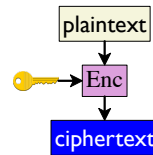


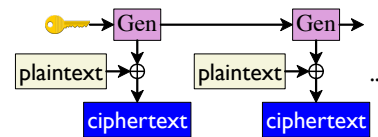
Cryptography Toolbox

block cipher



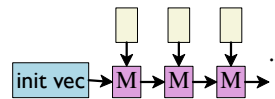
fixed-size

stream cipher



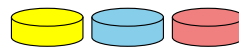
any length

hash function



$$H(x) = H(y) \\ \Rightarrow \\ x = y$$

Diffie-Hellman



shared secret



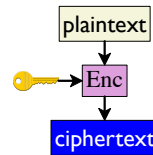
RSA



authentication

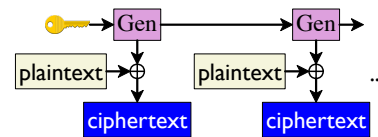
Cryptography Toolbox

block cipher



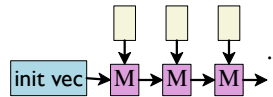
fixed-size

stream cipher



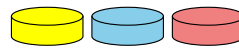
any length

hash function



$$H(x) = H(y) \\ \Rightarrow \\ x = y$$

Diffie-Hellman



shared secret



RSA

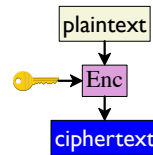


authentication

Cryptographic algorithms need to mix these ingredients correctly

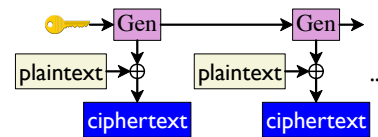
Cryptography Toolbox

block cipher



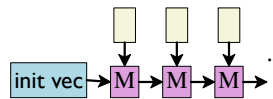
fixed-size

stream cipher



any length

hash function



$$H(x) = H(y) \\ \Rightarrow \\ x = y$$

Diffie-Hellman



shared secret



RSA



authentication

Cryptographic algorithms need to mix these ingredients correctly

... including for authentication

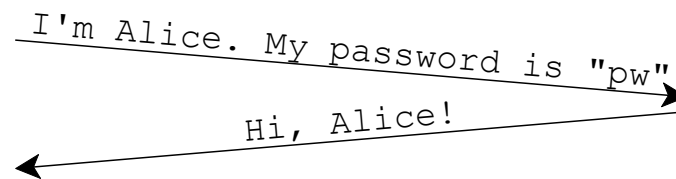
Sending a Password



Alice

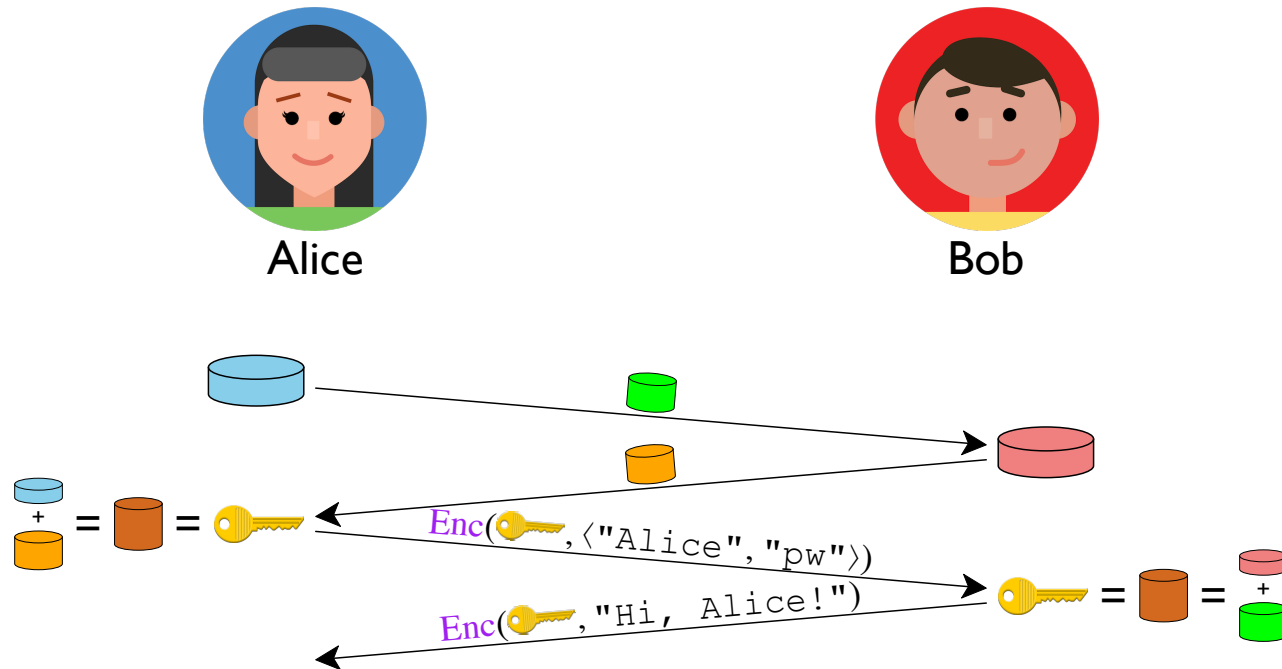


Bob



Obviously bad to send a password as plaintext

Sending a Password



Use Diffie-Hellman to get a shared secret key?

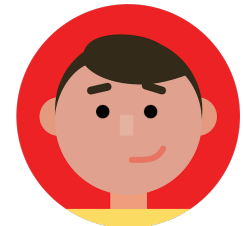
Sending a Password



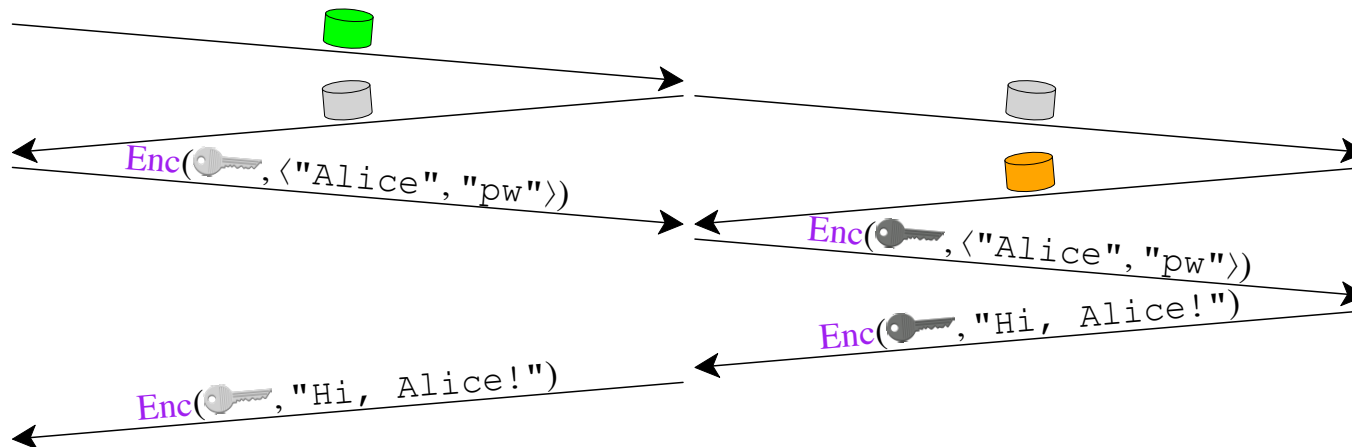
Alice



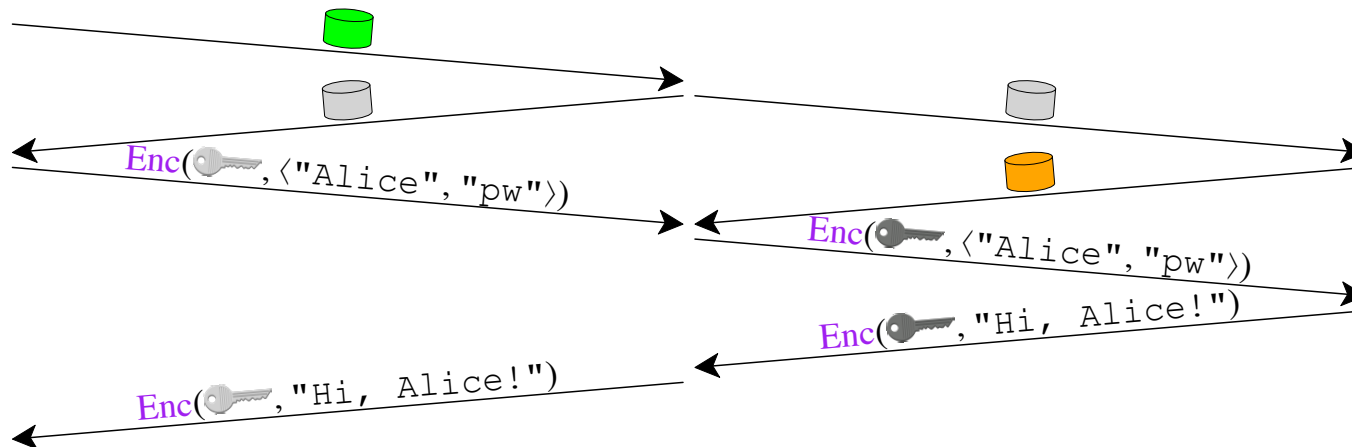
Mallory



Bob



Sending a Password



Don't send a password without authenticating first!

Authentication

Authentication requires some prior arrangement

Some options:

- previously shared secret 
- previously acquired RSA  and 

For now, assume a previously shared 

Some Bad Authentication Ideas



Alice



Bob



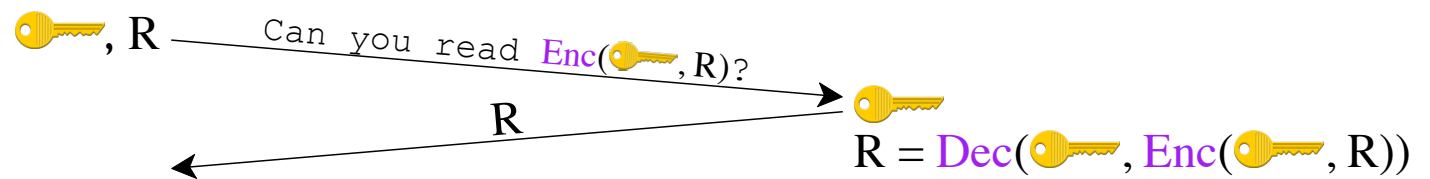
Some Bad Authentication Ideas



Alice



Bob



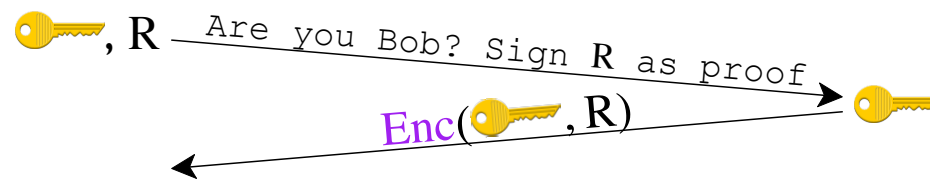
Some Bad Authentication Ideas



Alice



Bob



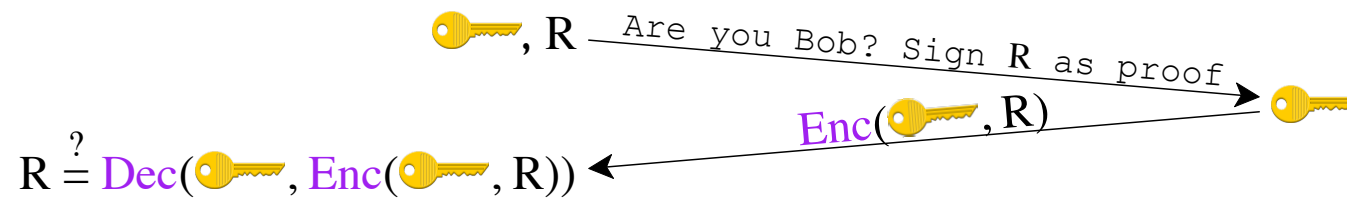
Some Bad Authentication Ideas



Alice



Bob



Some Bad Authentication Ideas



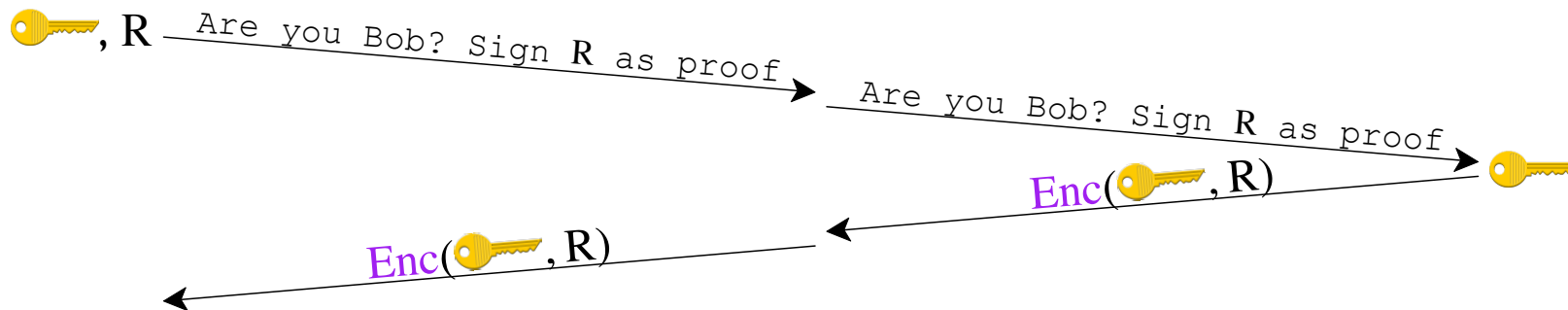
Alice



Mallory



Bob



Whether this is immediately bad depends on whether the rest of the conversation encrypts with 

Some Bad Authentication Ideas



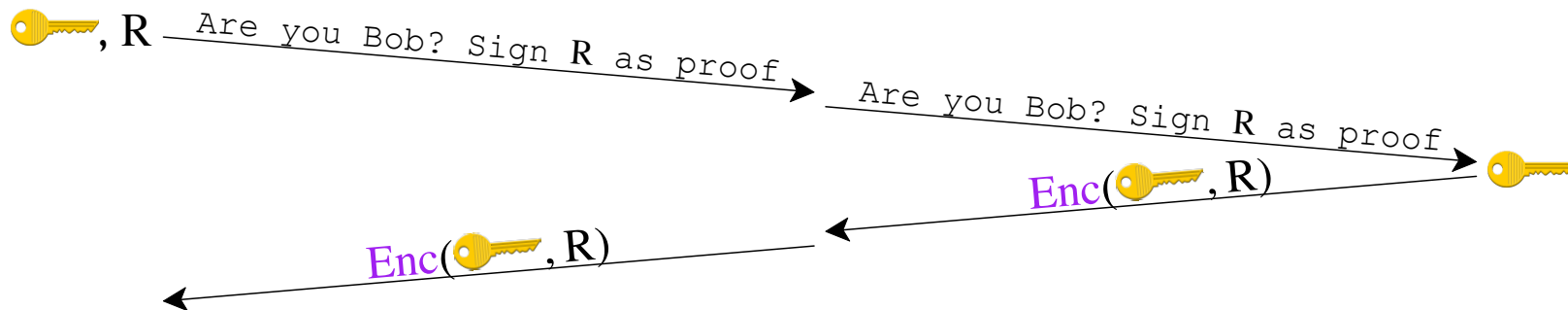
Alice



Mallory



Bob



More generally, the problem here is that Bob has blindly signed a value handed to him

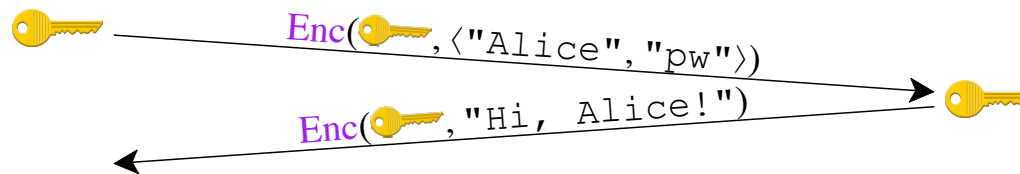
Some Bad Authentication Ideas



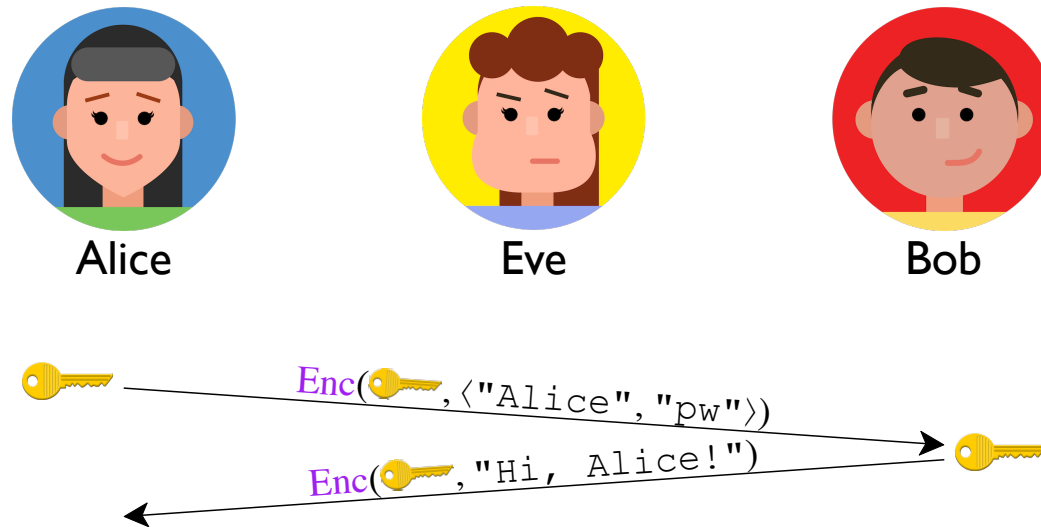
Alice



Bob

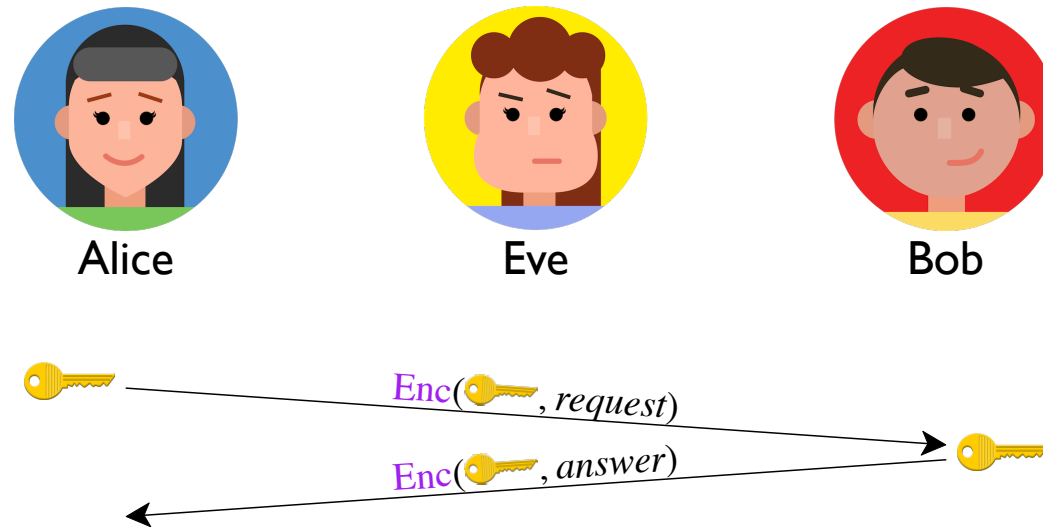


Some Bad Authentication Ideas



Replay attack: Eve can record $Enc(\text{key}, \langle \text{"Alice"}, \text{"pw"} \rangle)$ and play it back later

Some Bad Authentication Ideas



Anyway, no need for a separate password if you have , but beware of the same problem with other predictable messages

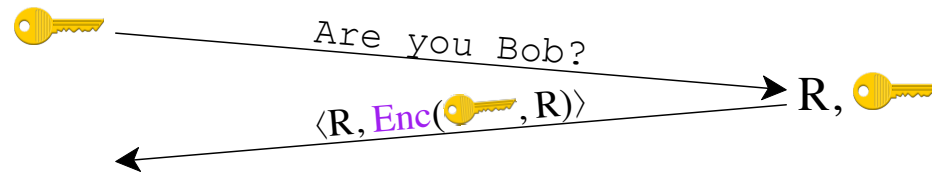
Authentication



Alice



Bob



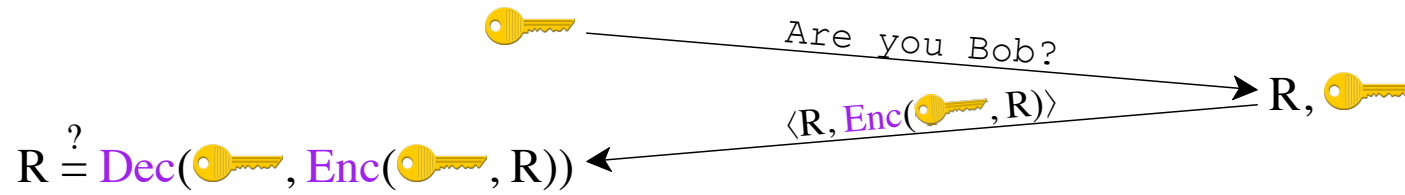
Authentication



Alice



Bob



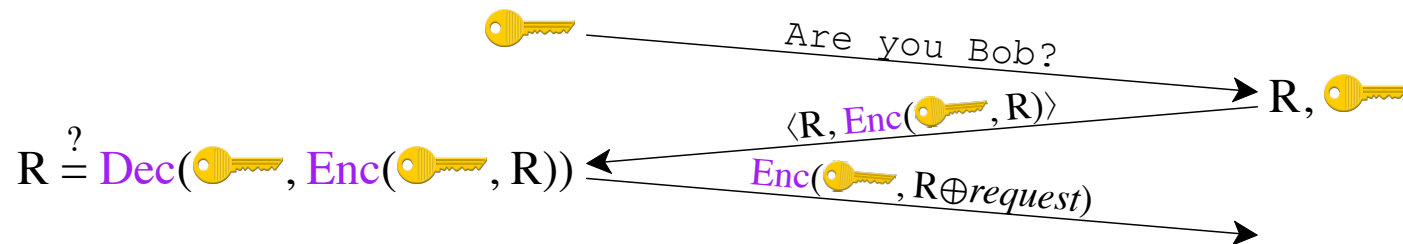
Authentication



Alice



Bob



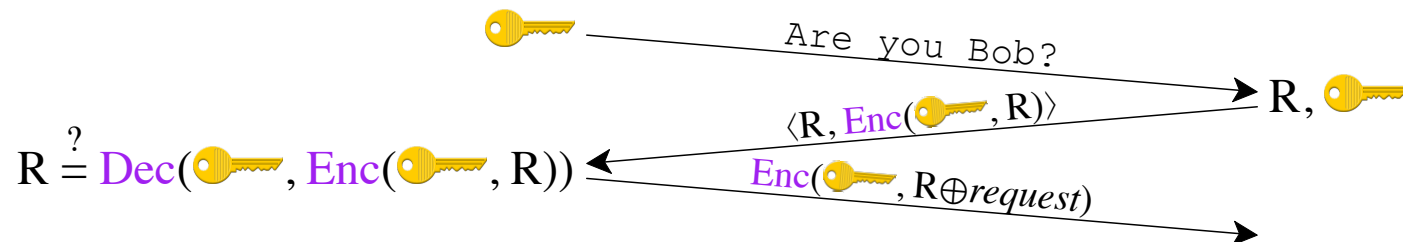
Authentication



Alice



Bob



As long as Bob picks a unique random R each time, this will work ok

... but there is still some room for improvement

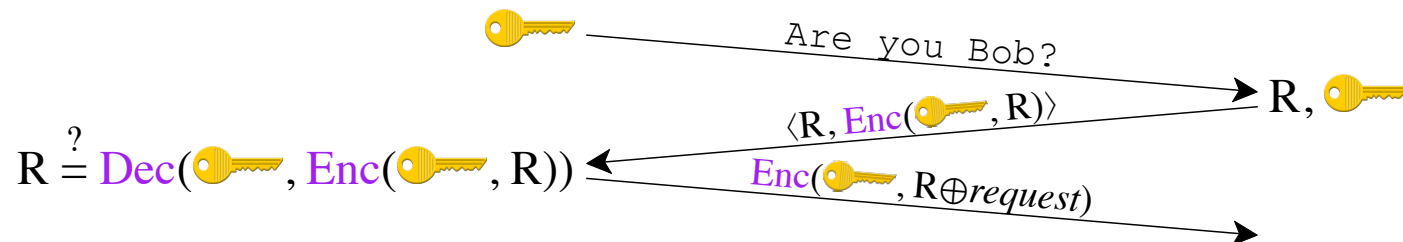
Authentication



Alice



Bob



Gives attackers lots of examples for key

As long as Bob picks a unique random R each time this will work ok

... but there is still some room for improvement

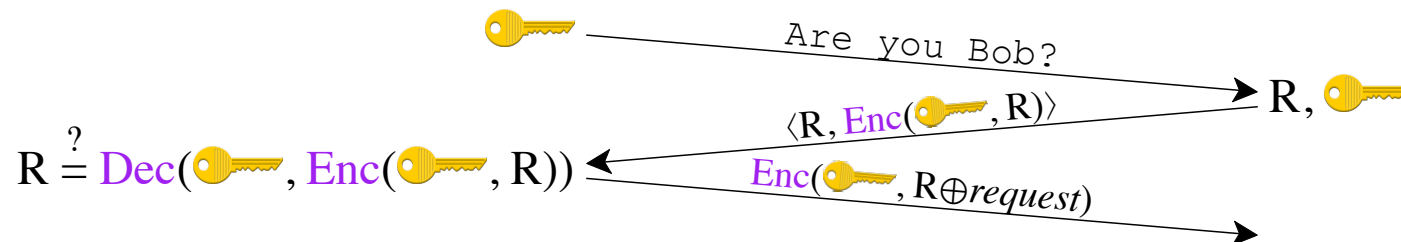
Authentication



Alice



Bob



If  is compromised, all history is readable

As long as Bob picks a unique random R each time this will work ok

... but there is still some room for improvement

Distributing Public Keys

Suppose that Alice acquires Bob's public key



Is it *really* Bob's key?



Alice



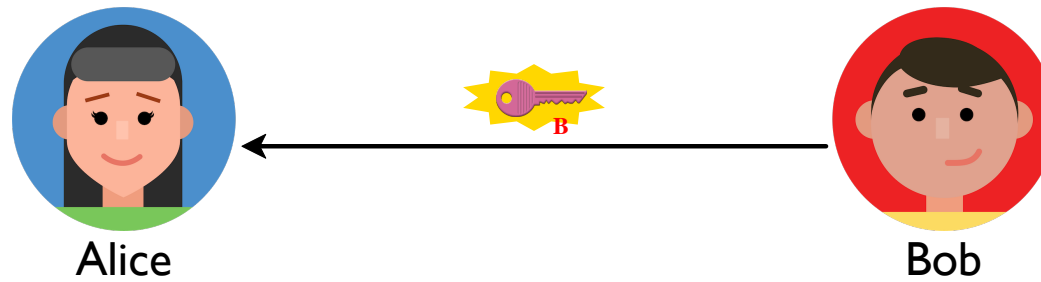
Mallory



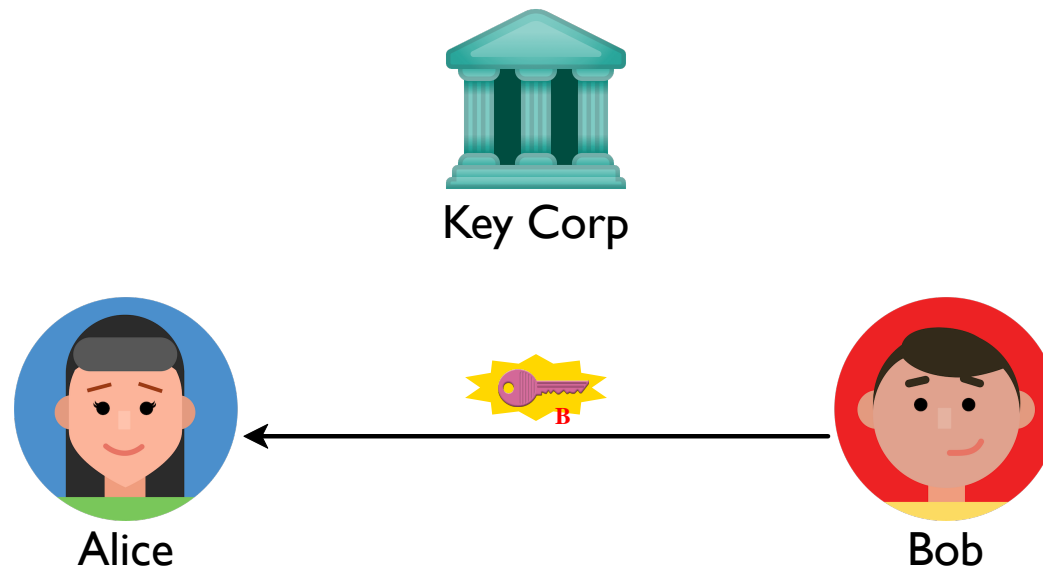
Bob



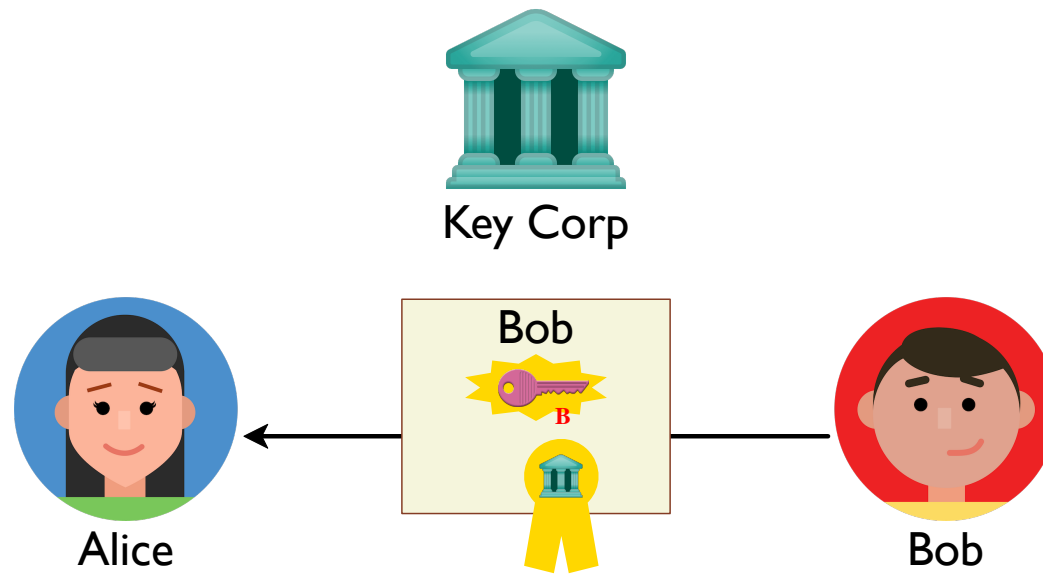
Distributing Public Keys



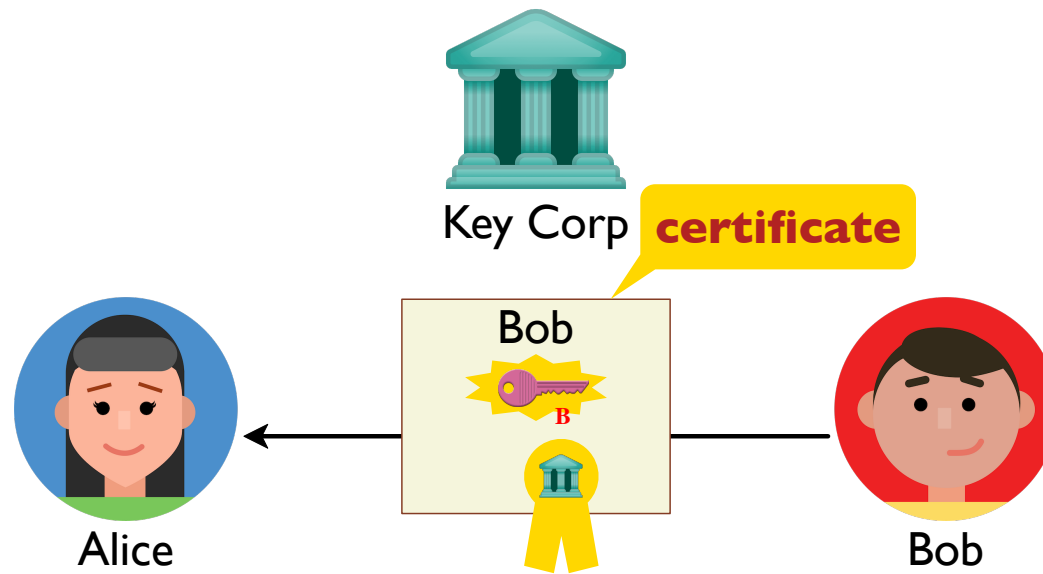
Distributing Public Keys



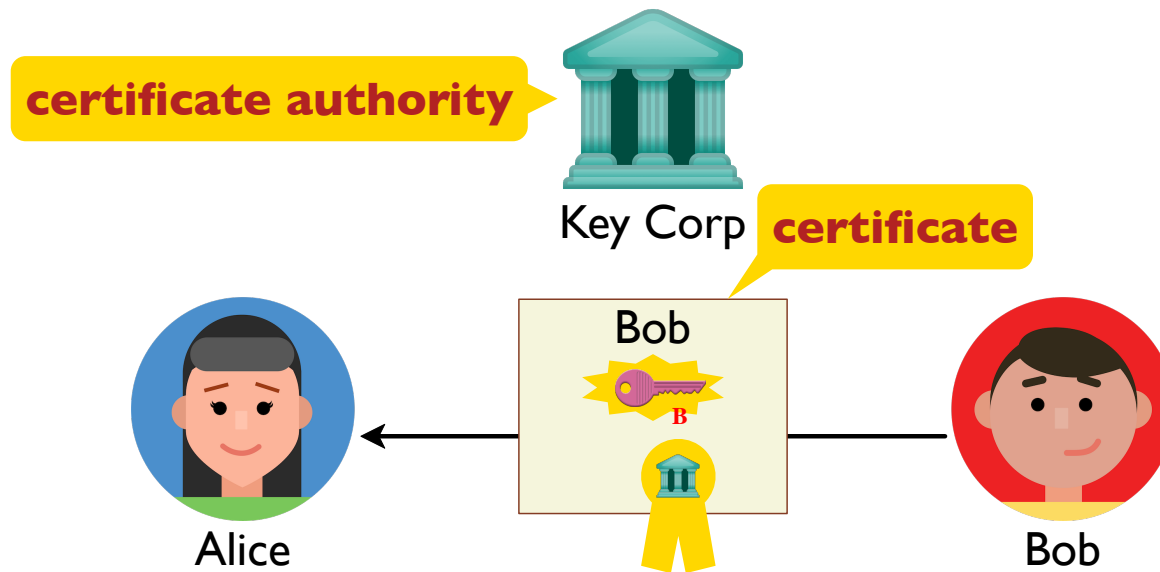
Distributing Public Keys



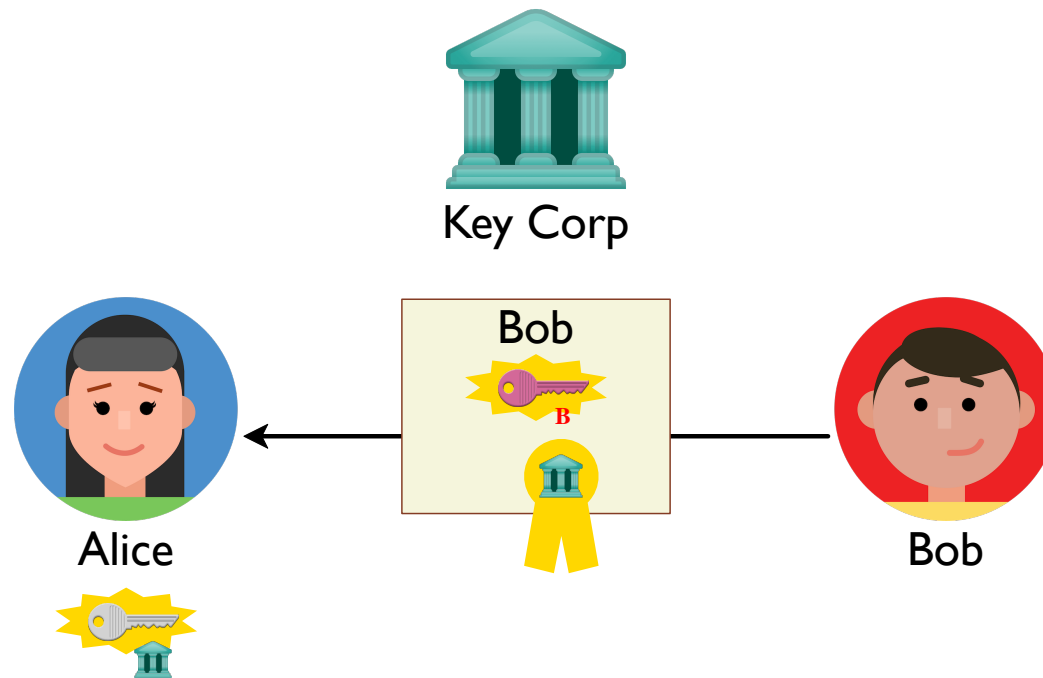
Distributing Public Keys



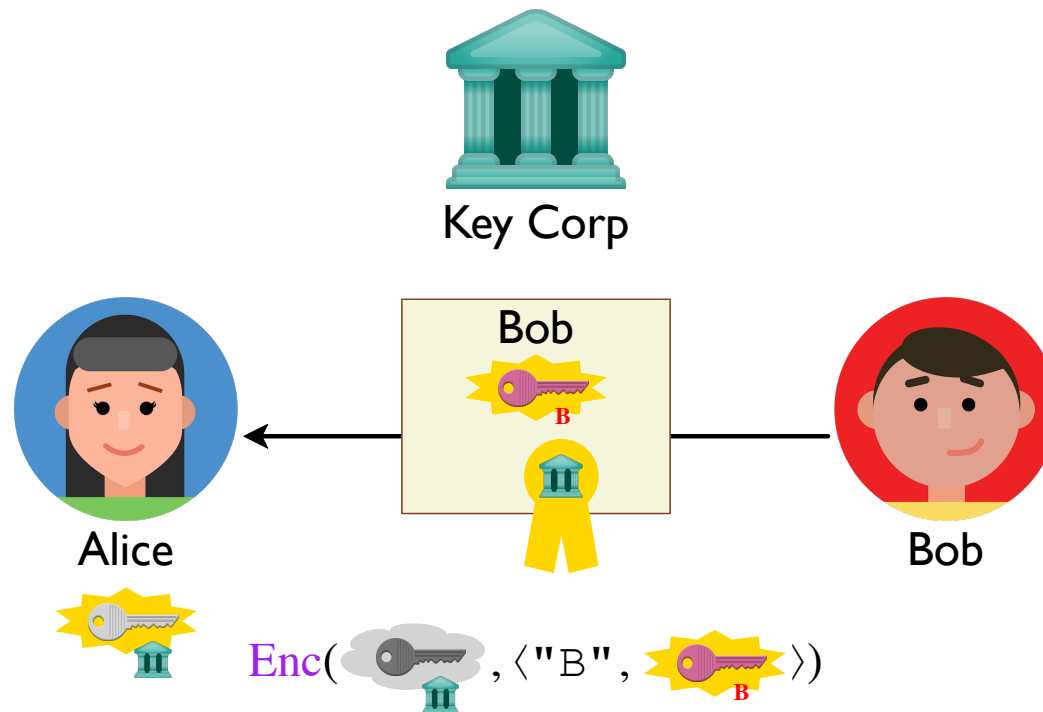
Distributing Public Keys



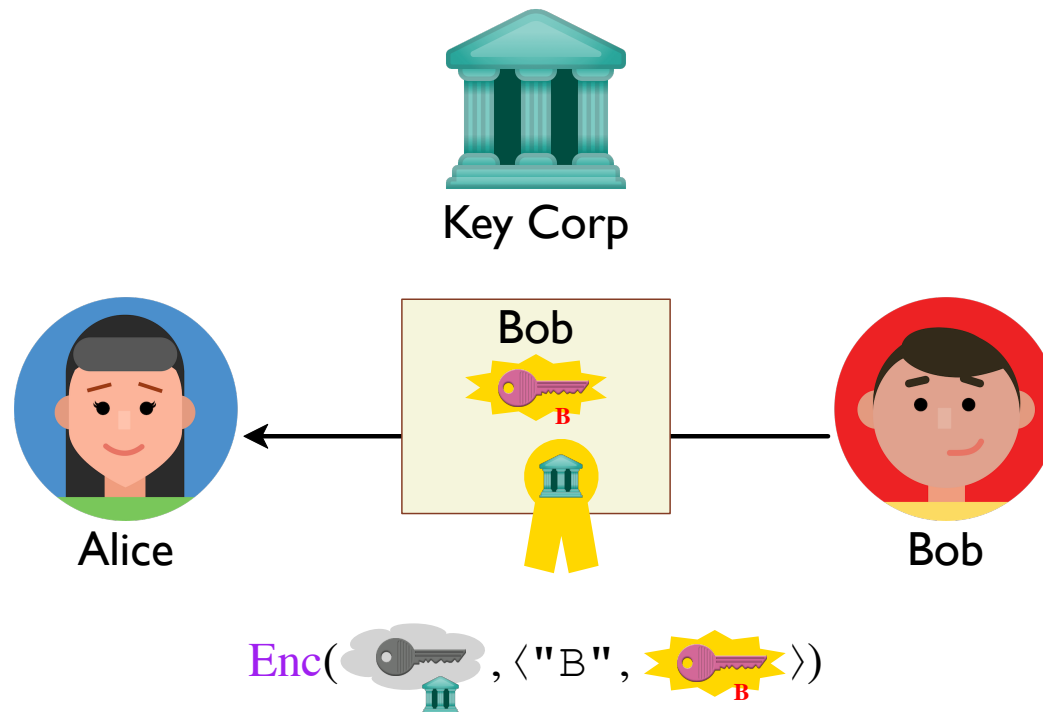
Distributing Public Keys



Distributing Public Keys



Distributing Public Keys



Distributing Public Keys



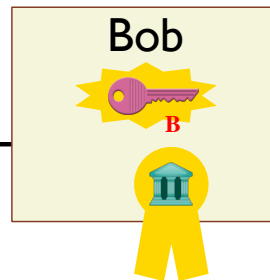
Meta Corp



Key Corp



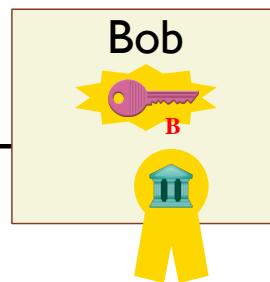
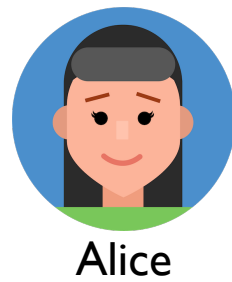
Alice



Bob

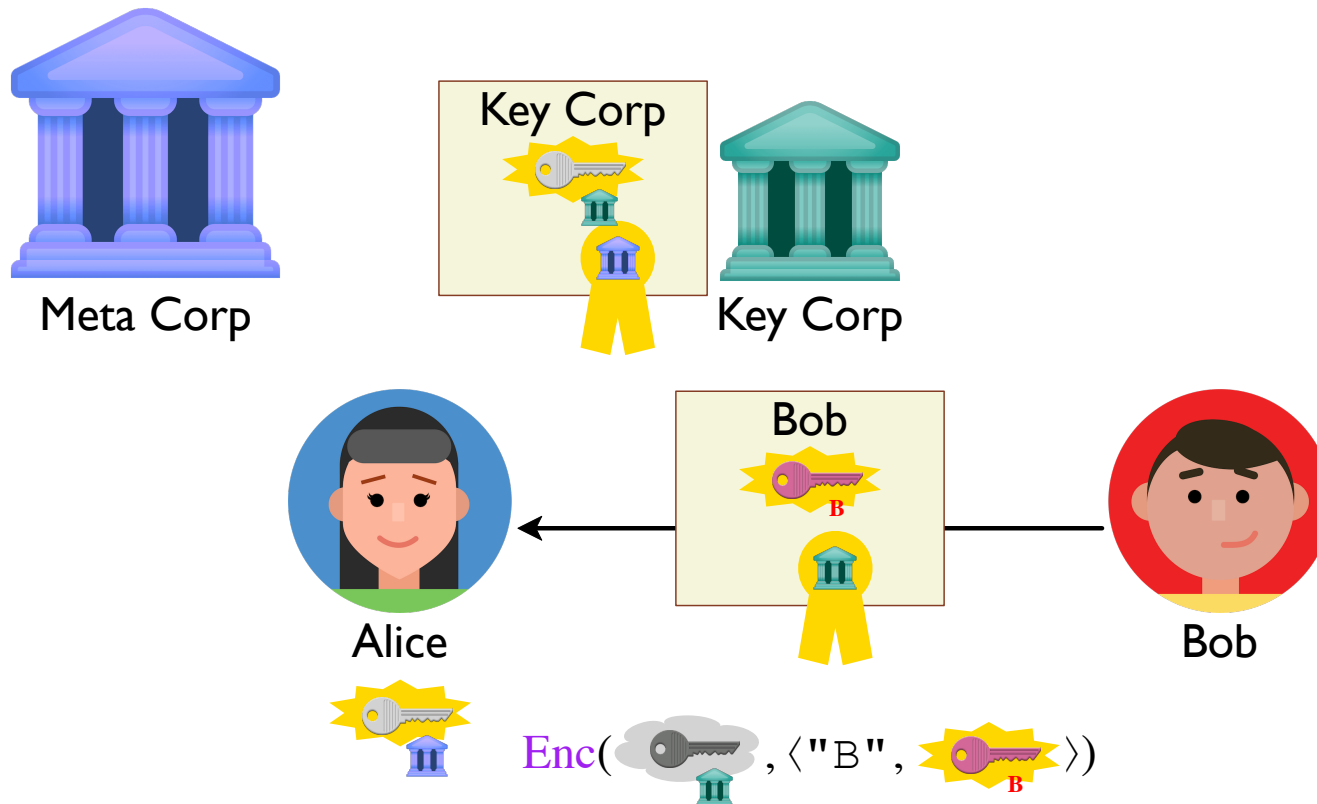
$Enc(\text{key icon}, \langle "B", \text{key icon} \rangle)$

Distributing Public Keys

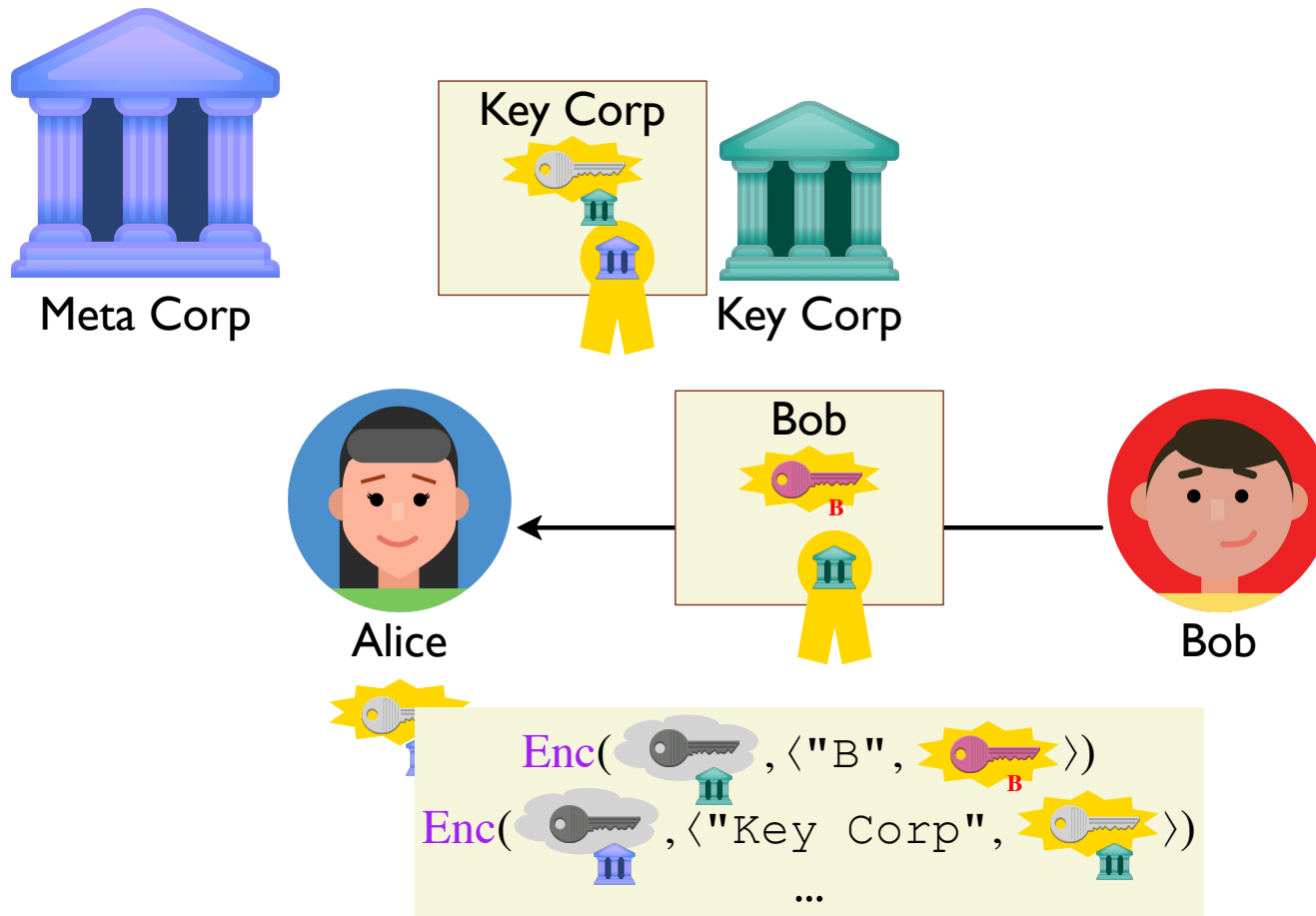


$Enc(\text{key}, \langle "B", \text{key} \rangle)$

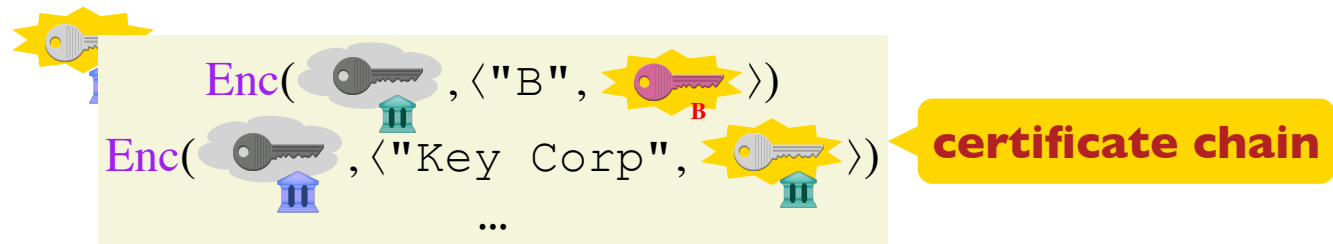
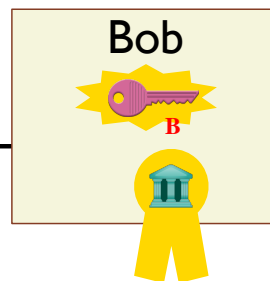
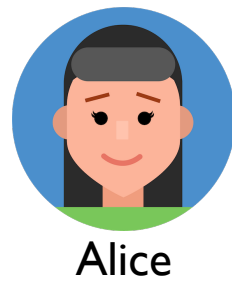
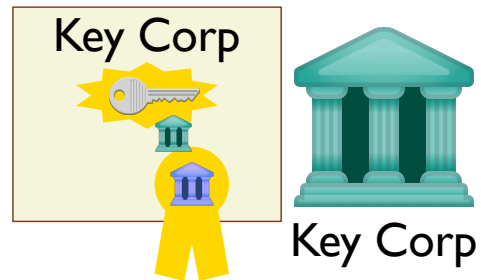
Distributing Public Keys



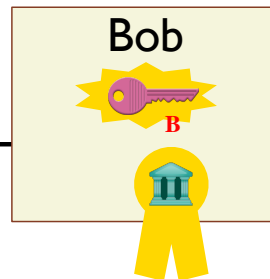
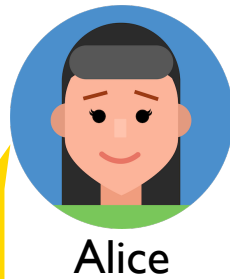
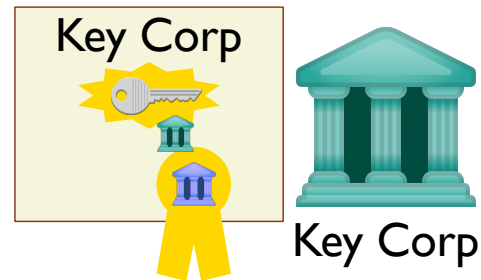
Distributing Public Keys



Distributing Public Keys



Distributing Public Keys

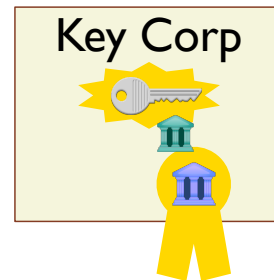


Alice's operating system includes **root certificates**

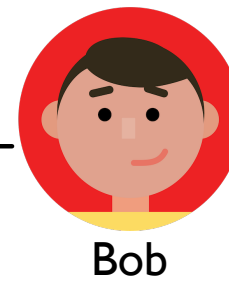
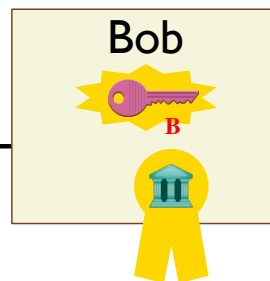
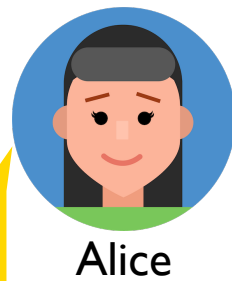
$\text{Enc}(\text{key}, \langle \text{"B"}, \text{key}_B \rangle)$
 $\text{Enc}(\text{key}, \langle \text{"Key Corp"}, \text{key}_{\text{Key Corp}} \rangle)$
...

certificate chain

Distributing Public Keys



Try
security dump-keychain
/System/Library/Keychains/SystemRootCertificates.keychain



Alice's operating system includes **root certificates**

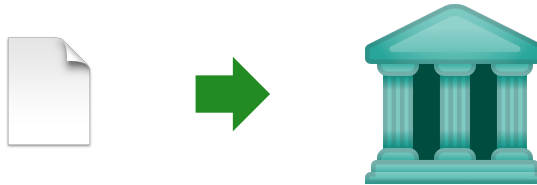
$\text{Enc}(\text{key}, \langle \text{"B"}, \text{key}_B \rangle)$
 $\text{Enc}(\text{key}, \langle \text{"Key Corp"}, \text{key}_{\text{Key Corp}} \rangle)$
...

certificate chain

Getting a Certificate

How does a service provider obtain a certificate?

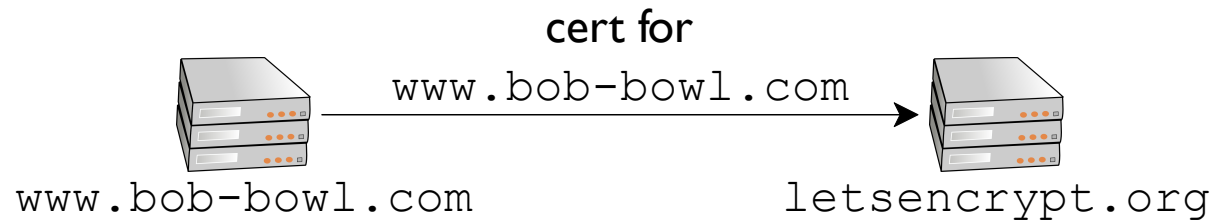
- Some kinds of certificates: **paperwork**



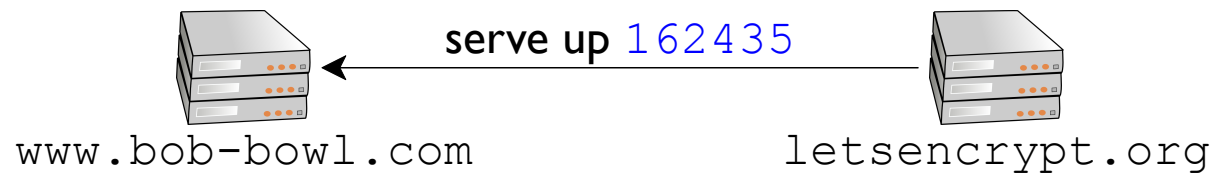
- Web sites: **Let's Encrypt**



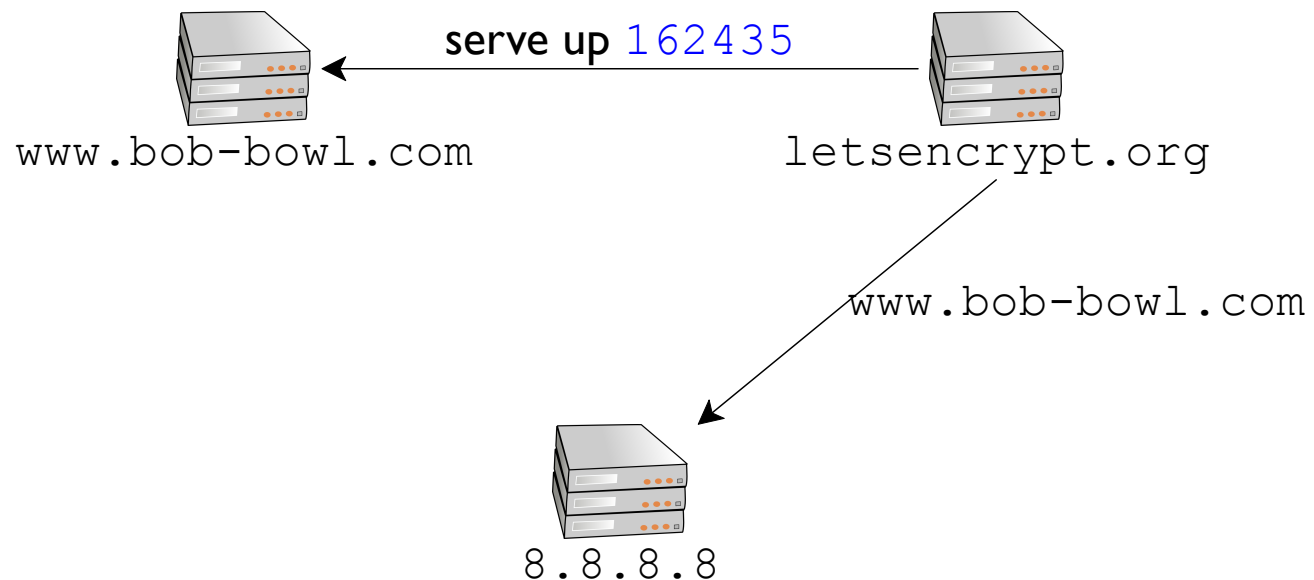
Getting a Certificate via Let's Encrypt



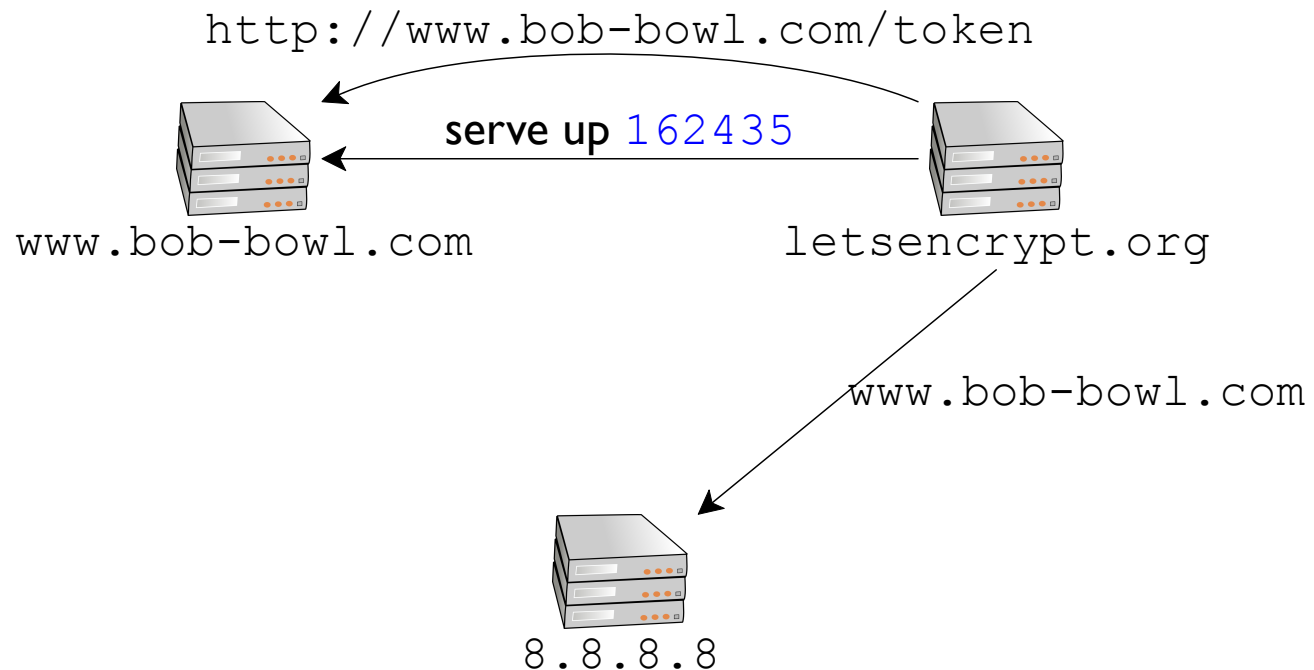
Getting a Certificate via Let's Encrypt



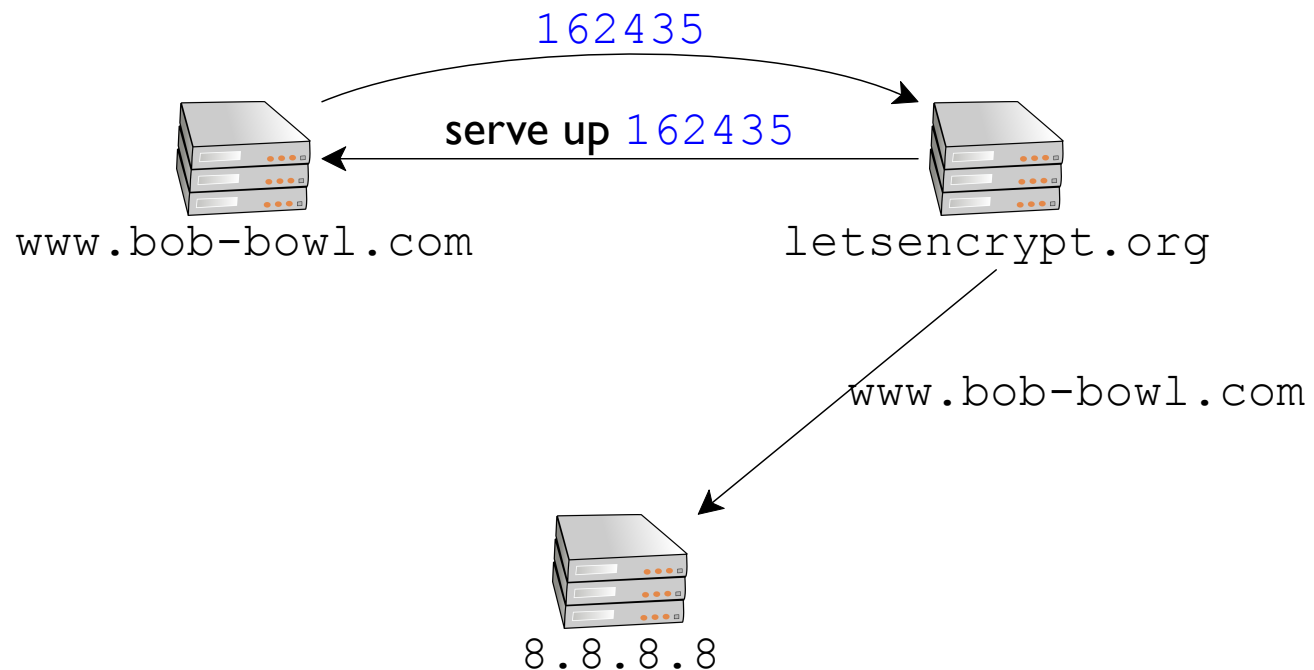
Getting a Certificate via Let's Encrypt



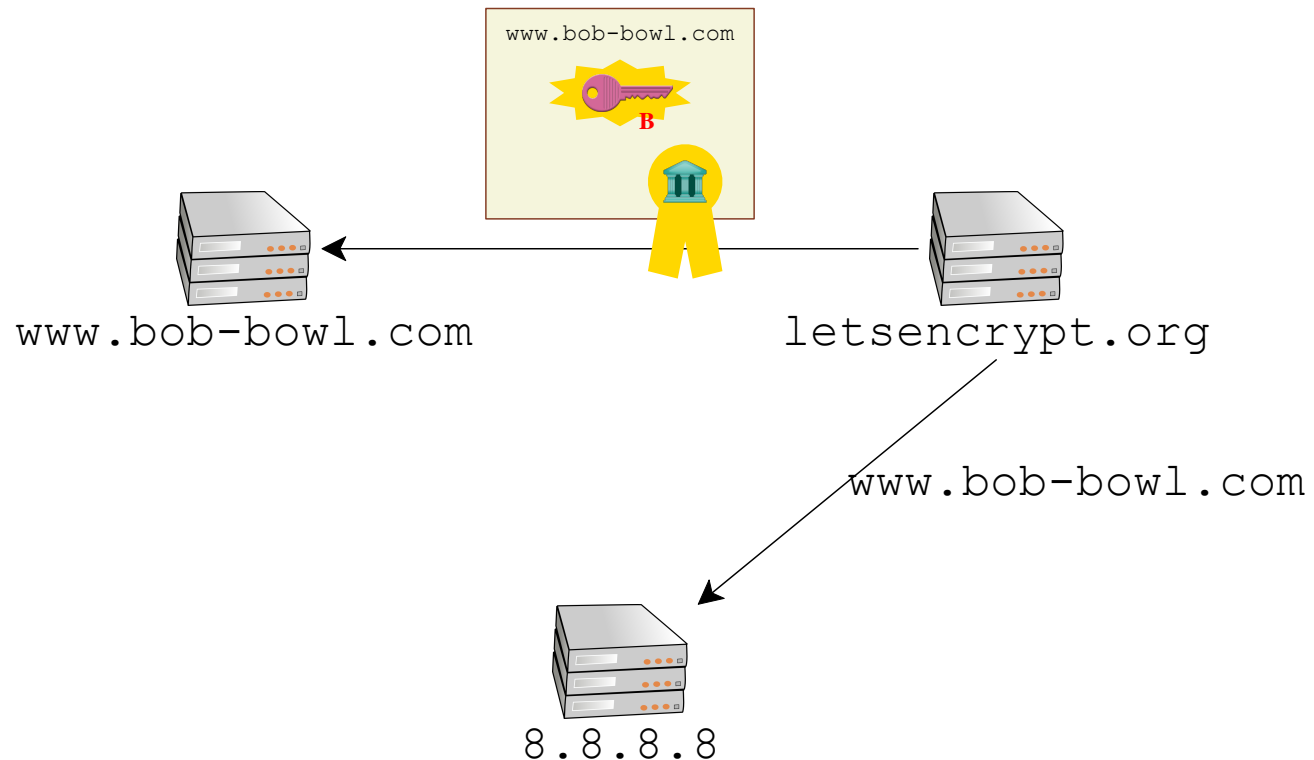
Getting a Certificate via Let's Encrypt



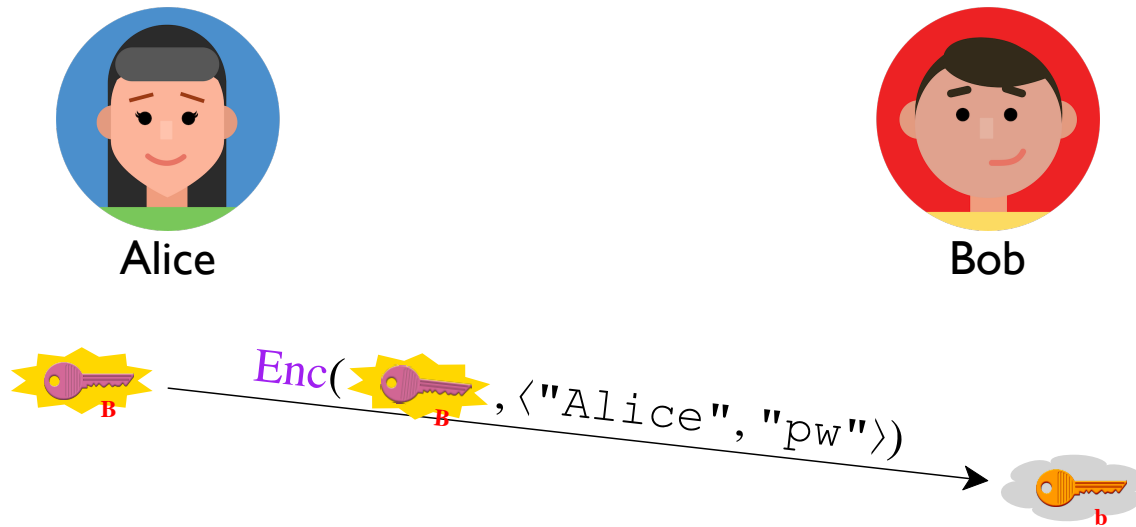
Getting a Certificate via Let's Encrypt



Getting a Certificate via Let's Encrypt

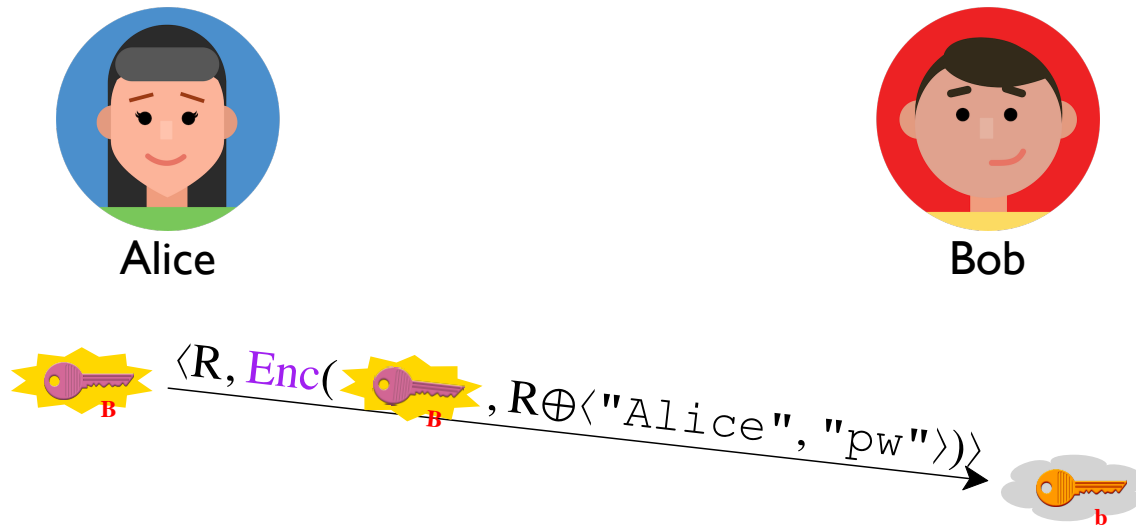


More Bad Authentication Ideas



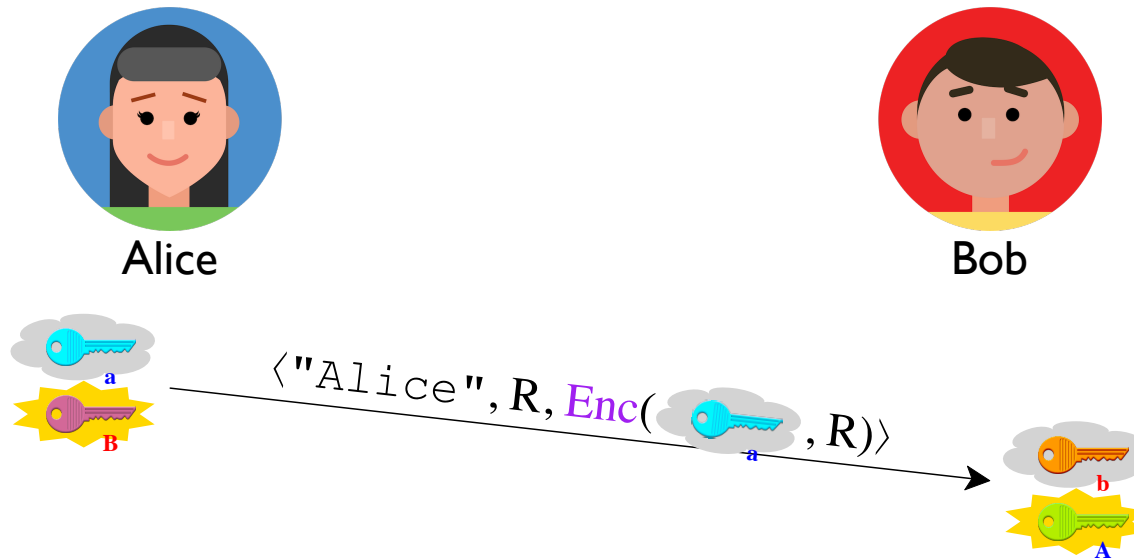
This is the replay-attack example, again

More Bad Authentication Ideas



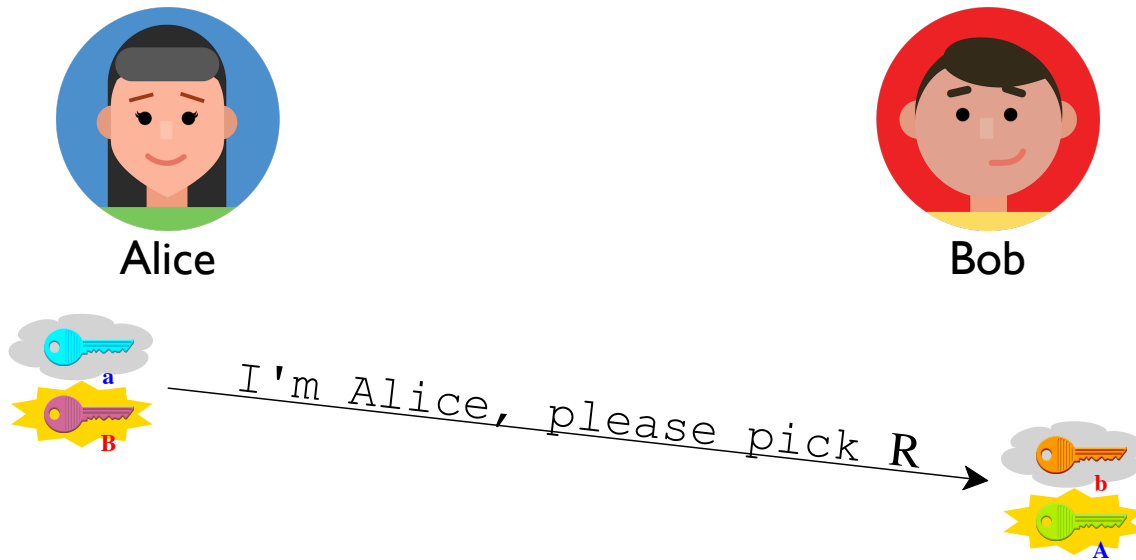
Still subject to a replay attack if Alice picks R

More Bad Authentication Ideas

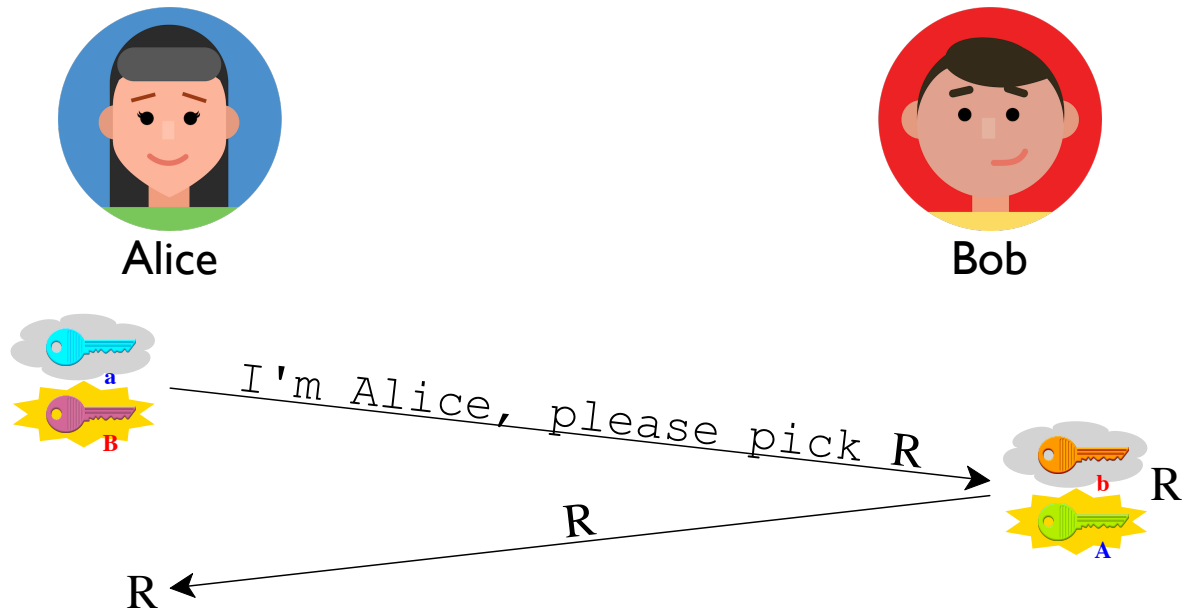


Still subject to a replay attack if Alice picks R

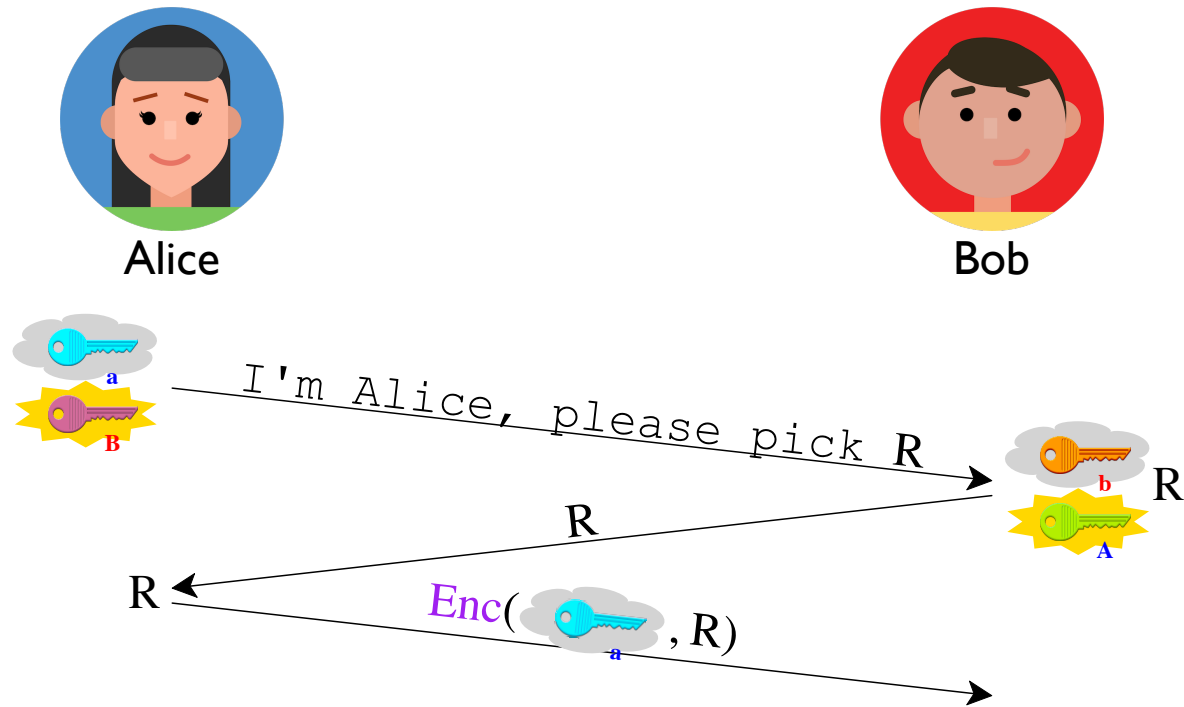
More Bad Authentication Ideas



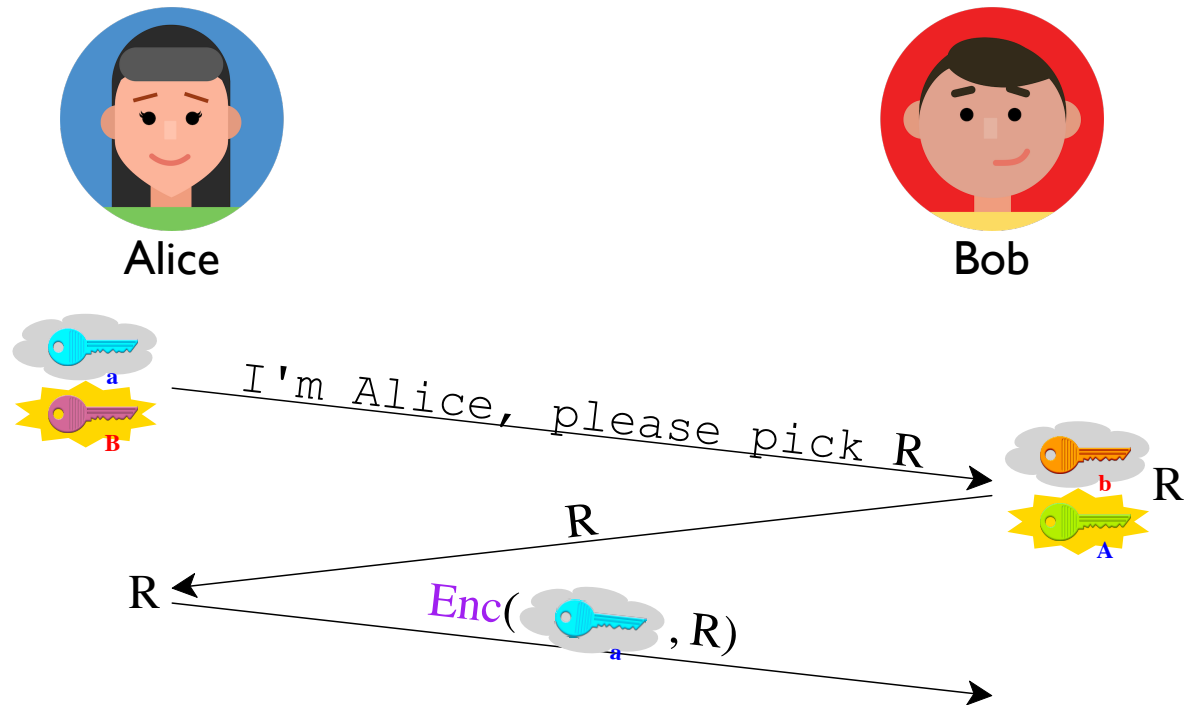
More Bad Authentication Ideas



More Bad Authentication Ideas



More Bad Authentication Ideas



Oops — Alice just signed an arbitrary R

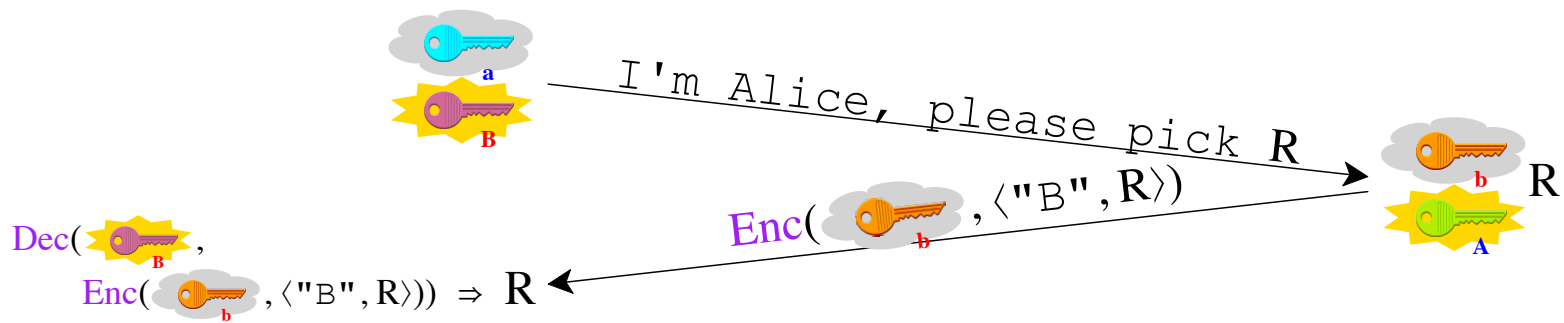
Authentication



Alice



Bob



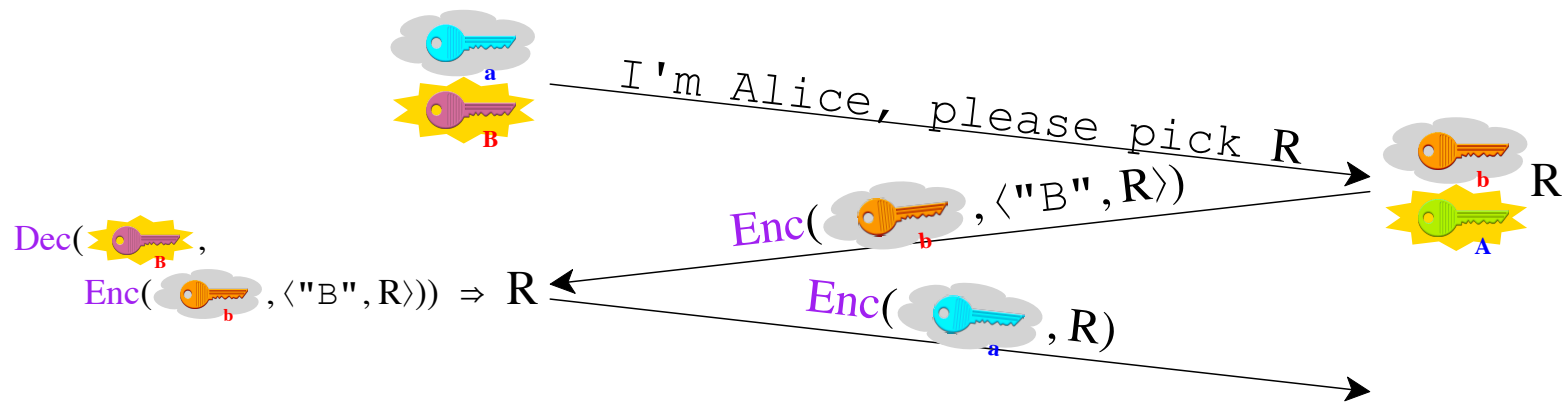
Authentication



Alice



Bob



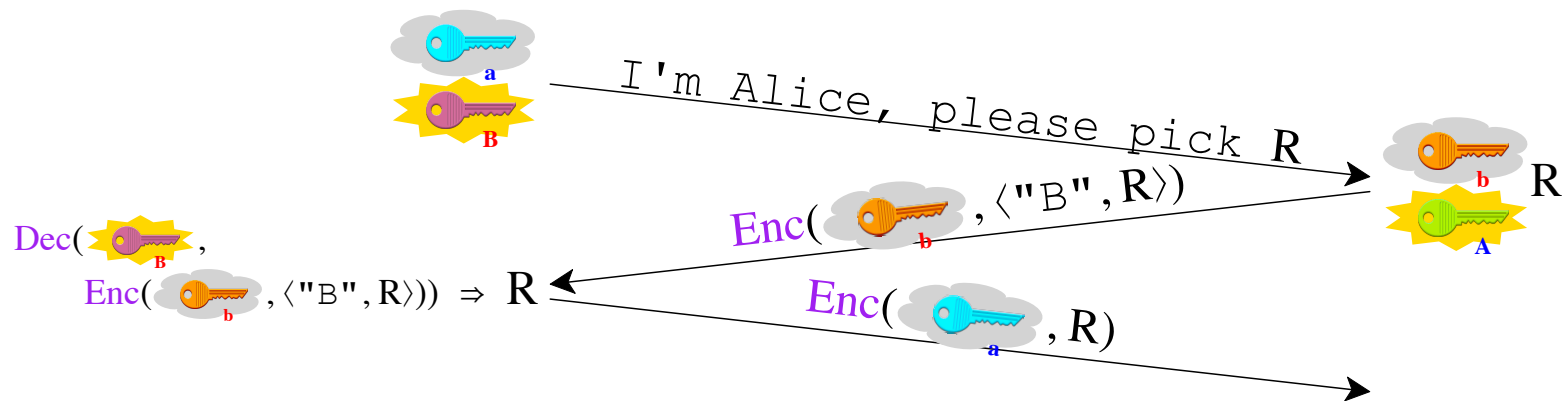
Authentication



Alice



Bob



Alice is signing R from Bob, but trusts Bob enough

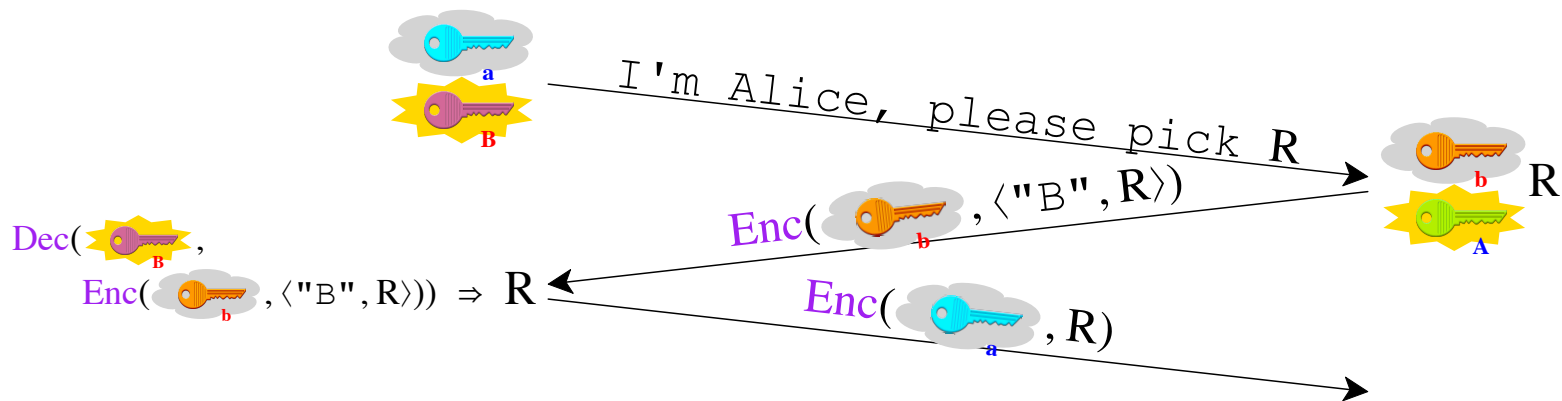
Authentication



Alice



Bob



SSH often works something like this, but not web sites

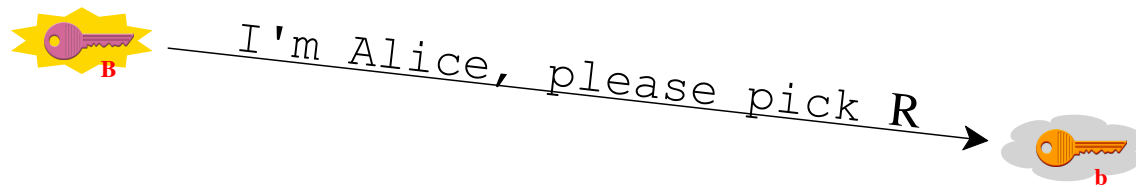
Authentication



Alice



Bob



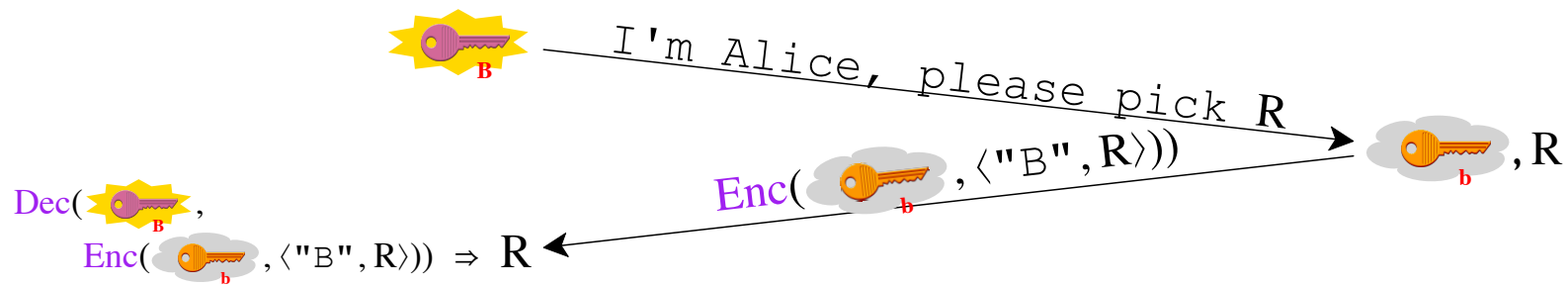
Authentication



Alice



Bob



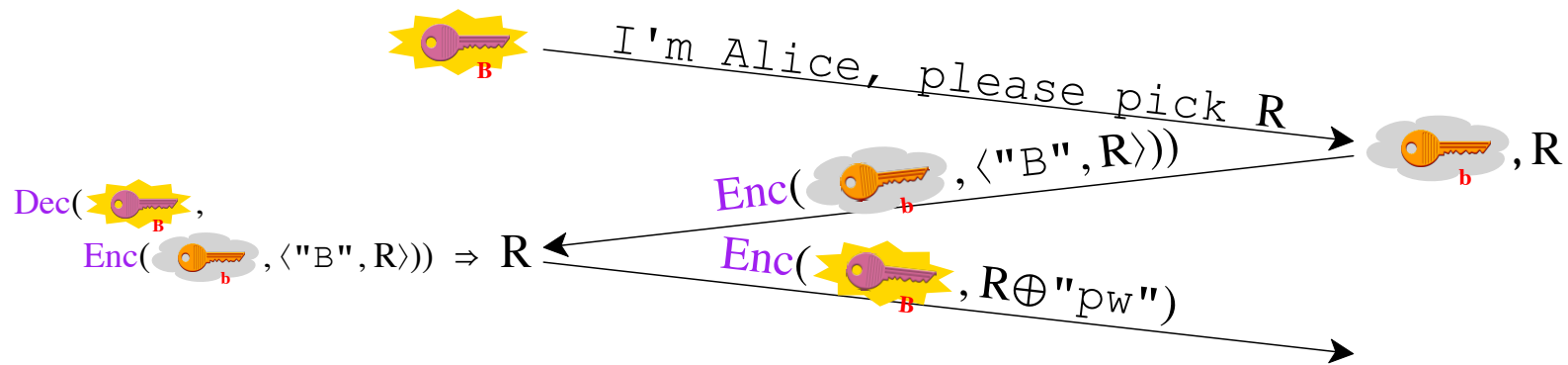
Authentication



Alice



Bob



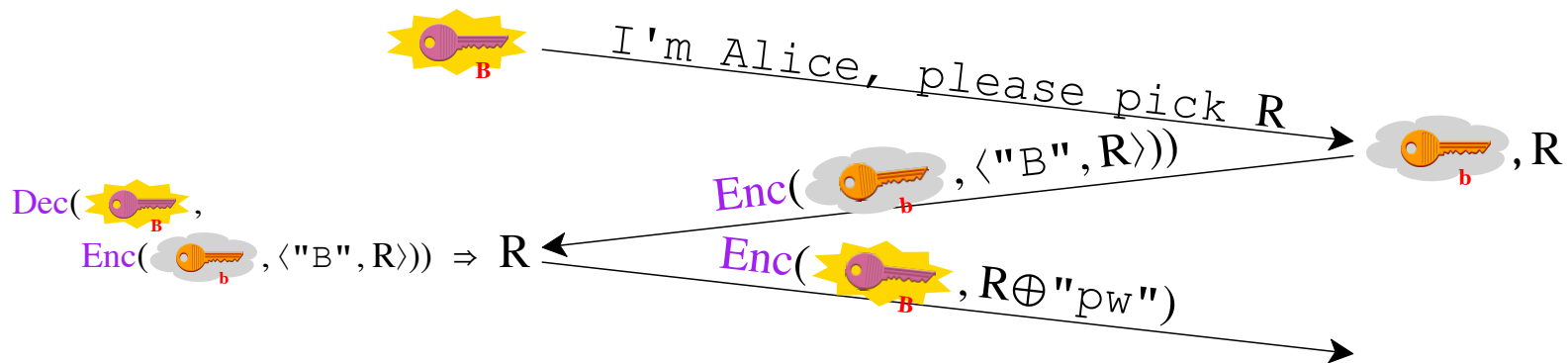
Authentication



Alice



Bob



As long as Bob picks a unique random R each time, this will work ok

... but there is still some room for improvement

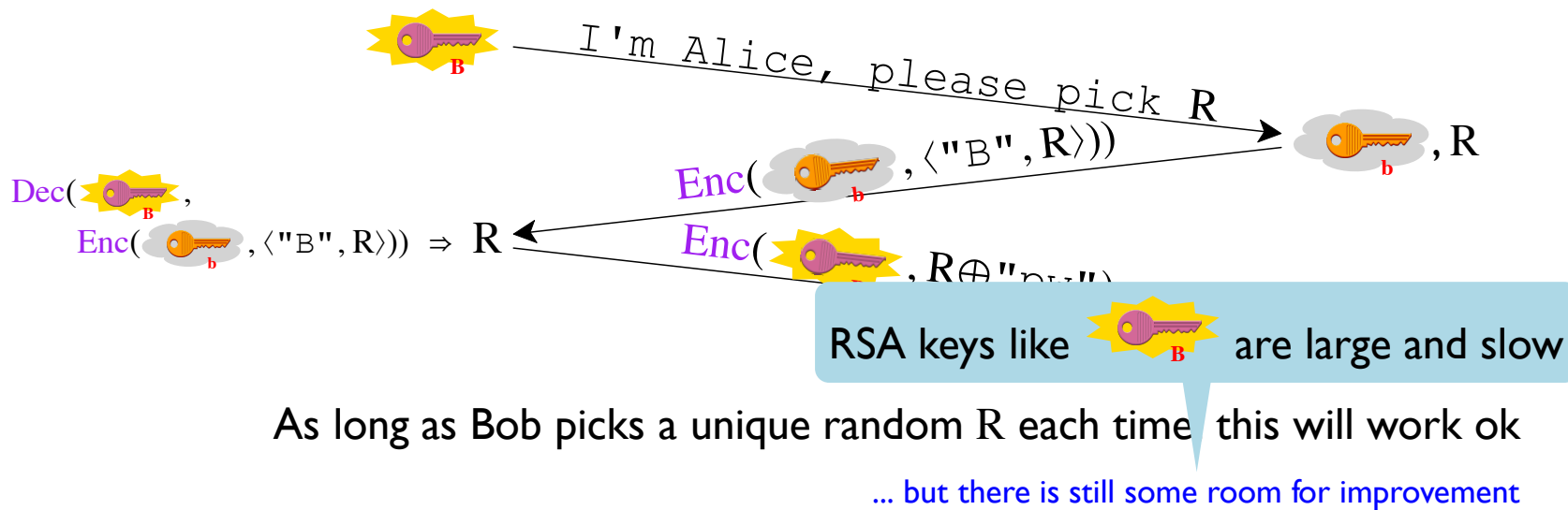
Authentication



Alice



Bob



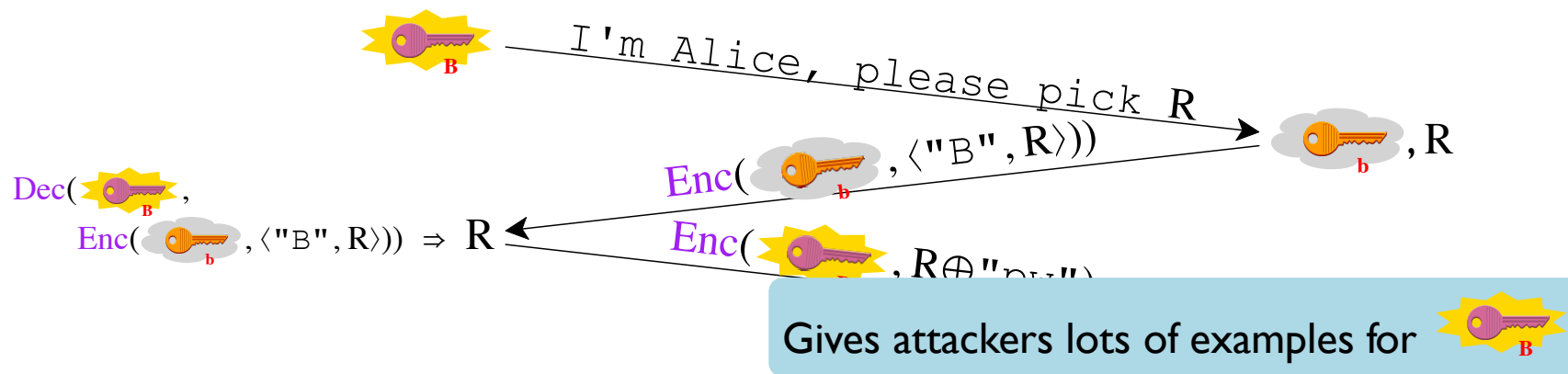
Authentication



Alice



Bob



As long as Bob picks a unique random R each time this will work ok

... but there is still some room for improvement

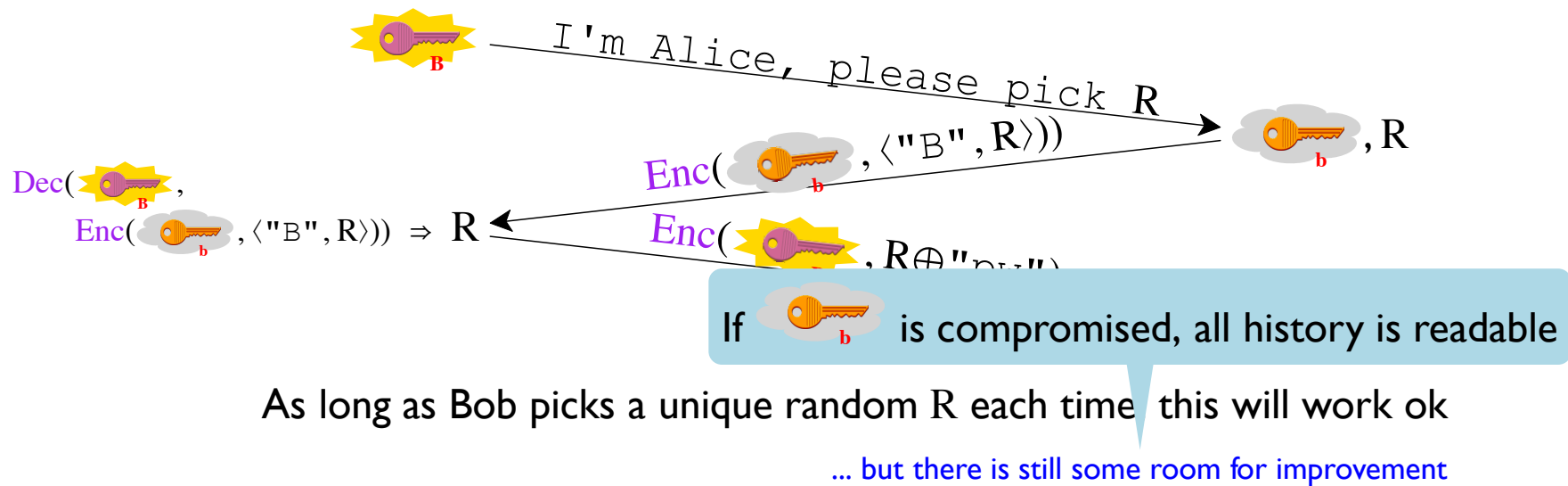
Authentication



Alice



Bob



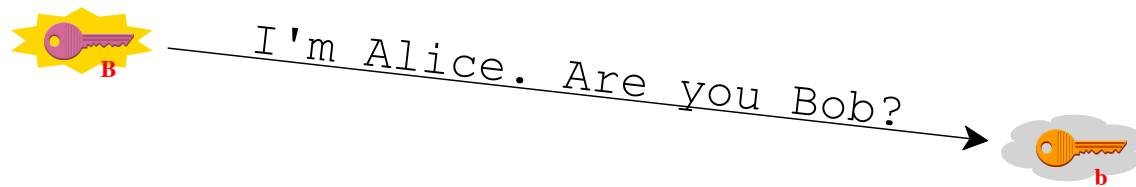
Session Keys



Alice



Bob



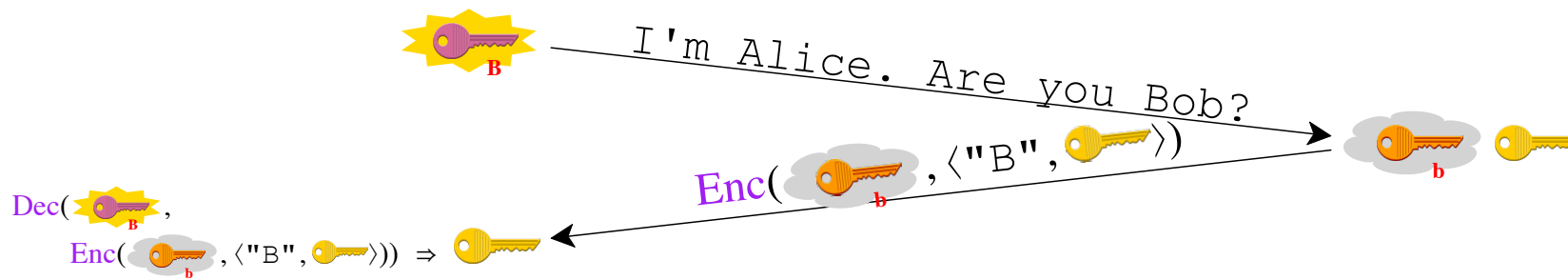
Session Keys



Alice



Bob



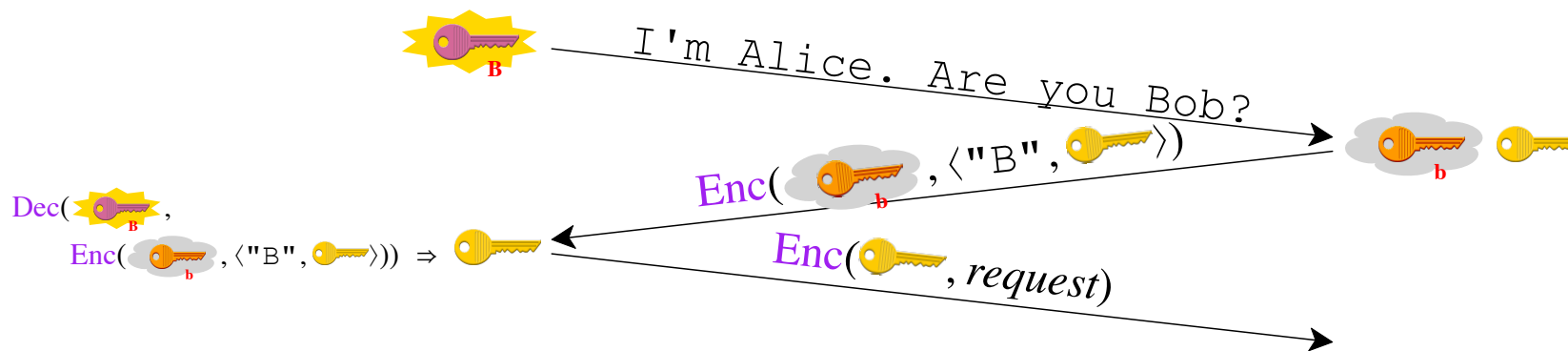
Session Keys



Alice



Bob



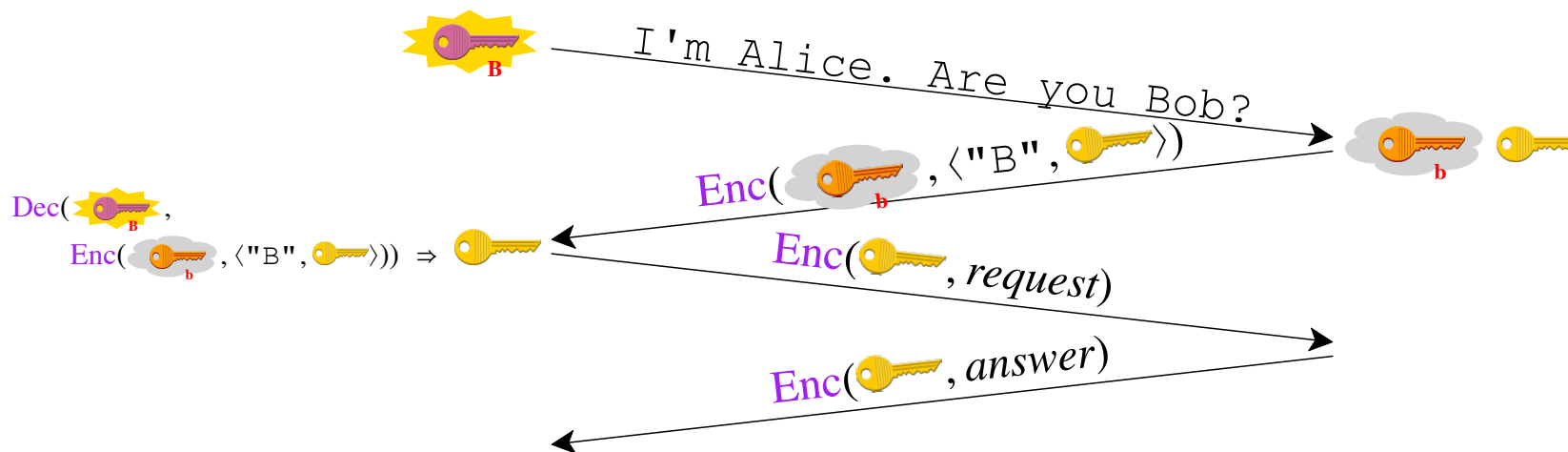
Session Keys



Alice



Bob



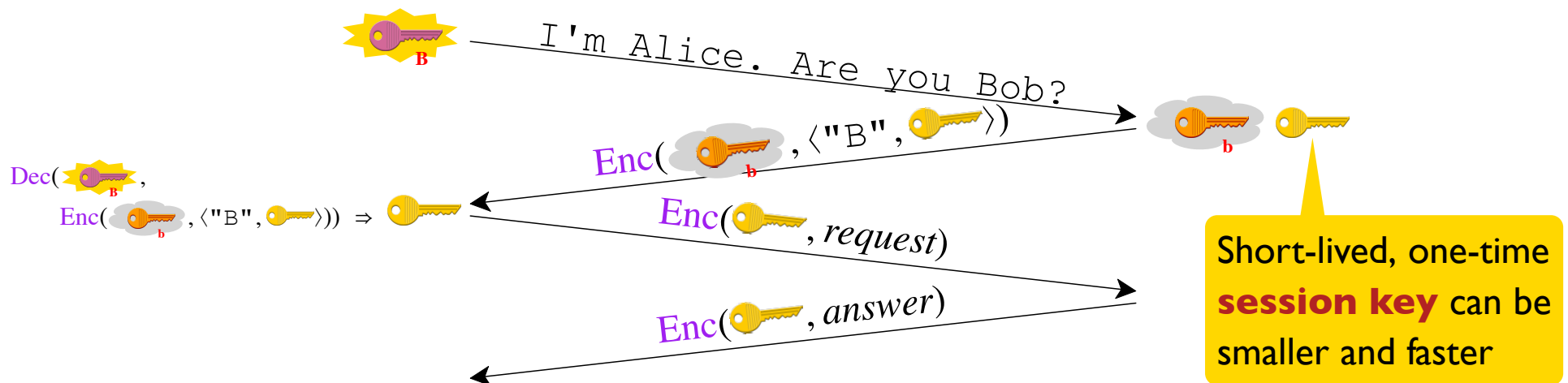
Session Keys



Alice



Bob



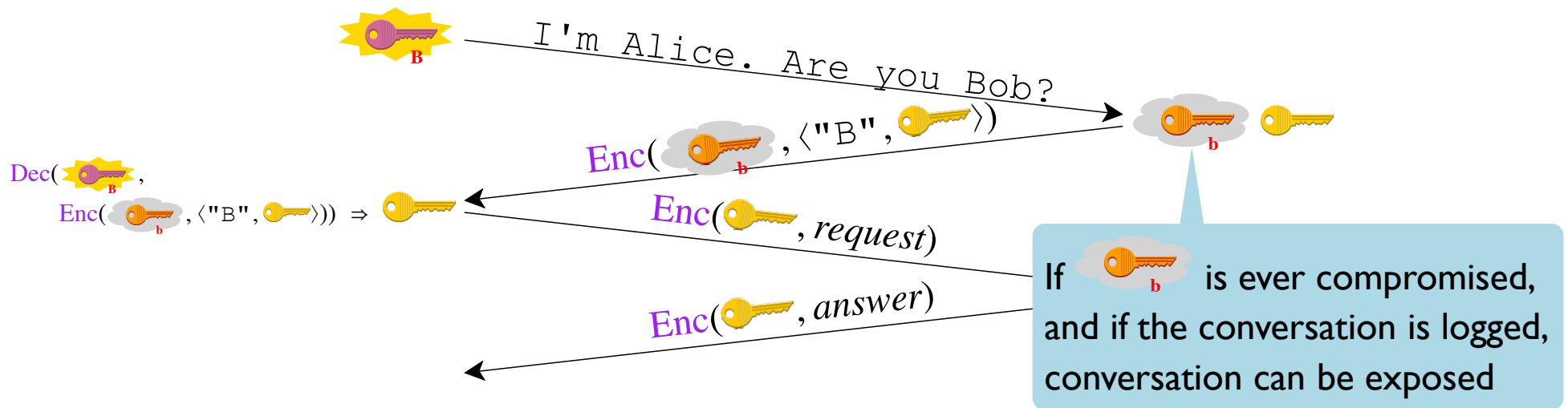
Session Keys



Alice



Bob



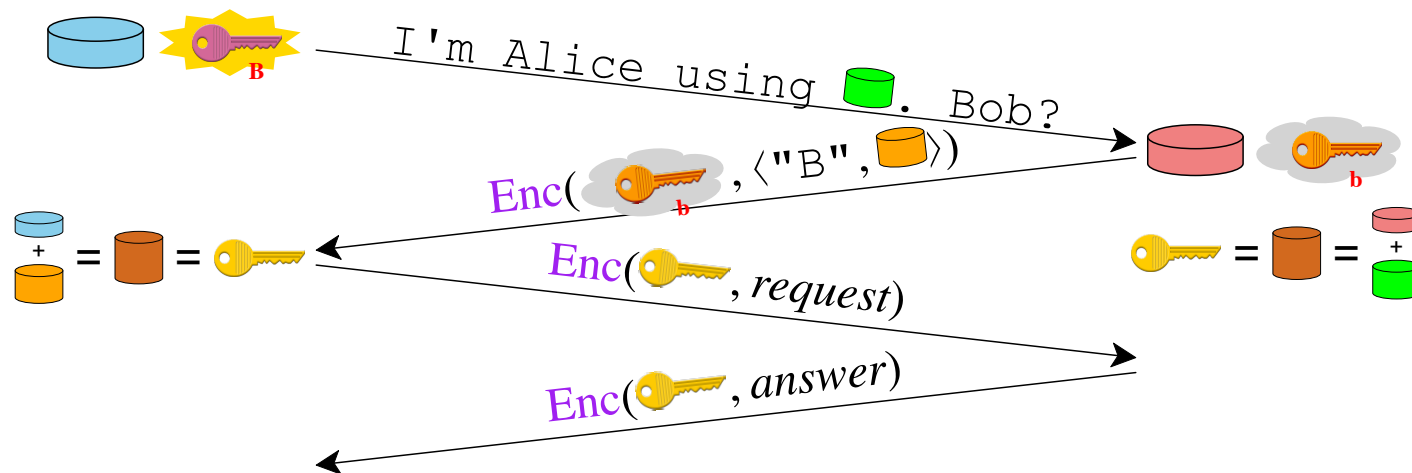
Session Keys with Perfect Forward Secrecy



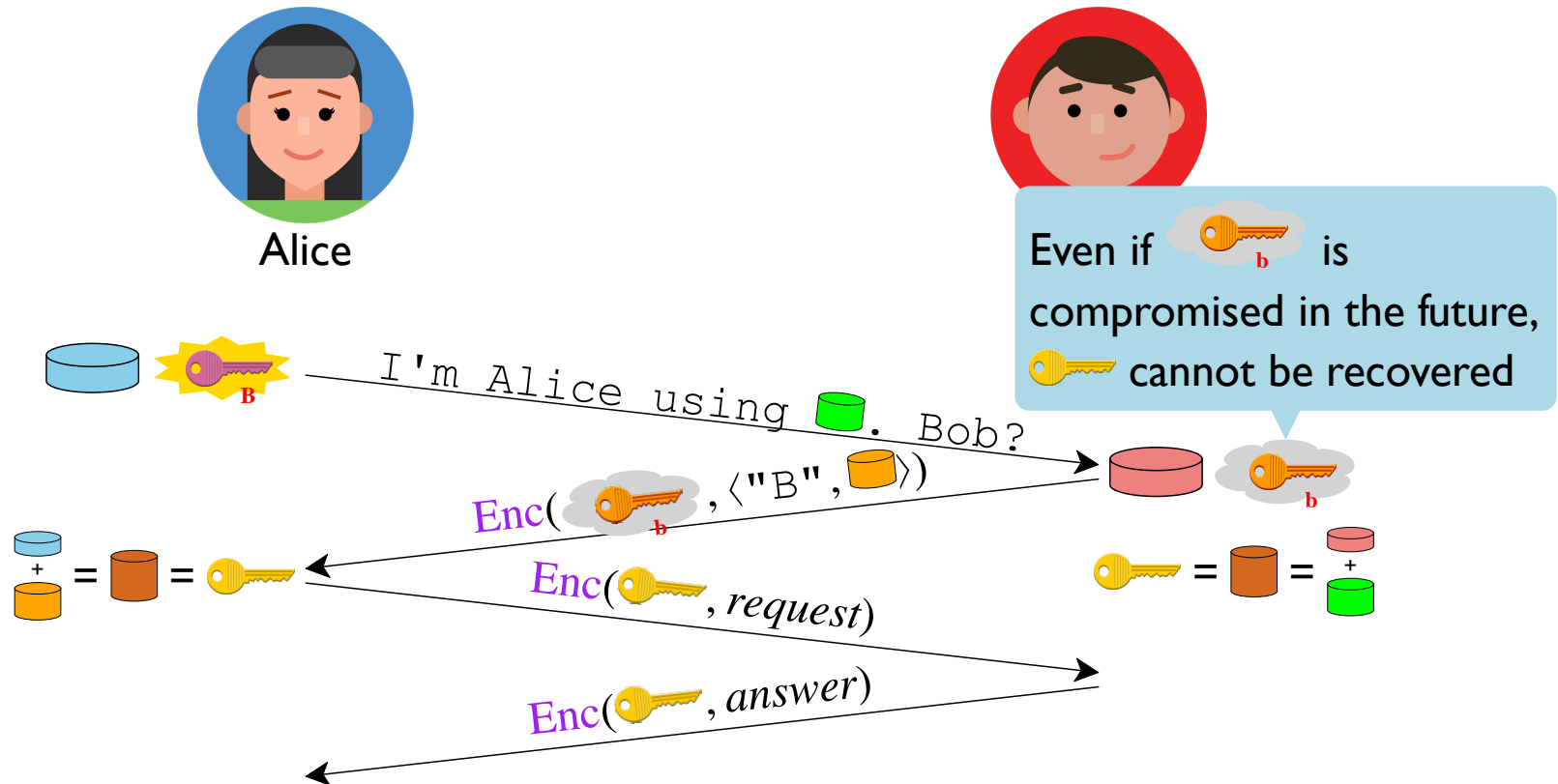
Alice



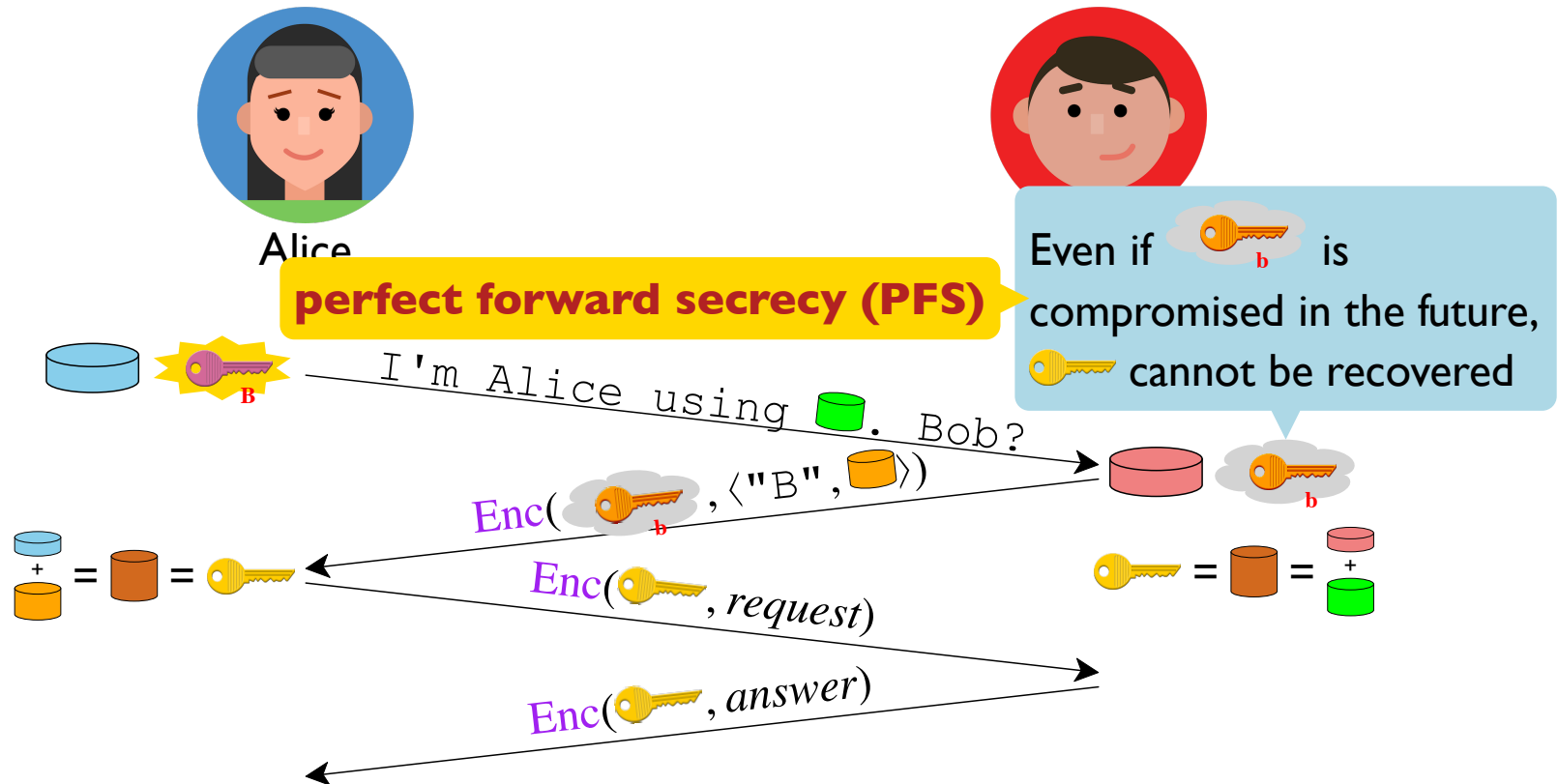
Bob



Session Keys with Perfect Forward Secrecy



Session Keys with Perfect Forward Secrecy



Summary

Authentication can mean logging in with a password, but it can also mean making sure that you're logging into the right server

We need our whole cryptography toolbox to get it right
... and there are still many pitfalls

Certificates and **certificate chains** address the problem of ahead-of-time sharing

Setting up **session keys** during authenticating is a best practice