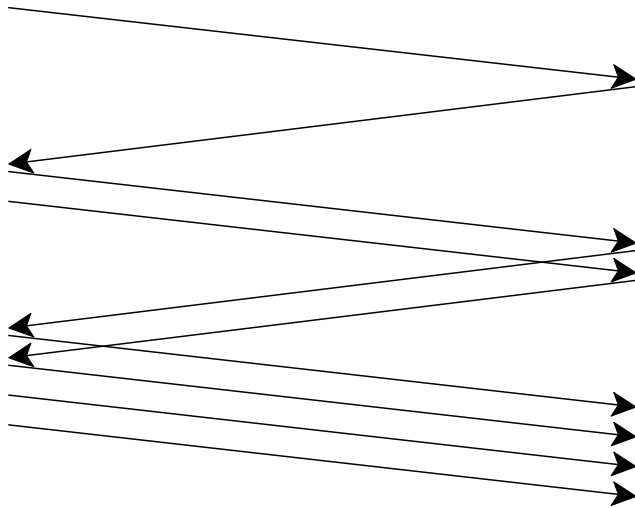
receiver



TCP Slow Start

sender

receiver

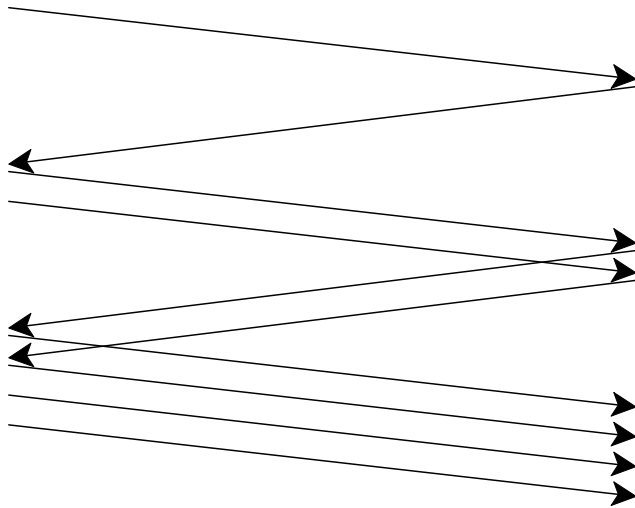


TCP Slow Start

sender

receiver

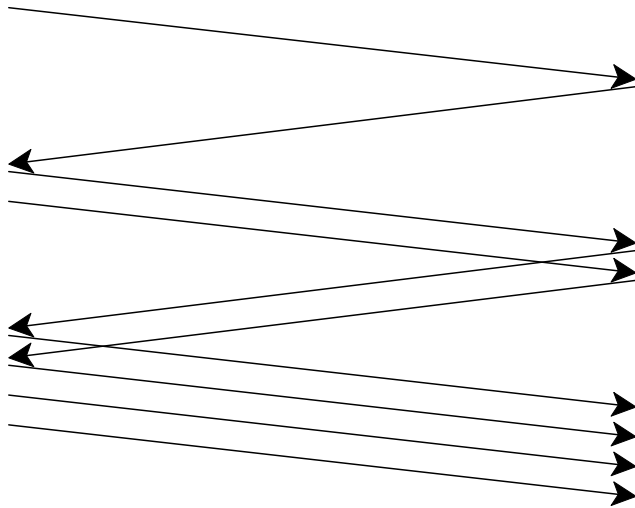
Initially, $\text{cwnd} = \text{MSS}$



TCP Slow Start

sender

receiver



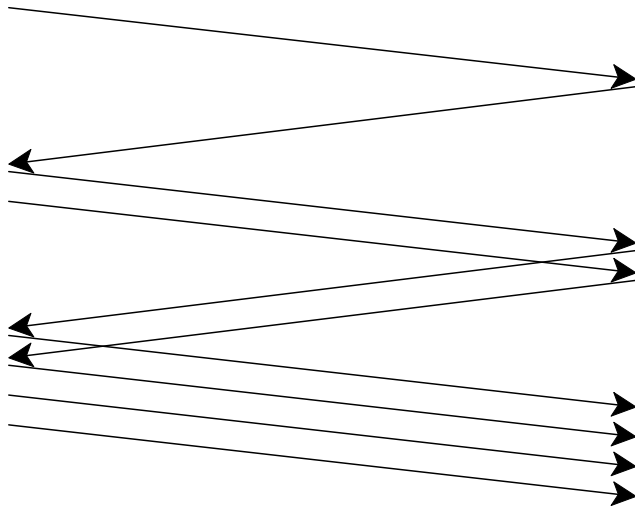
Initially, $\text{cwnd} = \text{MSS}$

After each ACK, $\text{cwnd} += \text{MSS}$

$\Rightarrow$ double cwnd each RTT

TCP Slow Start

sender **receiver**

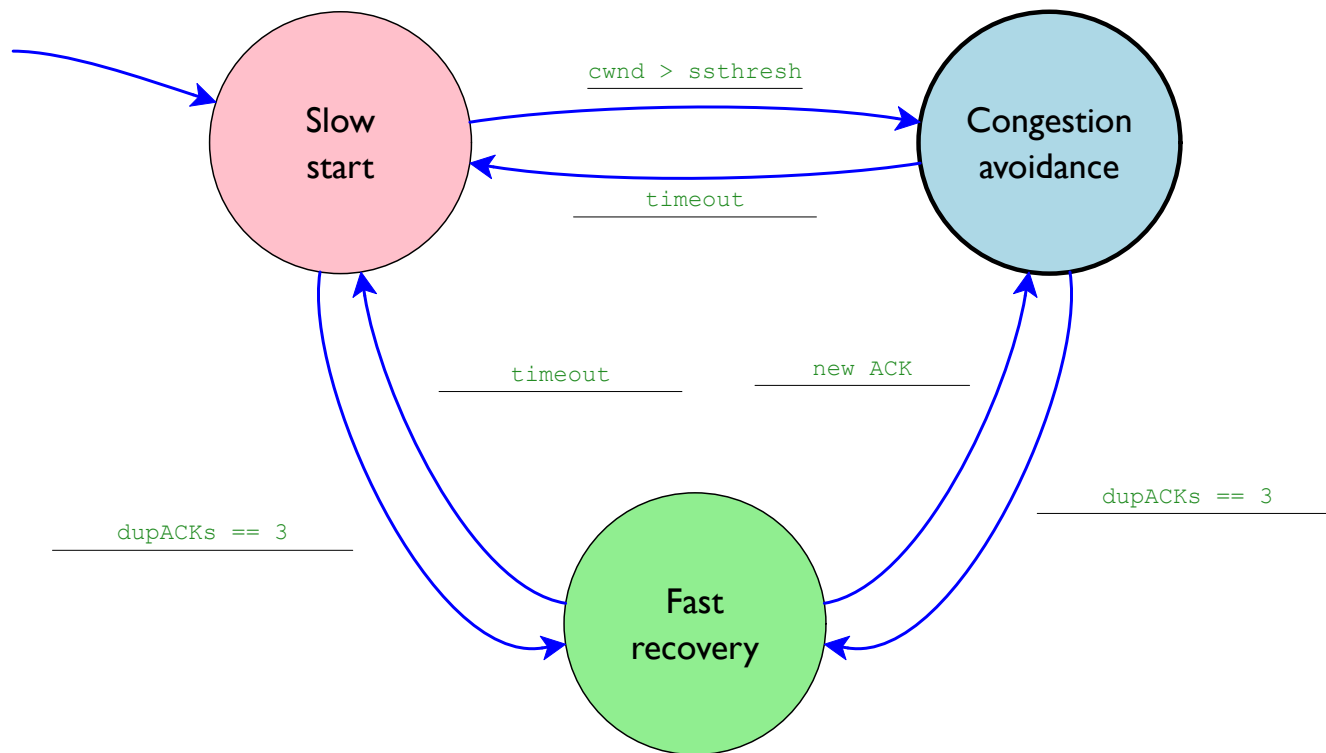


Initially, $cwnd = MSS$

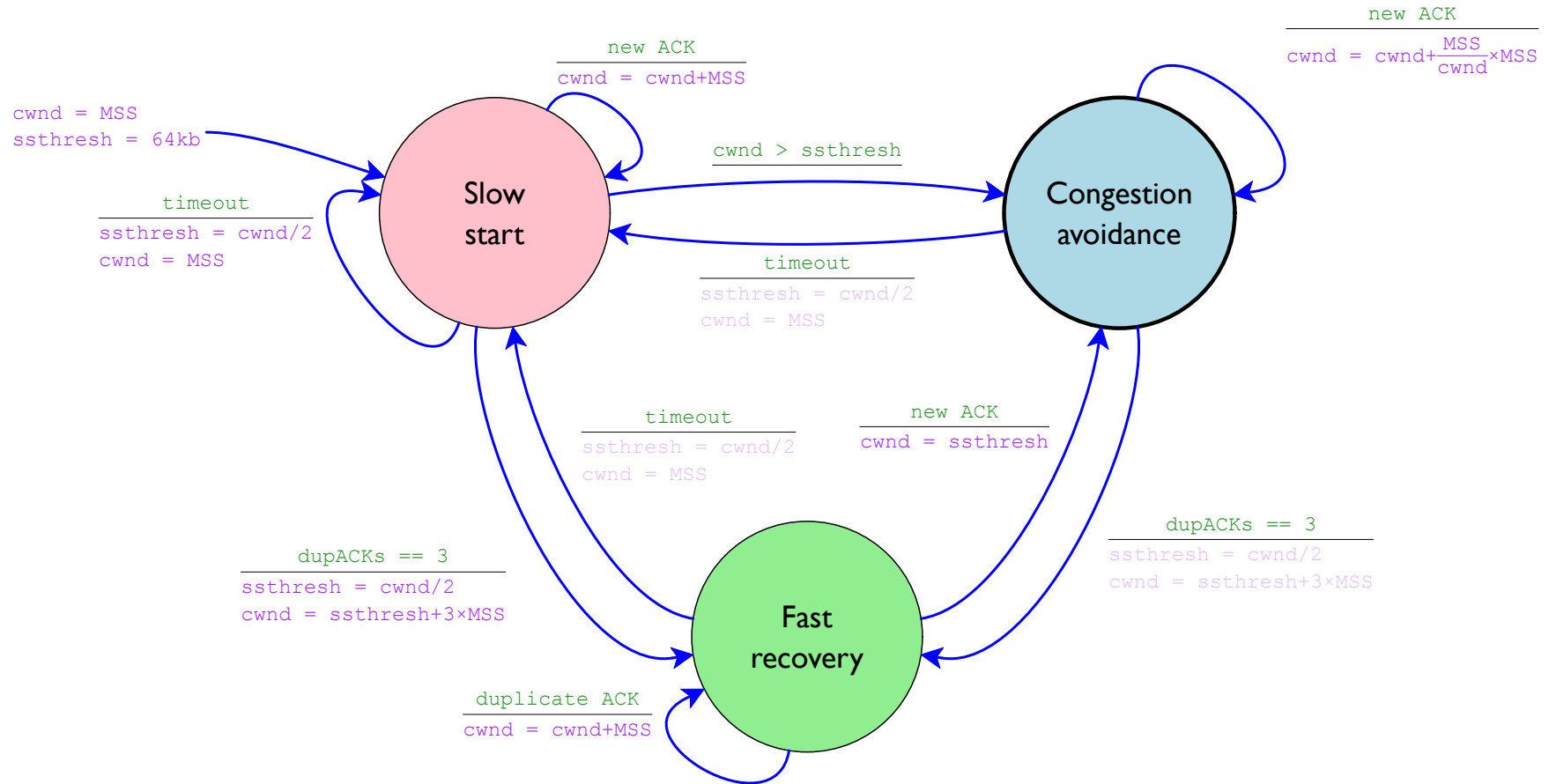
After each ACK, $cwnd += MSS$
 $\Rightarrow$ double $cwnd$ each RTT

End slow start when
 $cwnd = ssthresh$

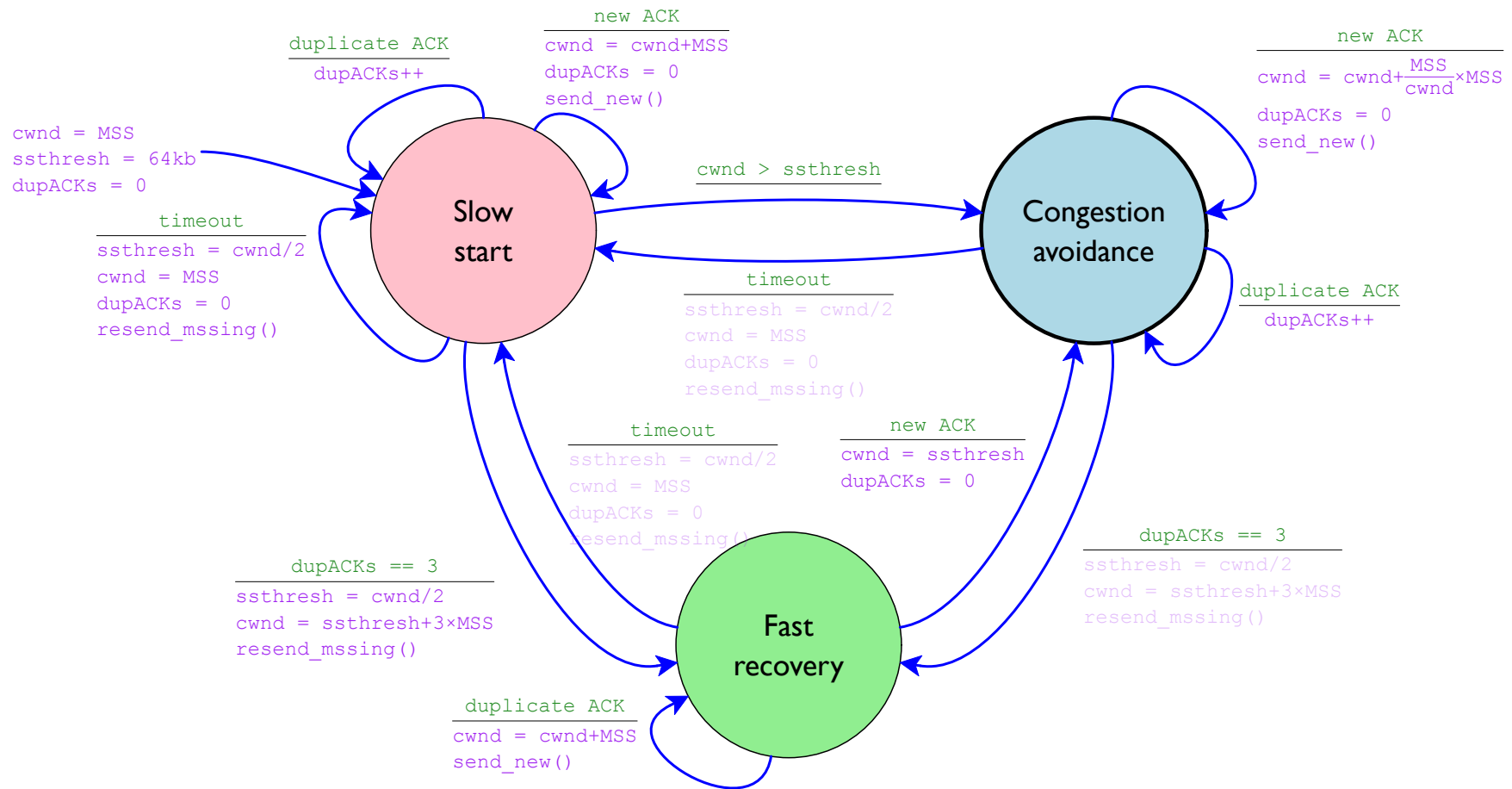
TCP Connection States



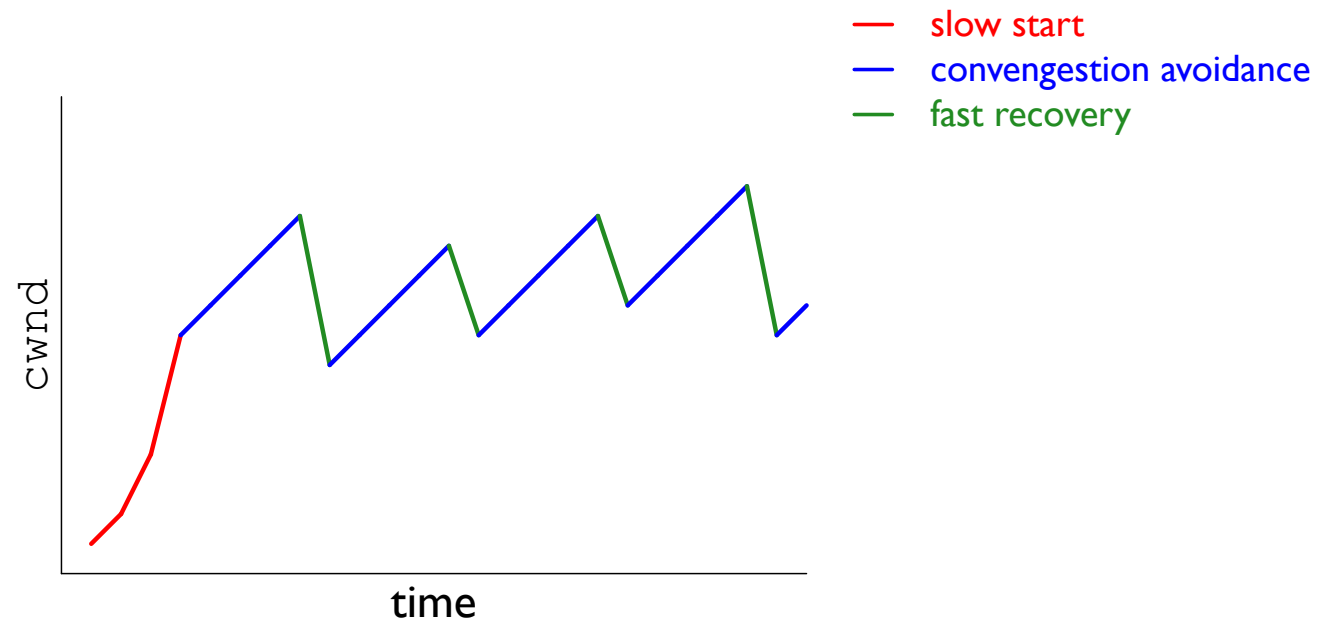
TCP Connection States



TCP Connection States

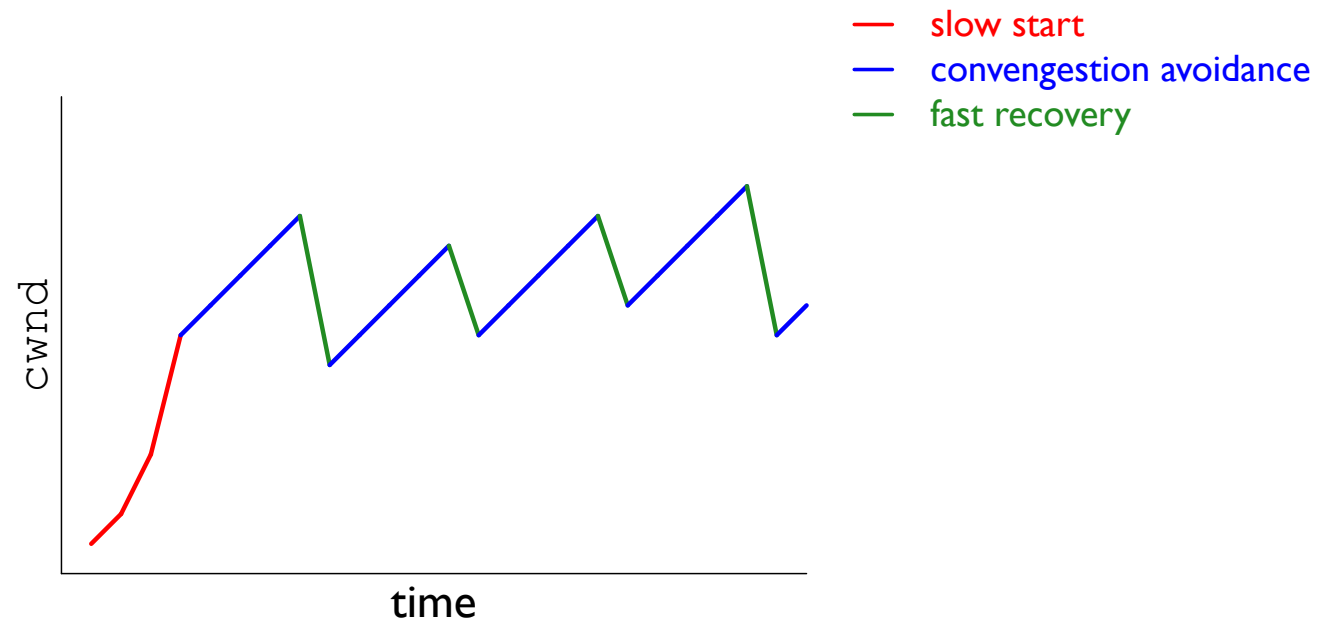


cwnd Adjustment over Time



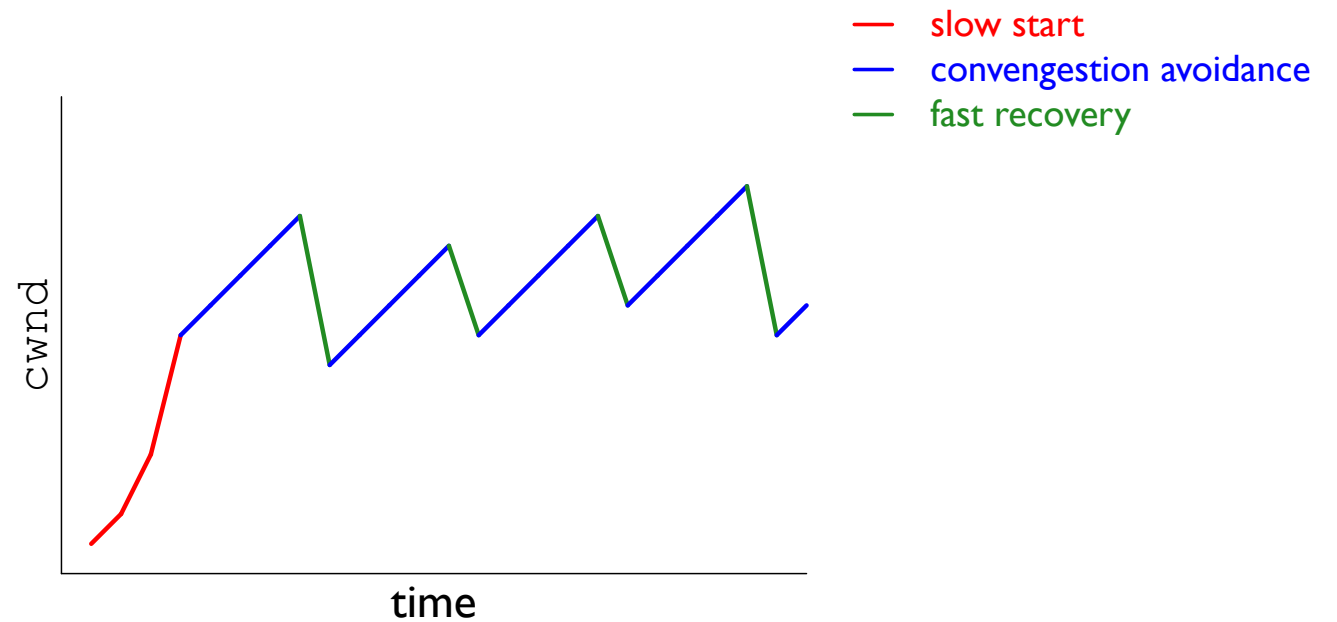
Additive increase, multiplicative decrease (AIMD)

cwnd Adjustment over Time



Expect average cwnd to be 75% of maximum

cwnd Adjustment over Time



Fair, because multiple senders tend toward same rate

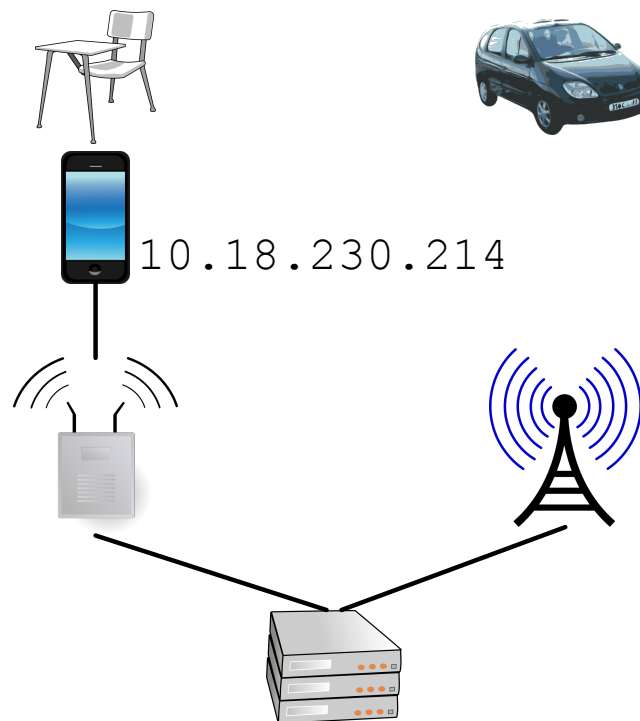
Issues with TCP

TCP is great for most purposes, but it's not perfect...

... for example, in the case of web services

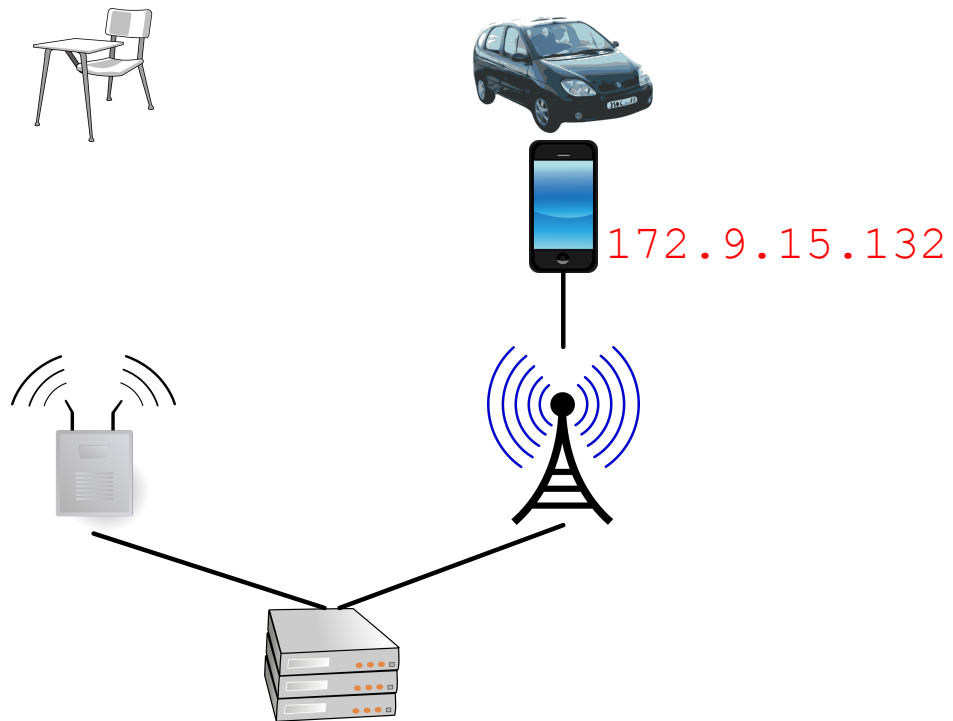
Issues with TCP

Parking-lot problem



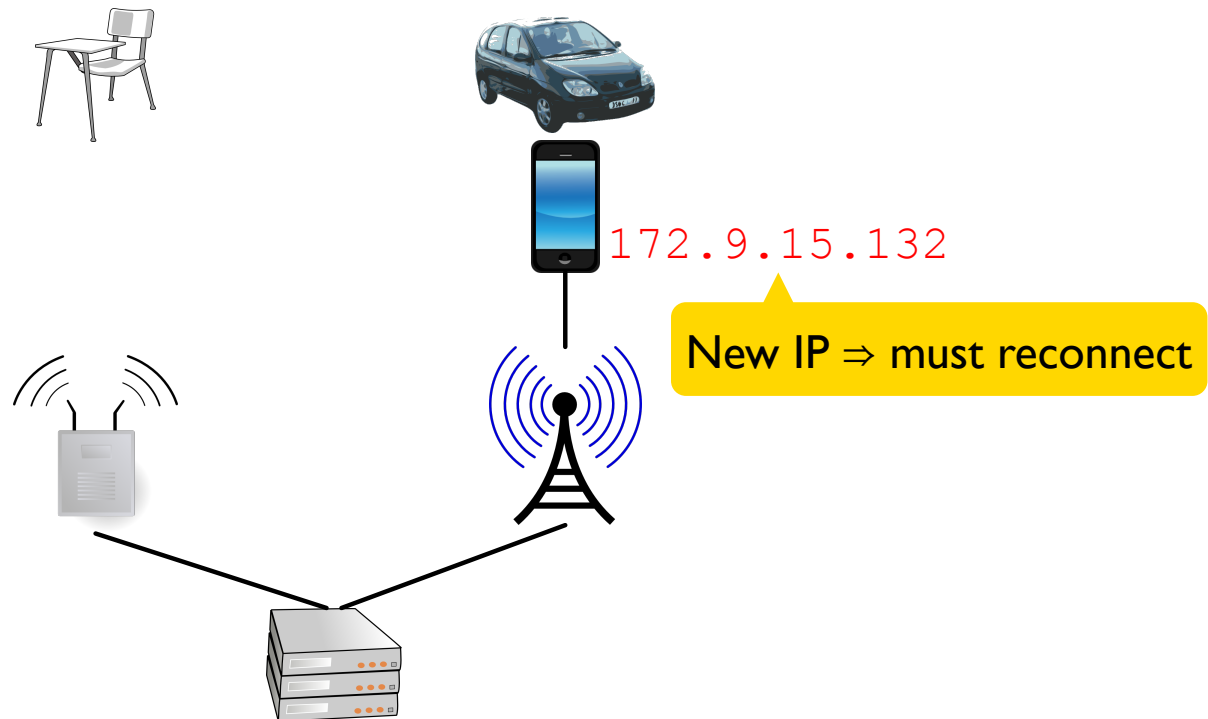
Issues with TCP

Parking-lot problem



Issues with TCP

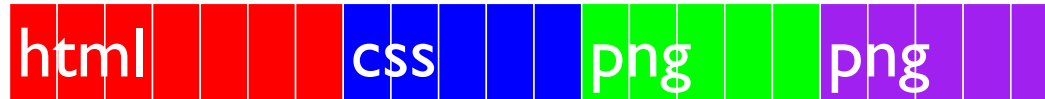
Parking-lot problem



Issues with TCP

Head-of-line problem

A typical web page needs multiple files:

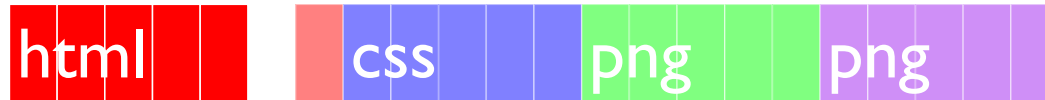


Getting parts in order can delay the whole page

Issues with TCP

Head-of-line problem

A typical web page needs multiple files:



A dropped packet delays everything further

Issues with TCP

Head-of-line problem

A typical web page needs multiple files:



HTTP/2 allows interleaving within a reply...

Issues with TCP

Head-of-line problem

A typical web page needs multiple files:



but that doesn't solve the dropped-packet problem

Issues with TCP

Head-of-line problem

A typical web page needs multiple files:

html

css

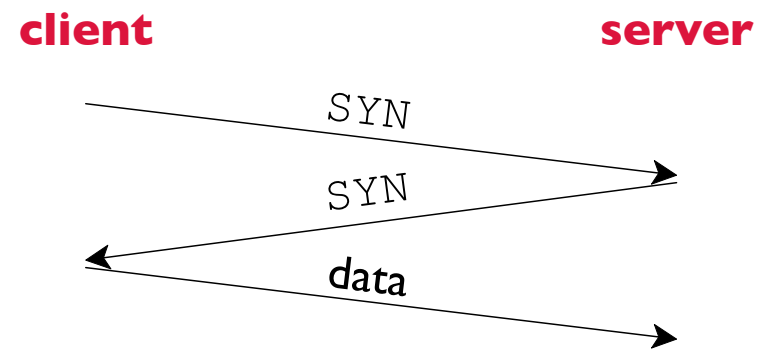
png

png

Multiple connections work, but each takes time to set up

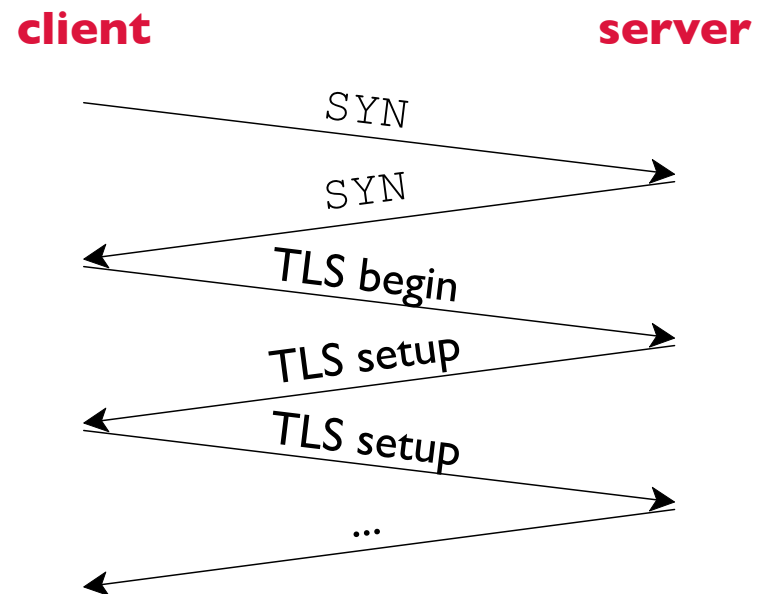
Issues with TCP

Handshake hell



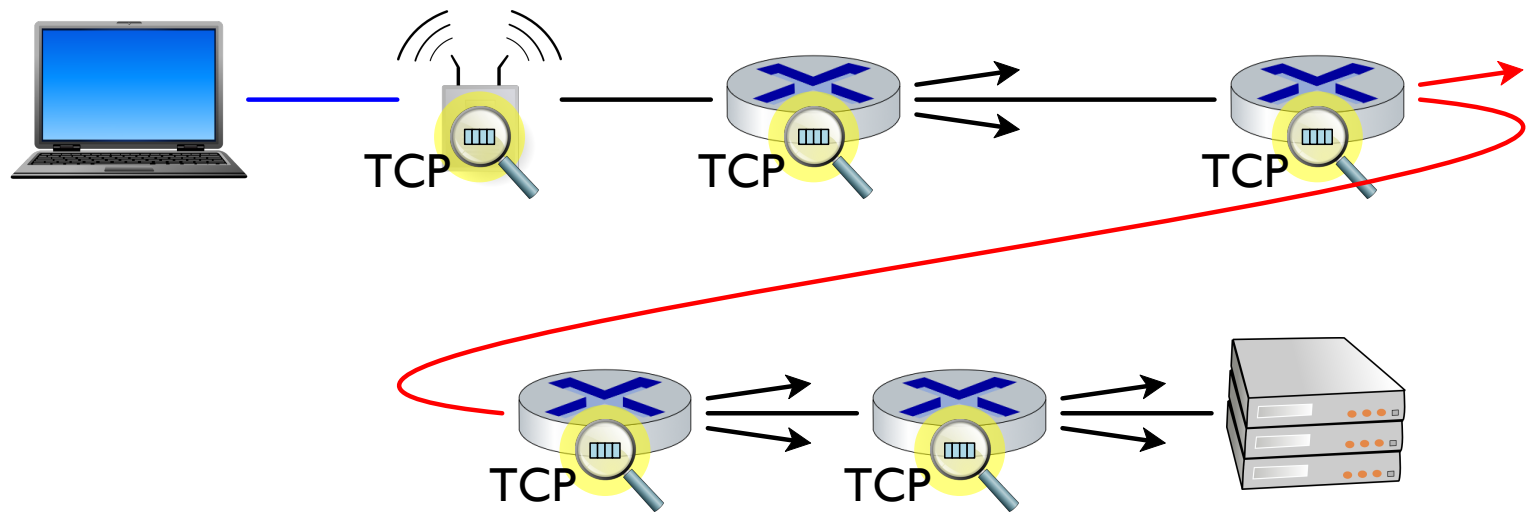
Issues with TCP

Handshake hell



Issues with TCP

Ossification



QUIC

QUIC: Quick UDP Internet Connections

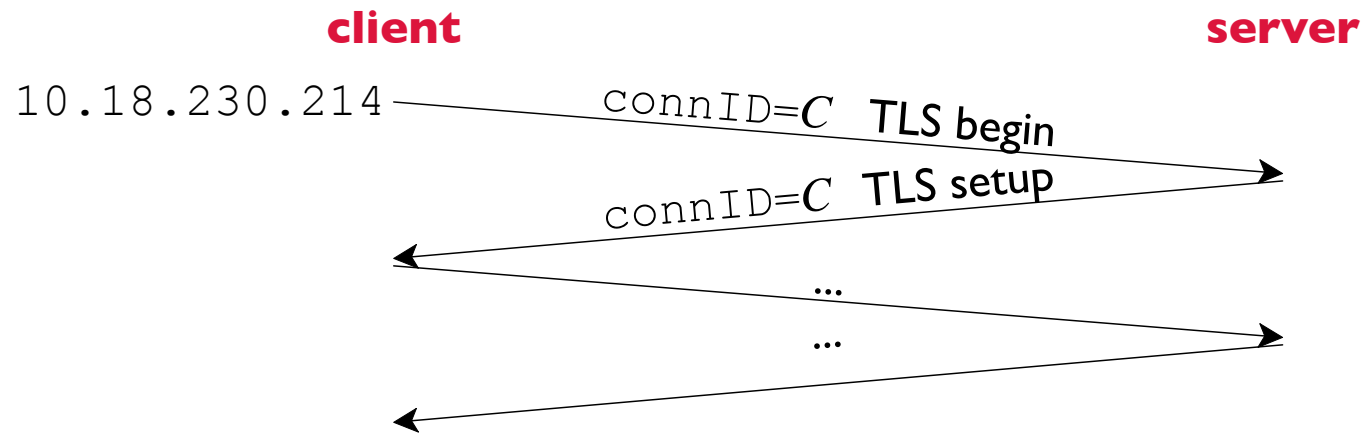
- implemented in Chrome in 2012
- standardized in 2021 as RFC 9000
- implemented in major browsers

QUIC

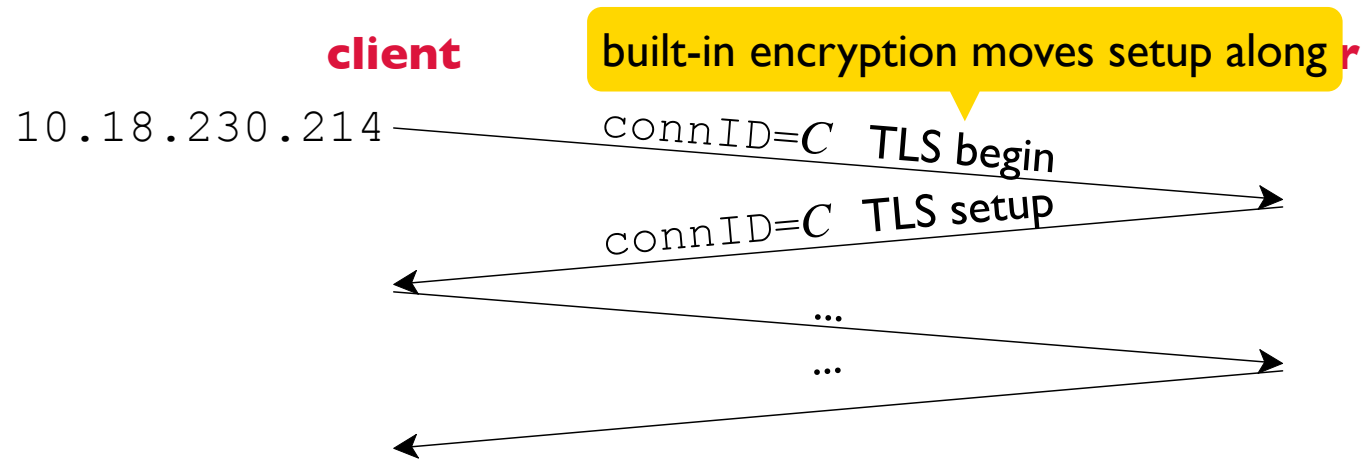
QUIC: Quick UDP Internet Connections

- builds on UDP
- connection-oriented based on a connection ID
not host and port
- built-in encryption, covers more headers
- can interleave files without dropped-packet interactions

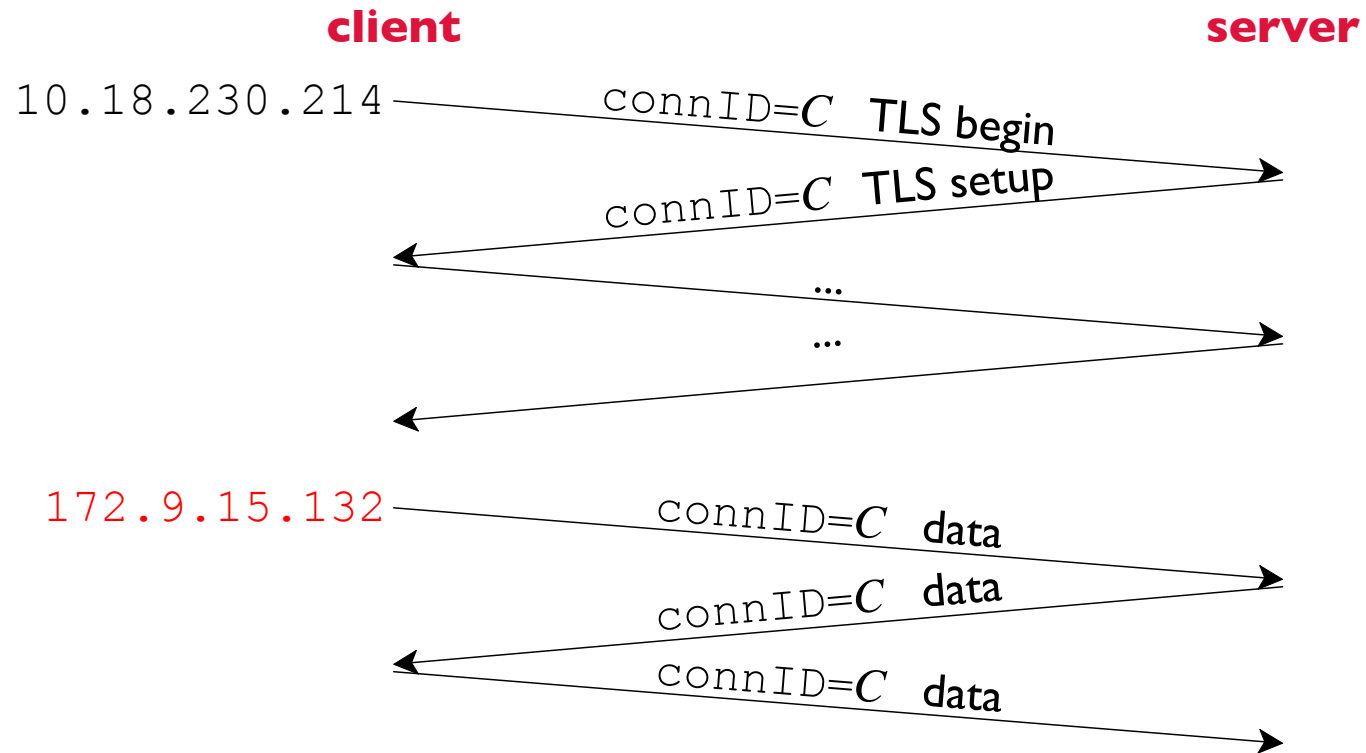
QUIC



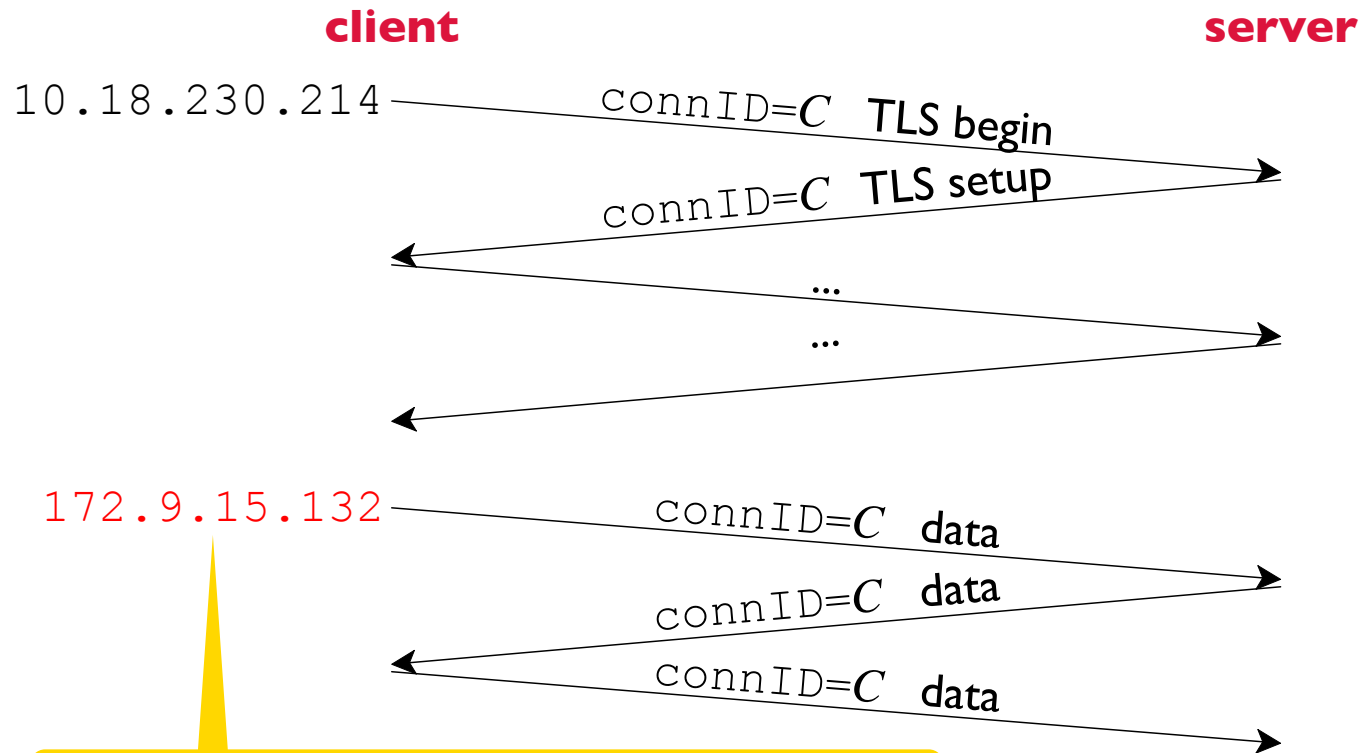
QUIC



QUIC



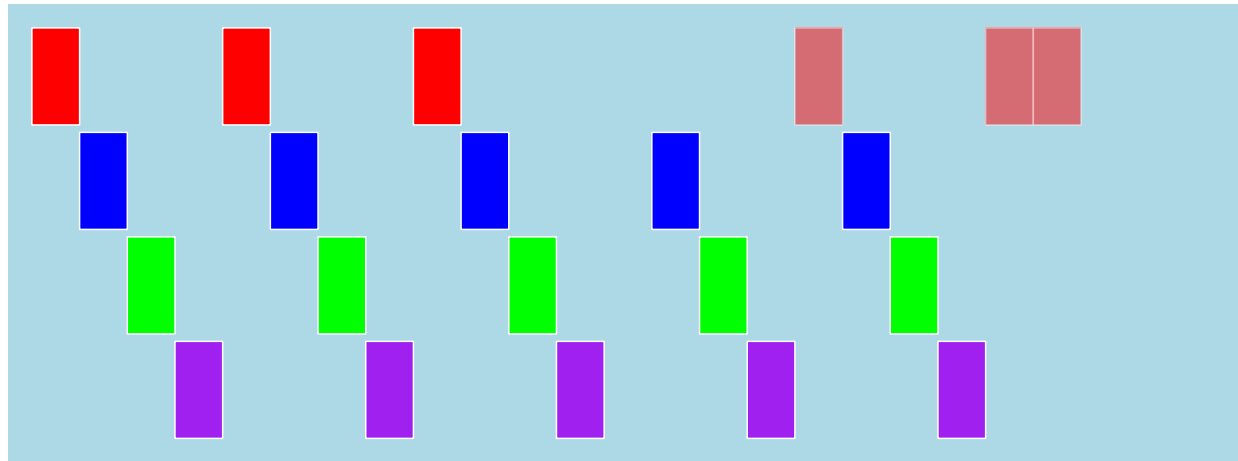
QUIC



Connection ID allows continue without setup from a new client address

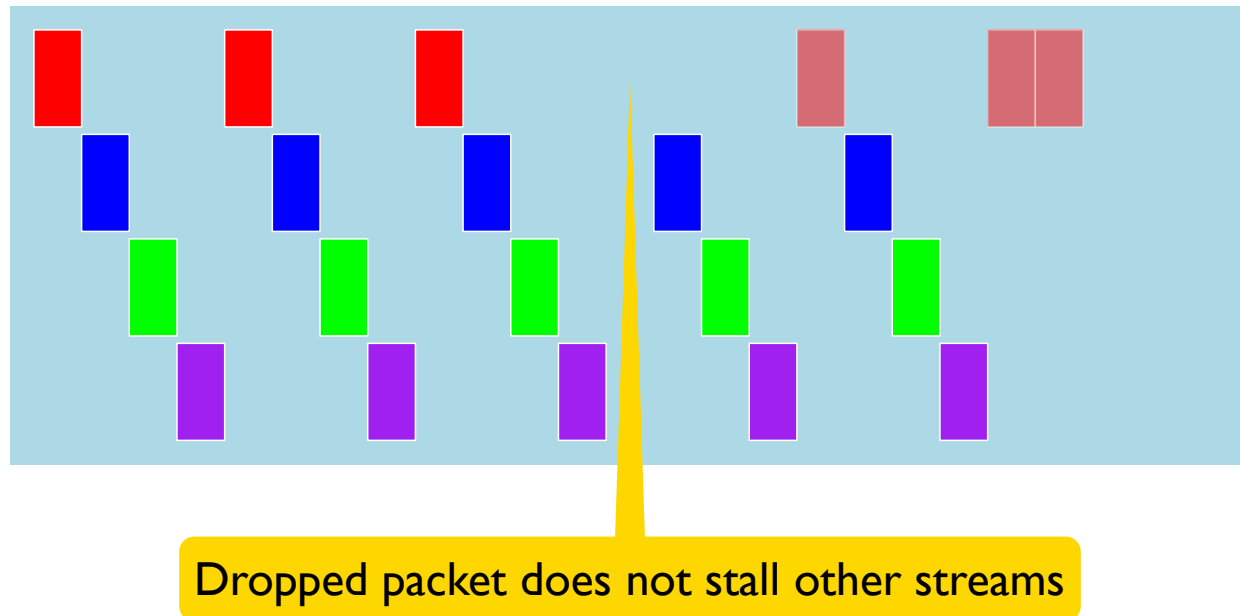
QUIC

Concurrent streams are handled at the packet level within a connection



QUIC

Concurrent streams are handled at the packet level within a connection



Summary

Flow control avoids overwhelming the other end of a connection

Congestion control avoids overwhelming the network as a whole

TCP states for congestion control:

- **slow start**
- **congestion avoidance**
- **fast recovery**

TCP isn't always the best solution, it can suffer from

handshake hell,

the **parking-lot problem**, and

the **head-of-line problem**.